# Stage-Aware Anomaly Detection Through Execution Flow Tracking

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#### Logging in Large Scale Systems

- Distributed storage systems of large-scale websites:
  - ▶ 1000s of servers
- Designed to be fault tolerant (automatic failure masking)
- Generate large volume of log data

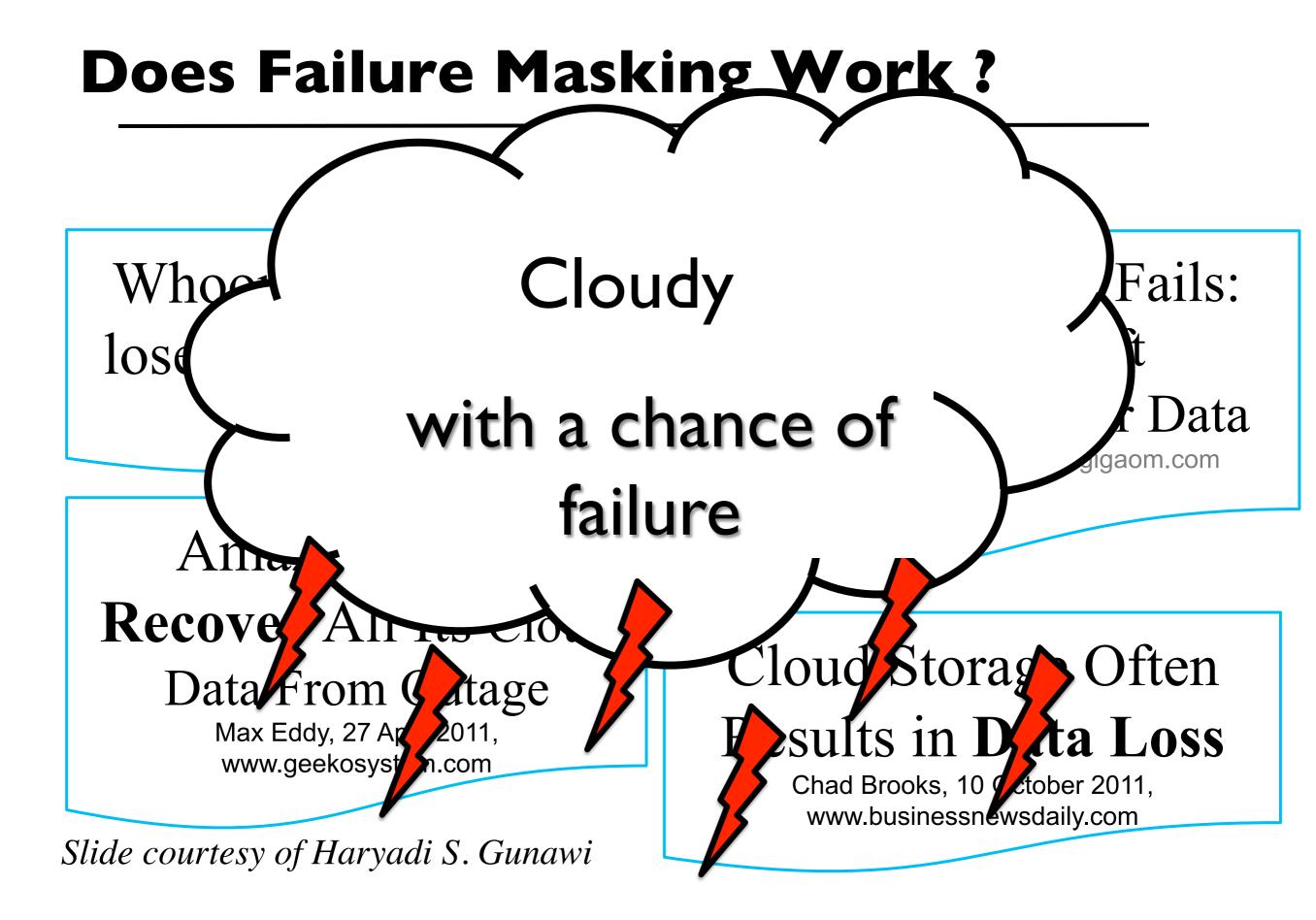






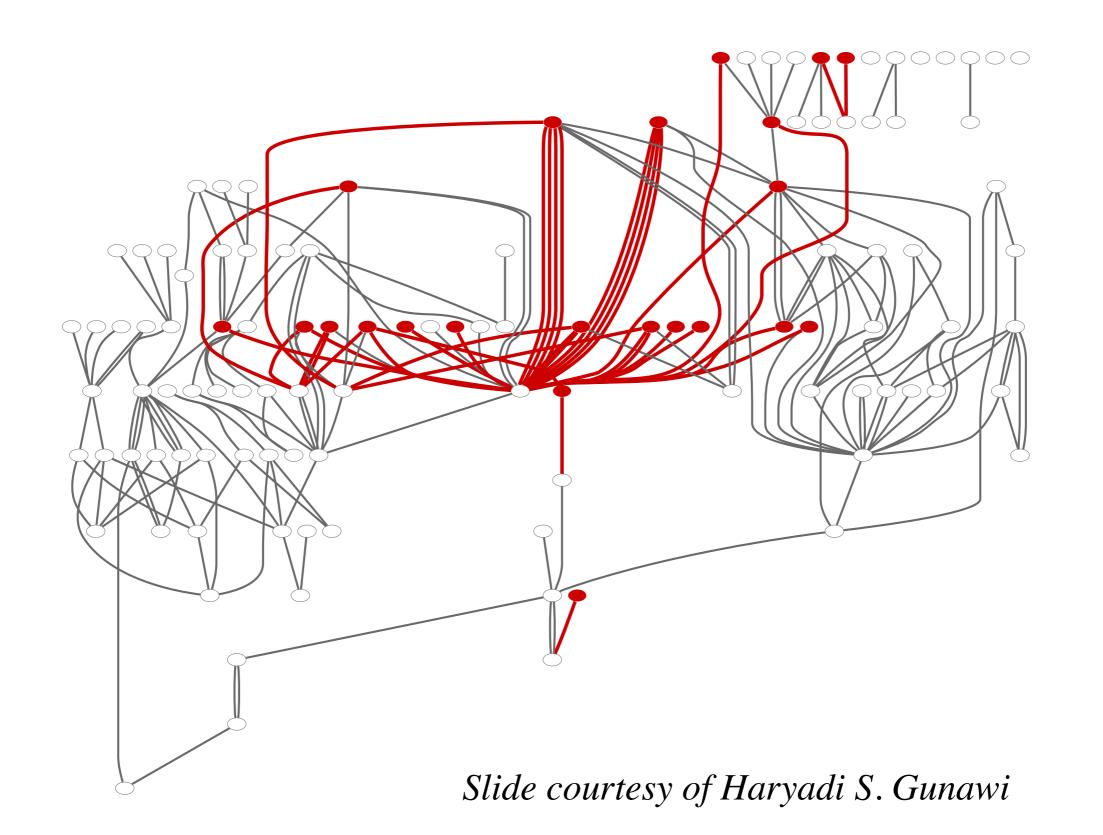
### Why Logging?

