# Tracing Interrupts in Embedded Software

#### Giovani Gracioli

Federal University of Santa Catarina giovani@lisha.ufsc.br

#### Sebastian Fischmeister

University of Waterloo sfischme@uwaterloo.ca



• Embedded => Interrupts (Especially in low power systems)

- Embedded => Interrupts (Especially in low power systems)
- Interrupts => Complexity (Non-linear program execution)

- Embedded => Interrupts (Especially in low power systems)
- Interrupts => Complexity
   (Non-linear program execution)
- Complexity => Nasty bugs
   (You only need one try, right?)

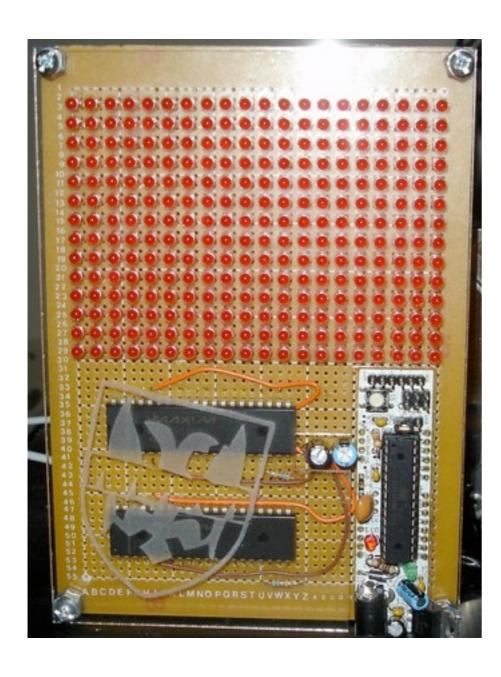
```
$ echo 'int main() { printf ("Hello,
world\n"); }' | gcc -xc - -o out
```

- Embedded => Interrupts (Especially in low power systems)
- Interrupts => Complexity
   (Non-linear program execution)
- Complexity => Nasty bugs
   (You only need one try, right?)
- Nasty bugs require good debugging support

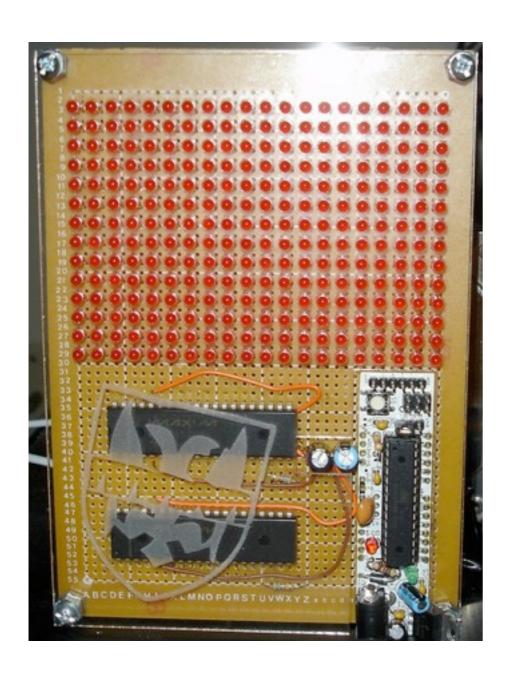