- We log detailed information for anomaly diagnosis
- Verbose logging: Calls, network, control flow, errors
- Goal: Understand call graph, execution flow, root cause analysis



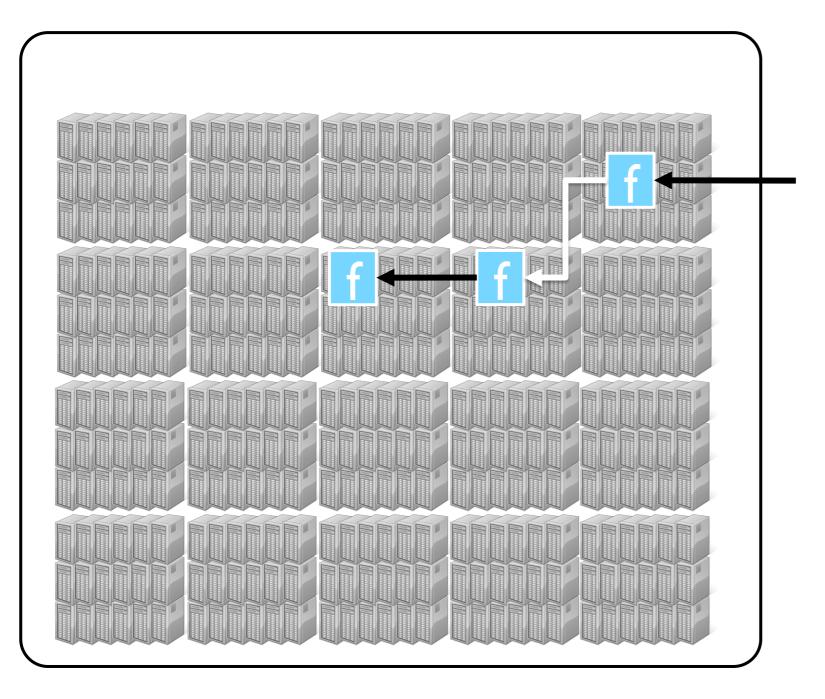
### Call Graph for FS

#### 37 errors out of 188 calls



### HDFS: 3-way Replication (3 nodes)

 Storage for Hadoop MapReduce Jobs

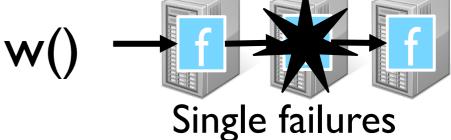


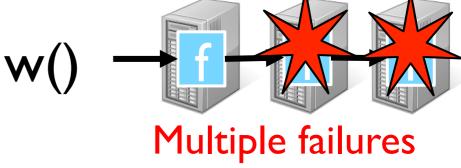
▶ Yahoo!: 4000 nodes

▶ Facebook: 2000 nodes

▶ Ebay: 700 nodes







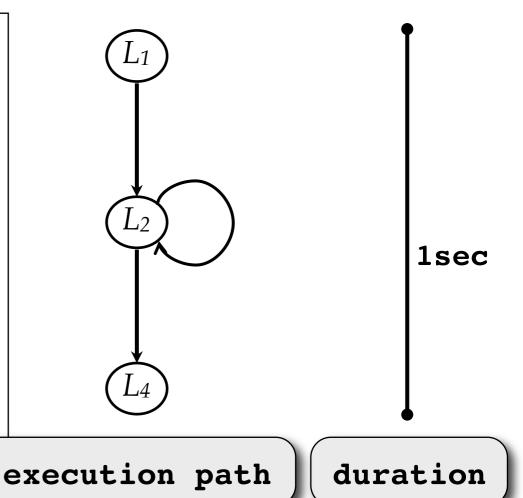
#### In An Ideal World ...



### In the Real World of Logs, Logs, ...

#### FileTransfer

```
LOG("Sending file "+file.name); //L1
int i=0;
while(i < file.size){
   int ack=send_recv(dest,file.block(i));
   if( ack == 0){
        i++;
        LOG("Sent block "+ i); //L2
   }else{
        LOG("Failed to send file "+ file.name); //L3
        return -1;
   }
}
LOG("Sent file "+ file.name); //L4
return 0;</pre>
```



Logs

1:10.032 Sending file A

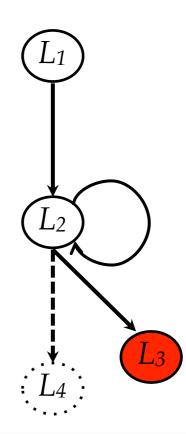
Logs contain information about the

execution flow and duration

#### Logs Contain Detailed Information

#### FileTransfer

```
LOG("Sending file "+file.name); //L1
int i=0;
while(i < file.size){
   int ack=send_recv(dest,file.block(i));
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return 0;</pre>
```



execution path

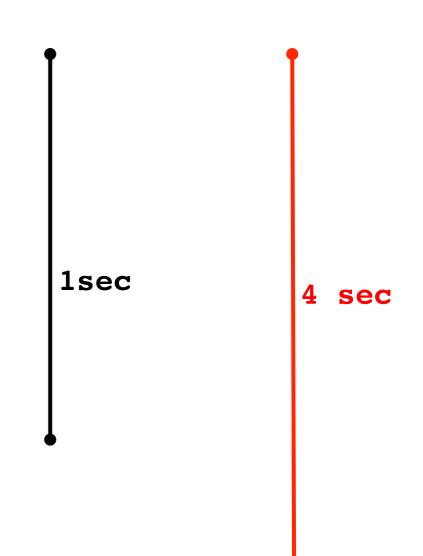
#### Logs

```
1:10.023 Sending file F
1:10.042 Sent block 1
1:10.099 Sent block 2
....
1:11.130 Sent block 300
1:11.131 Failed to send file F
```

### Logs Contain Detailed Information

#### FileTransfer

```
LOG("Sending file "+file.name); //L1
int i=0;
while(i < file.size){
   int ack=send_recv(dest,file.block(i));
   if( ack == 0){
      i++;
      LOG("Sent block "+ i); //L2
   }else{
      LOG("Failed to send file "+ file.name); //L3
      return -1;
   }
}
LOG("Sent file "+ file.name); //L4
return 0;</pre>
```



duration

#### Logs

```
1:10.032 Sending file B
1:10.042 Sent block 1
1:10.099 Sent block 2
....
1:14.030 Sent block 1000
1:14.032 Sent file B
```

### Inferring Execution Flow from Logs

```
1:10.023 Sending file
1:10.042 Sent block 1
1:10.099 Sent block 2
1:11.130 Sent block 300
1:11.131 Failed to send file F
                                                   execution path
                                                                            duration
```

#### Challenges

- Identify tasks from logs
- Distinguish normal vs. anomalous tasks
- Overheads
  - Storage and processing

### Challenge 1: Identify tasks from logs

#### Normal Execution Flow

```
1:10.023 Start sending file A
1:10.042 Sent block 1
1:10.099 Sent block 2
....
1:11.130 Sent block 1000
1:11.132 Sent file A
```

#### **Anomalous Execution Flow**

```
1:10.033 Start sending file F
1:10.032 Sent block 1
1:11.049 Sent block 2
....
1:11.050 Sent block 300
1:11.131 Failed to send file F
```

```
1:10.023 Start sending file A
1:10.033 Start sending file F
1:10.032 Sent block 1
1:10.042 Sent block 1
1:11.049 Sent block 2
1:10.099 Sent block 2
...
1:11.130 Sent block 1000
1:11.131 Failed to send file F
1:11.132 Sent file A
```

Server runs forever, how to delimit tasks?

### Challenge 1: Identify tasks from logs

#### Normal Execution Flow

```
1:10.023 Start sending file A
1:10.042 Sent block 1
1:10.099 Sent block 2
....
1:11.130 Sent block 1000
1:11.132 Sent file A
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#### **Anomalous Execution Flow**

```
1:10.033 Start sending file F
1:10.032 Sent block 1
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```
1:10.023 Start sending file A
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...
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1:11.131 Failed to send file F
1:11.132 Sent file A
```

## Thread reuse, data flow tracking may obscure task execution flow

### Challenge 2: Normal vs Anomaly

- No baseline
- Anomalies do not always generate an explicit error/warning log message

How to identify normal vs. rare execution flows?

#### Challenge 3: Overhead

- Log messages are text to be read by human
- Verbose logging generates huge log volume!
  - 2600 times more than INFO-level (default) logging
- Text processing is expensive and imprecise

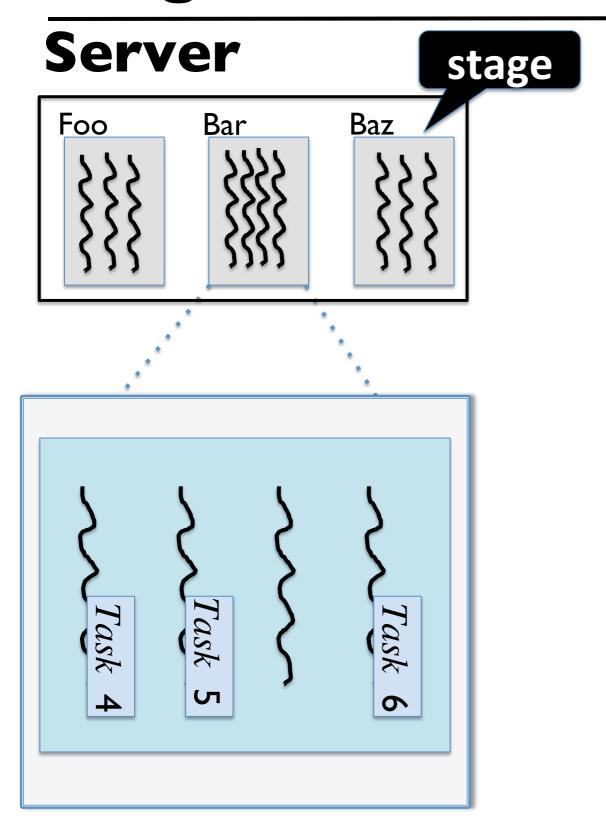
# 2.4TB per day for a 100 node Cassandra cluster Not many people can look into a 2TB log file

### **Key Observation**

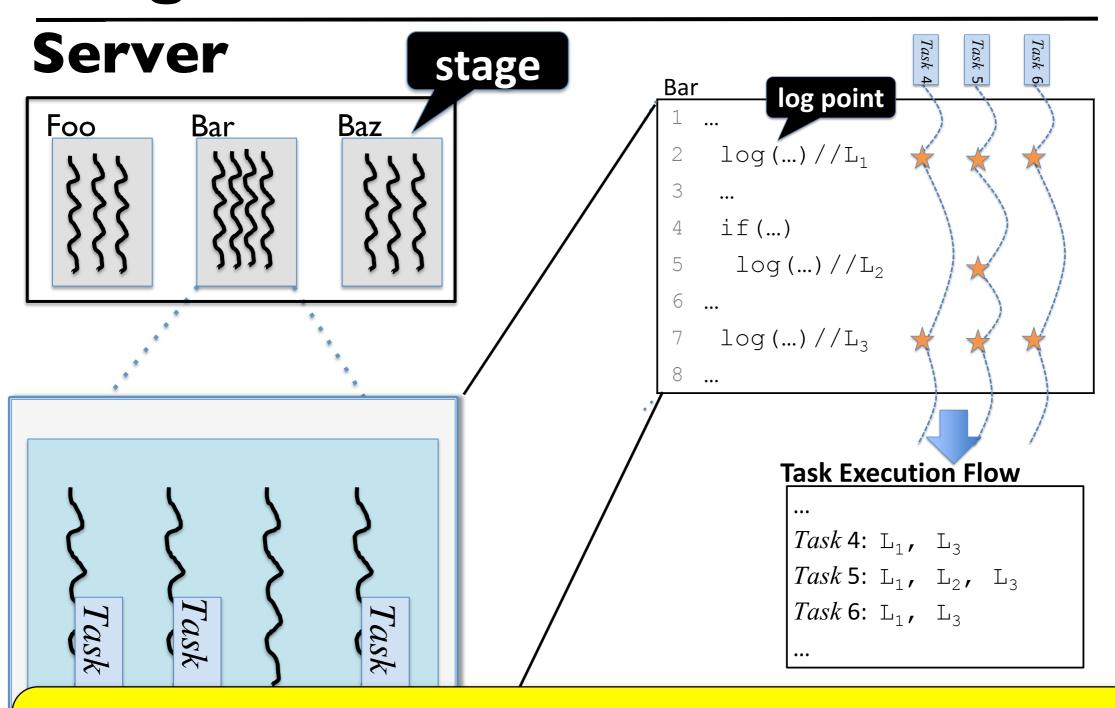
Many server codes have a modular or stage-based architecture

- Stage-aware anomaly detection
  - ▶ on-the-fly
  - with low overhead
  - almost completely automated

### Staged-Architecture

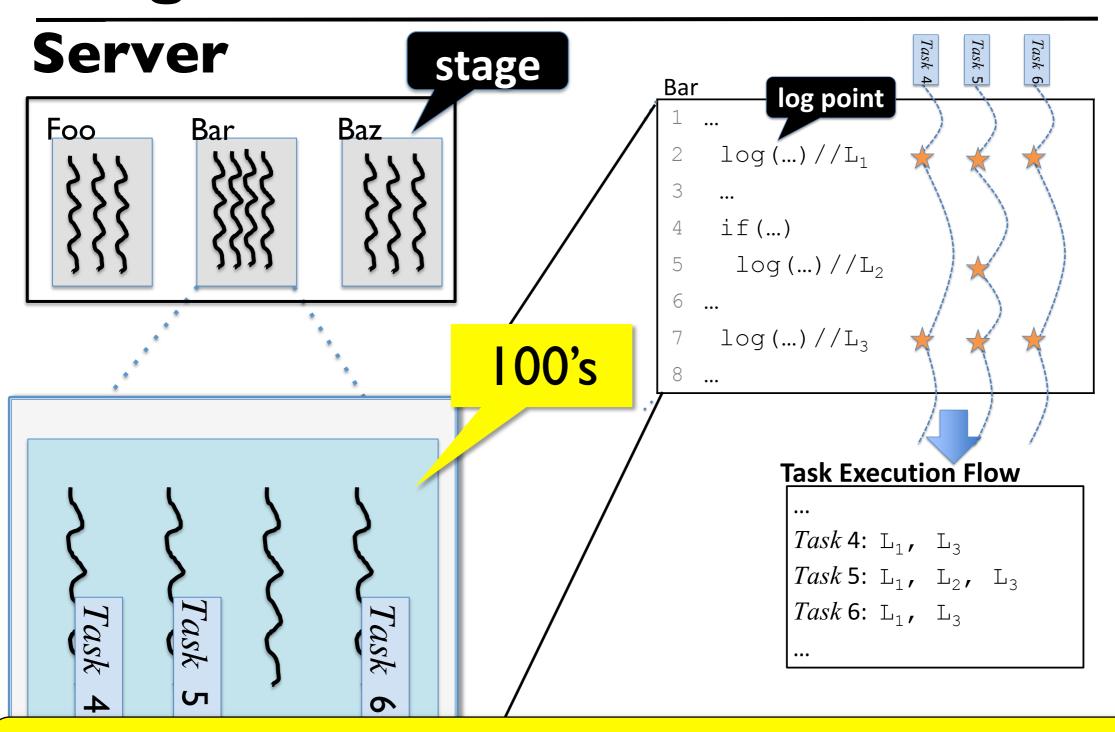


#### Staged-Architecture



We capture execution flow from log points onthe-fly without generating logs

#### Staged-Architecture



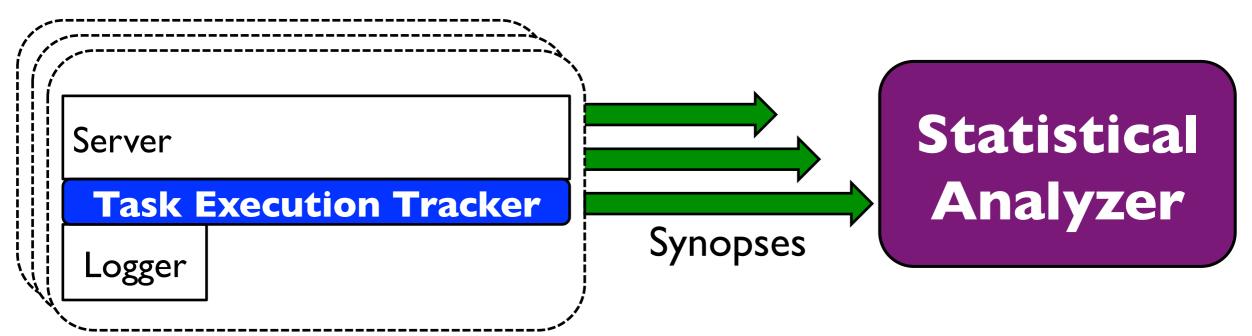
We exploit similarity between tasks for statistical anomaly detection

### Stage-Aware Anomaly Detection

- Leverage the staged code structure
  - ▶ To track start and end of each task
- Log statements as trace points
  - ▶ To track execution flow of the tasks
- Exploit the statistical similarity of tasks to detect
  - Flow anomalies: rare execution paths
  - Performance anomalies: unusually high duration

#### **Prototype**

#### Hosts

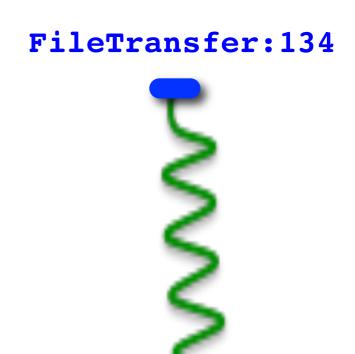


- Task Execution Tracker on each sever
  - tracks log points encountered by each task
  - generates synopses of task execution flows
  - streams to Statistical Analyzer for real-time anomaly detection

#### Instrumentation

#### FileTransfer

```
Tracker.setContext("FileTransfer", task_id);
LOG("Sending file "+file.name); //L1
int i=0;
while(i < file.size){
  int ack=send_recv(dest,file.block(i));
  if( ack == 0){
    i++;
    LOG("Sent block "+ i); //L2
  }else{
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    return -1;
  }
}
LOG("Sent file "+ file.name); //L4
return 0;</pre>
```

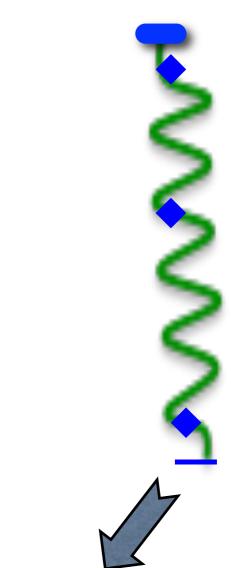


### **Tracking Log Points**

#### FileTransfer

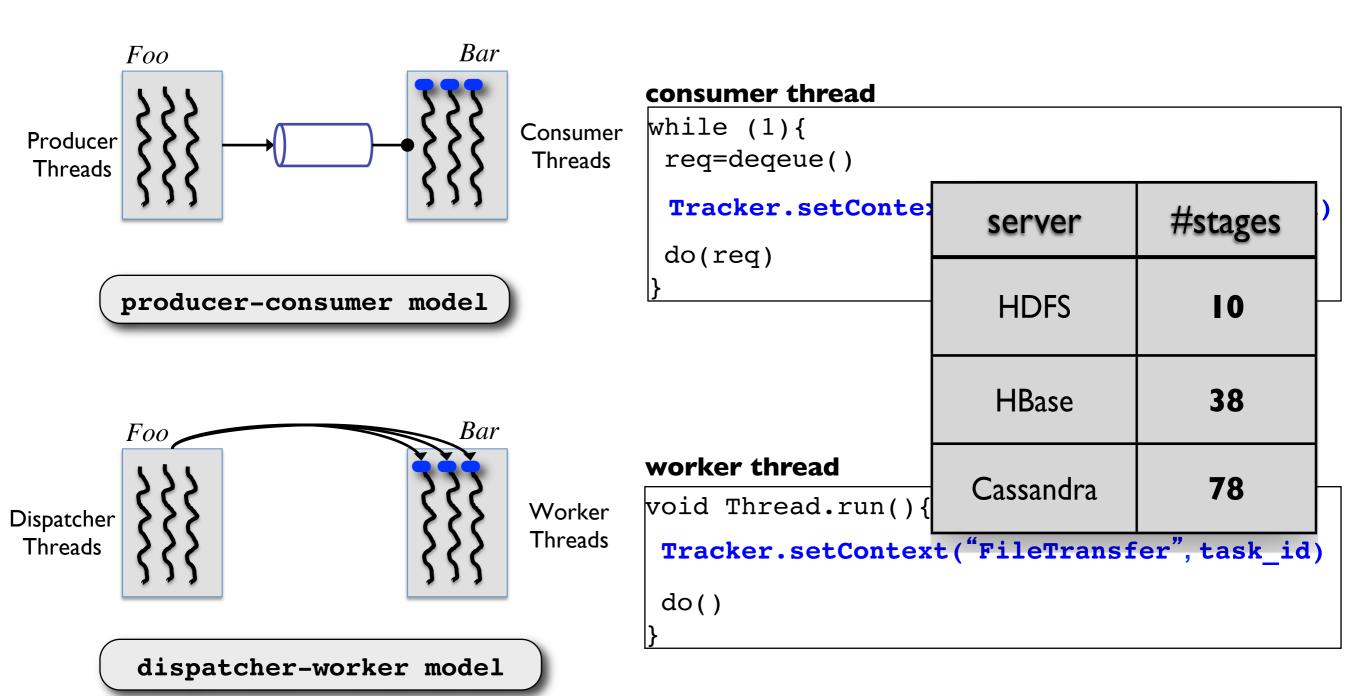
```
Tracker.setContext("FileTransfer", task_id)
LOG(1, "Sending file "+file.name); //L1
int i=0;
while(i < file.size){
   int ack=send_recv(dest,file.block(i));
   if( ack == 0){
        i++;
        LOG(2, "Sent block "+ i); //L2
   }else{
        LOG(3, "Failed to send file "+ file.name); //L3
        return -1;
   }
}
LOG(4, "Sent file "+ file.name); //L4
return 0;</pre>
```

#### FileTransfer:134

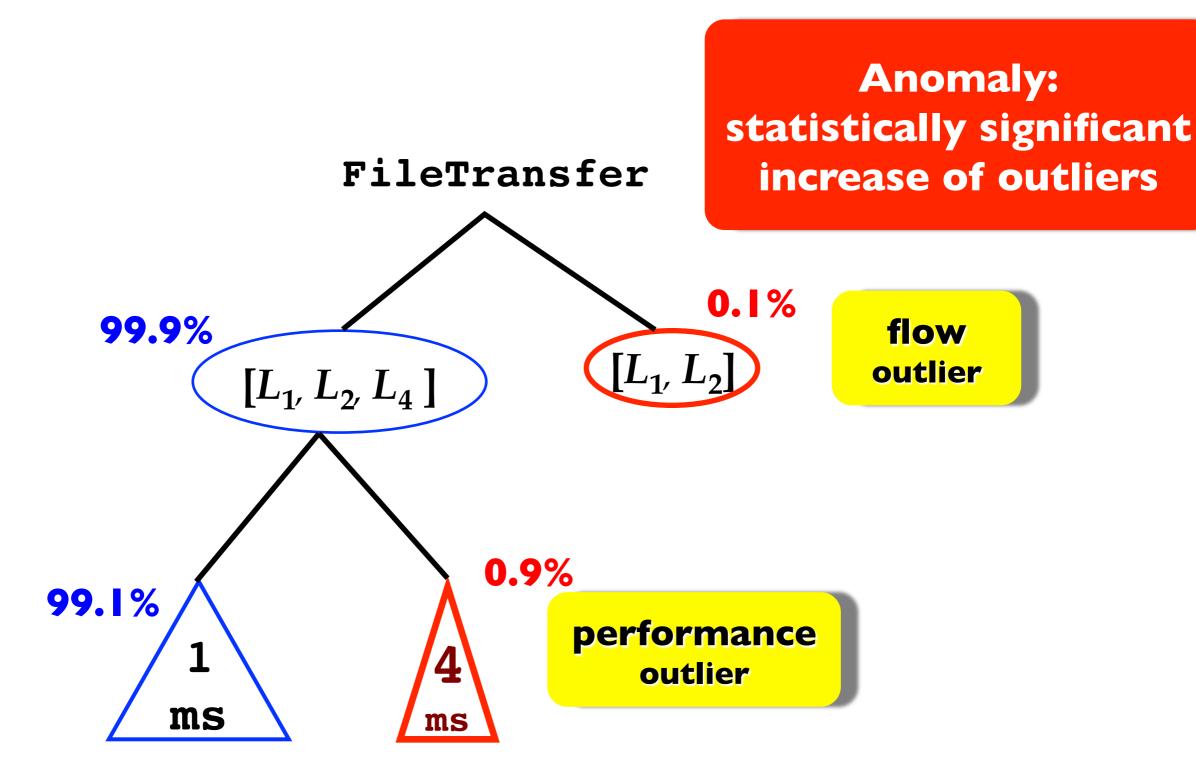


**Synopsis**: <134, "FileTransfer", 3.5 ms,  $[L_1, L_2, L_4]$ >

#### **Automatic Instrumentation**



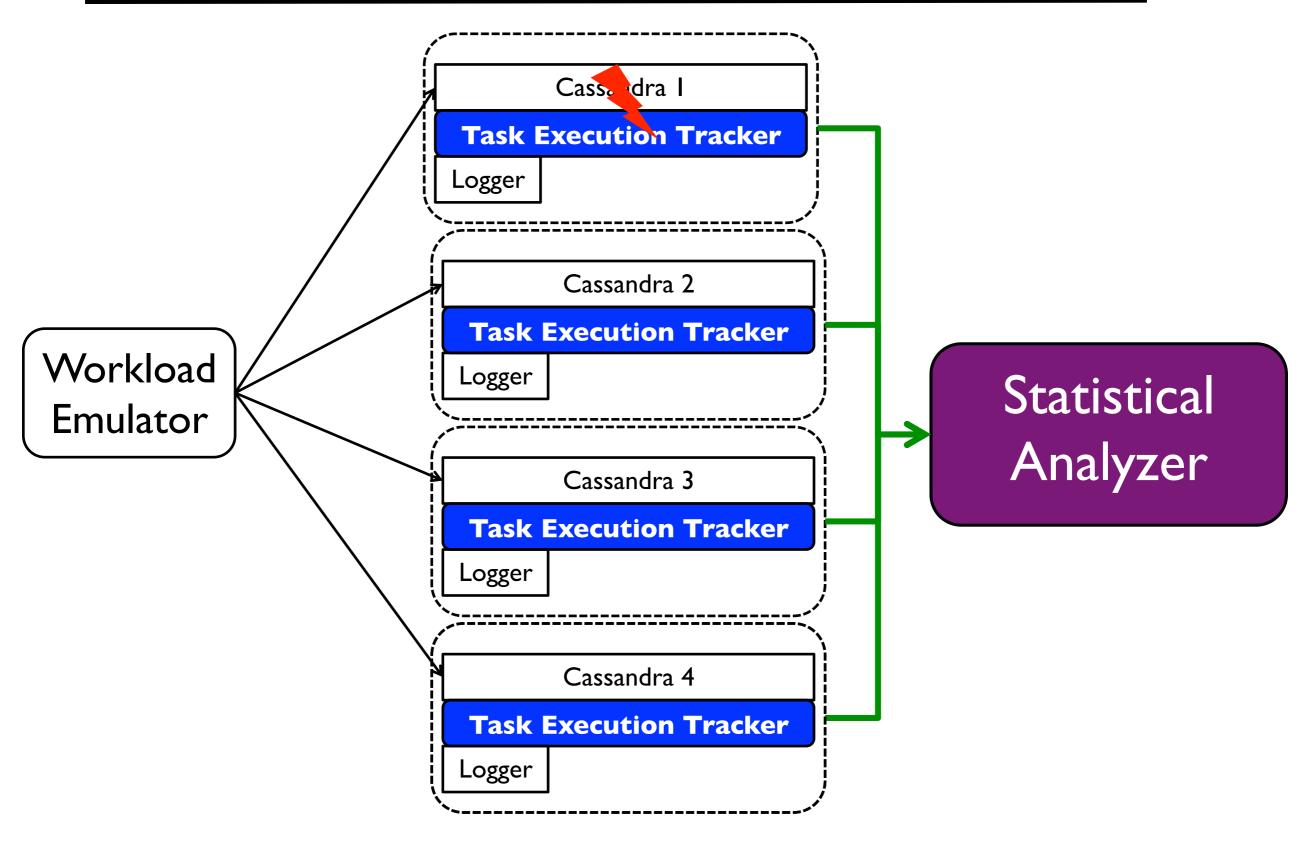
### Model Building (per Stage)



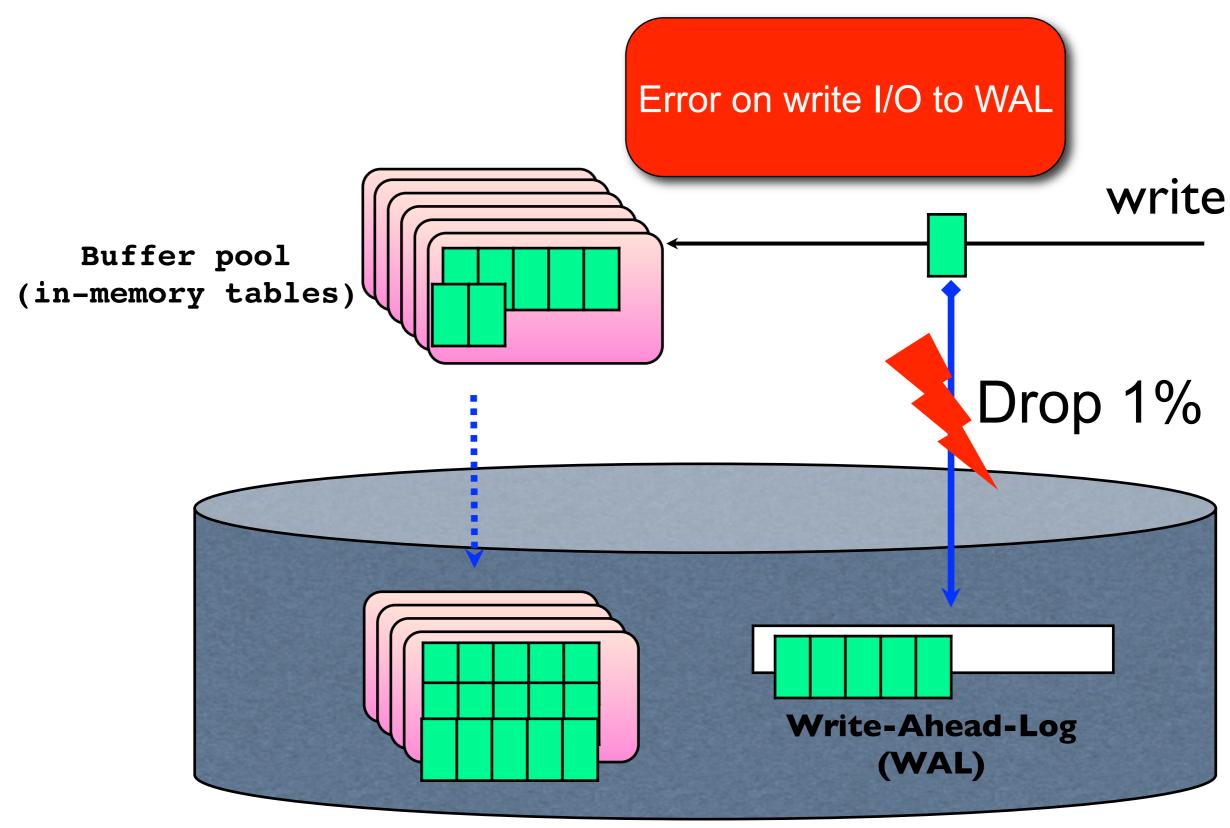
#### **Evaluation**

- •Three distributed storage systems:
  - ▶ Hadoop Distributed Filesystem (HDFS)
  - ▶ HBase
  - **▶** Cassandra
- Write intensive workload of Yahoo Cloud Serving Benchmark (YCSB)
  - ▶ Similar to real-world scenarios
  - Stress the system

### **Experiment Setup**



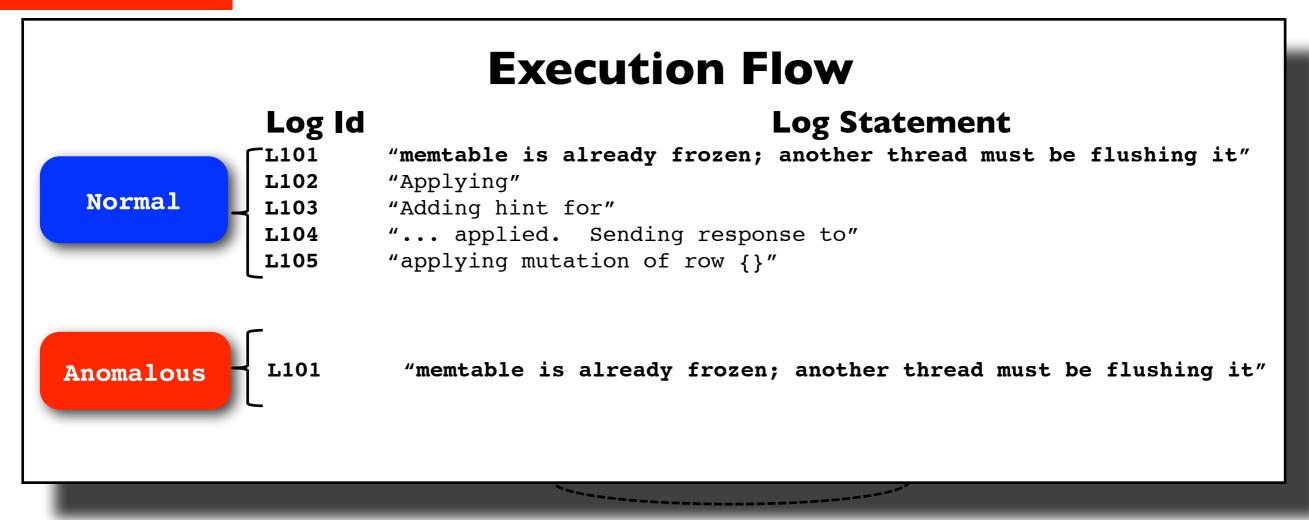
#### Cassandra write I/O path

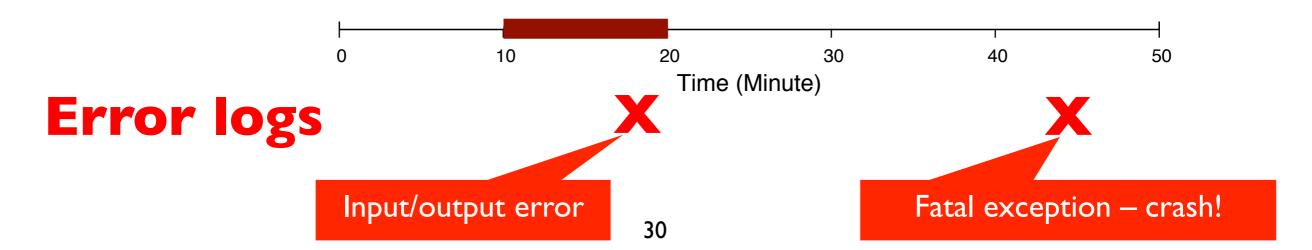


### Error on write I/O ()

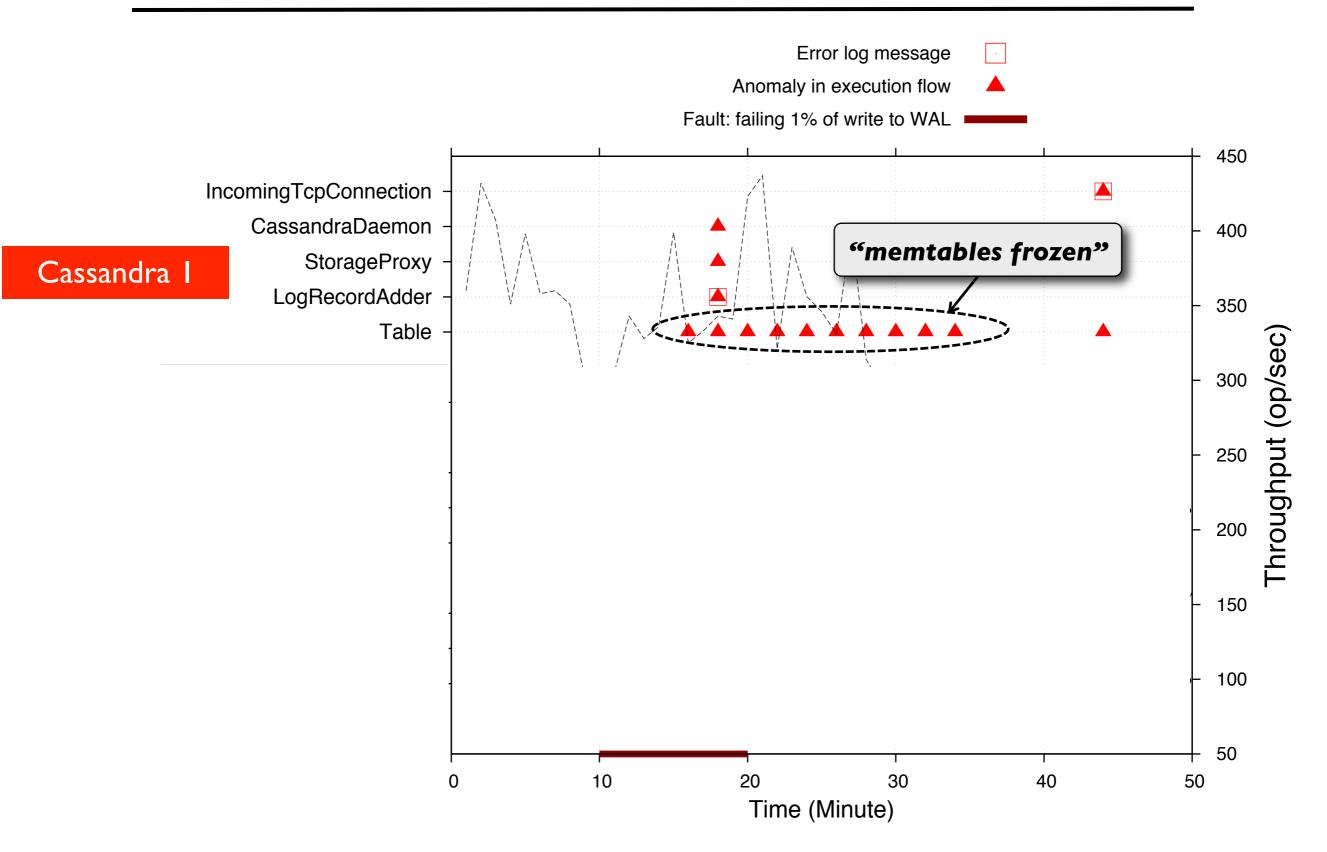
Fault: failing 1% of write to WAL

#### Cassandra I





#### Error on write I/O



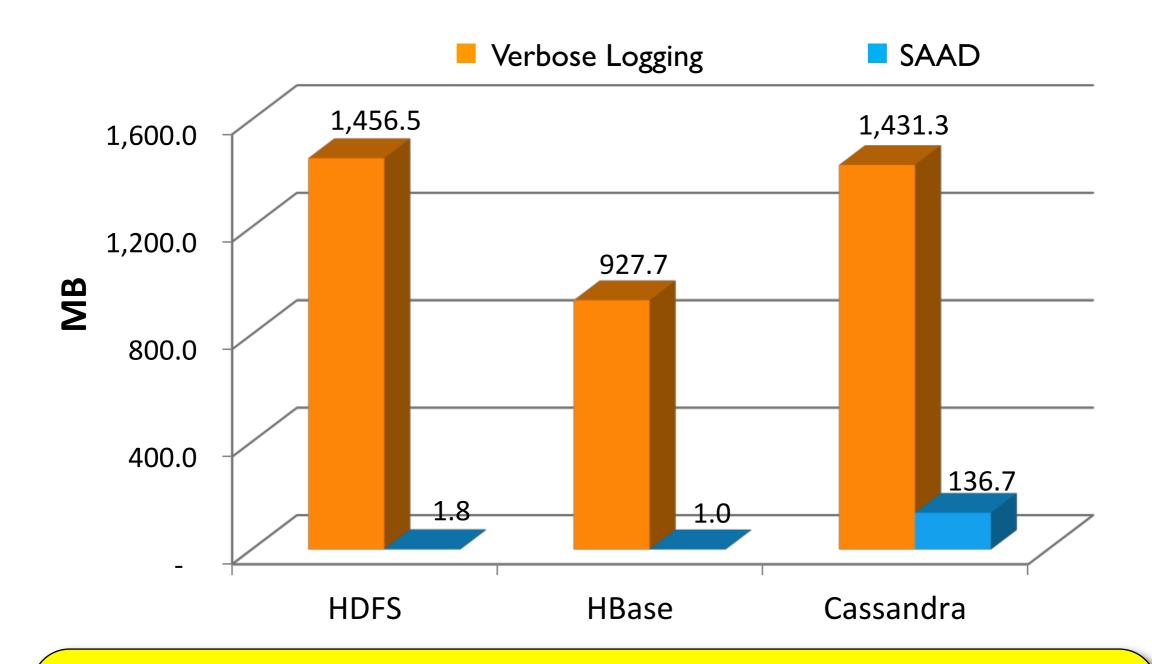
#### **Uncovering Bugs & Misconfigurations**

#	Туре	Component	Description
1	Bug	HDFS Data Node	Empty Packet
2	Bug	Hbase Regionserver	Distributed Log splitting gets stuck
3	Misconfig.	HBase Regionserver	No live nodes contain current block
4	Misconfig.	Hbase Regionserver	Zookeeper missed heartbeat due to lengthy GC

### **Analysis Cost**

- •State of the Art, based on regular expression matching:
  - Offline processing of 1.6GB log data
    - 12 millions log messages
  - ▶ 12 Minutes on 8 cores full utilization
- Our solution
  - ▶ Real time on one core
  - ▶ Average of 3% CPU utilization

#### **Storage Overhead**



Up to 1000x storage saving vs. verbose logging

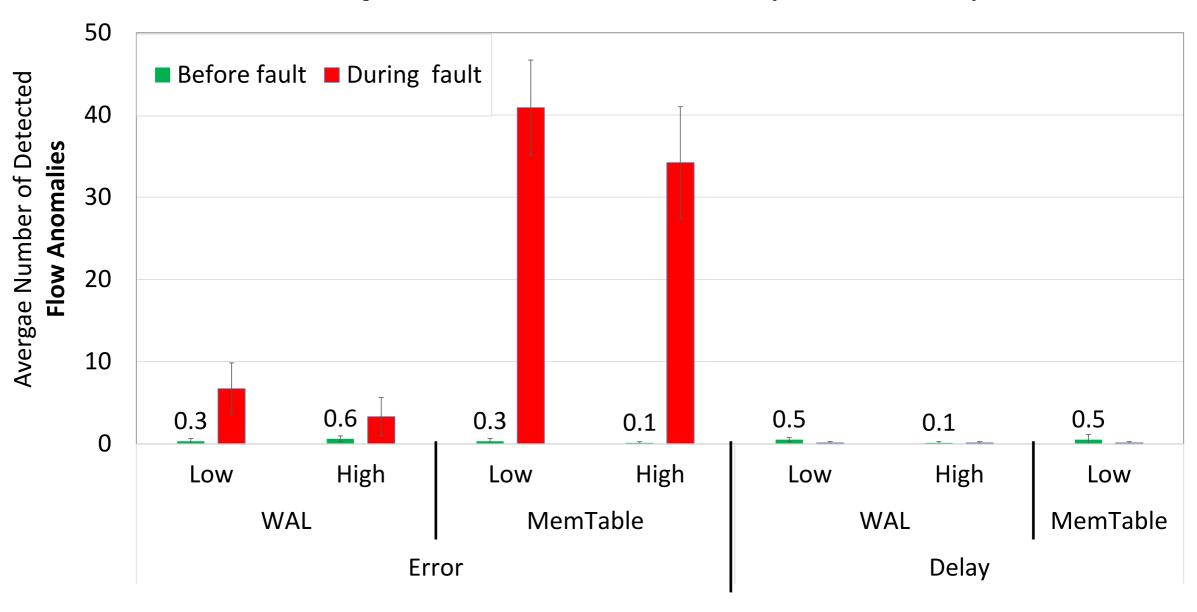
#### **Runtime Overhead**



No overhead when using SAAD vs. default (INFO) logging

#### False Positive: Flow anomalies

#### Each experiment: 90min (10 times)



1 false alarm per hour

### Stage Aware Anomaly Detection

Provides context to the anomaly

- ▶ Code module (stage)
- Execution flow (log points)

In Real Time, with Low overhead

- Minimal runtime overhead
- Low storage cost (synopses instead of logs)
- Low processing cost

Portable, evaluated on 3 distributed storage systems

#### Questions

Details in Saeed Ghanbari's PhD Thesis (2014)

and in ACM/IFIP/Usenix Middleware 2014

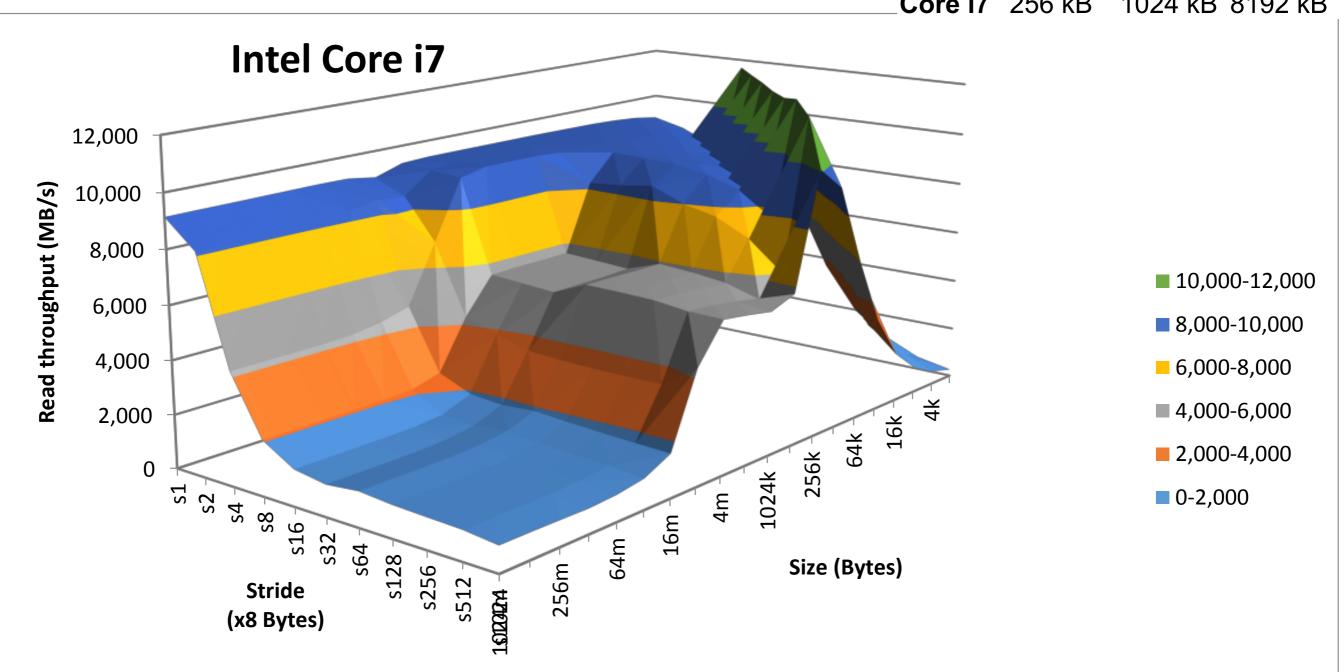
In collaboration with

Saeed Ghanbari and Ali Hashemi

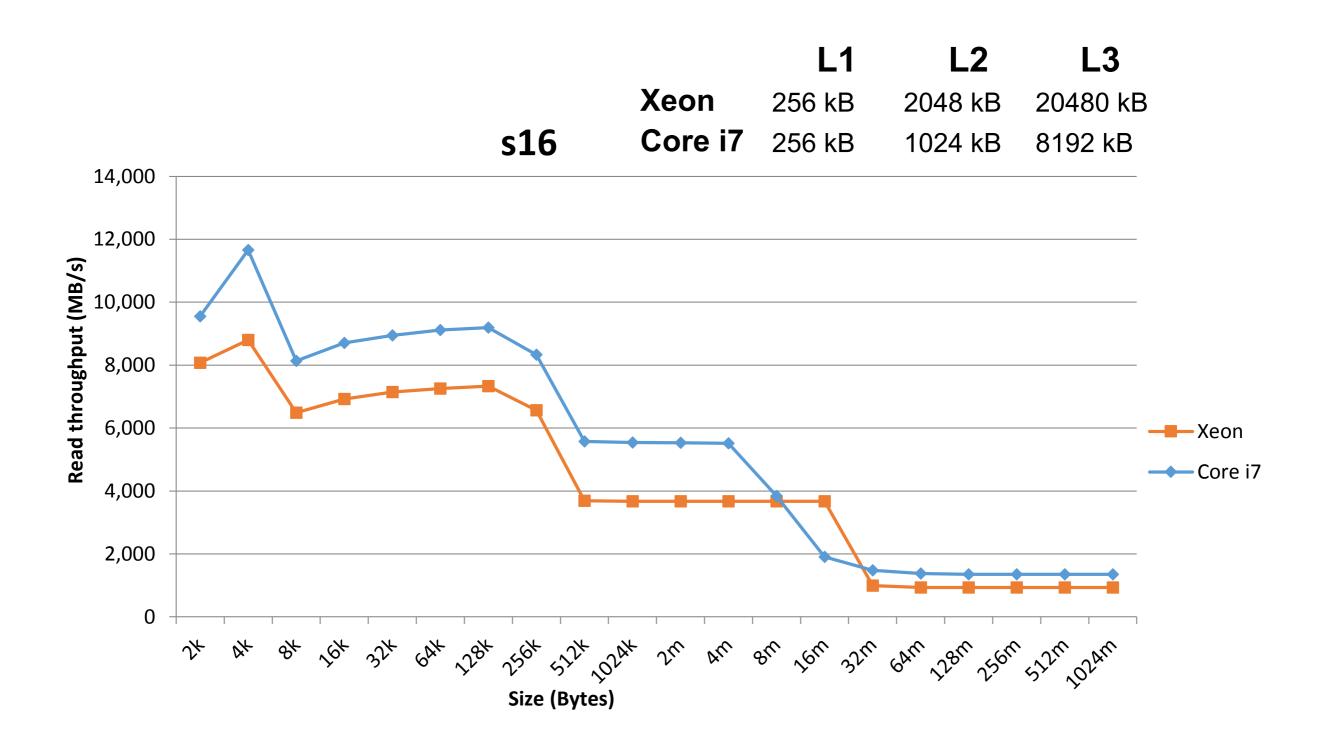
#### Ongoing Work: Model Transformation

Cache L1 L3 **L2** 

**Core i7** 256 kB 1024 kB 8192 kB



#### Guidance: e.g., 3D to 2D reduction



### With minimum new samples ...

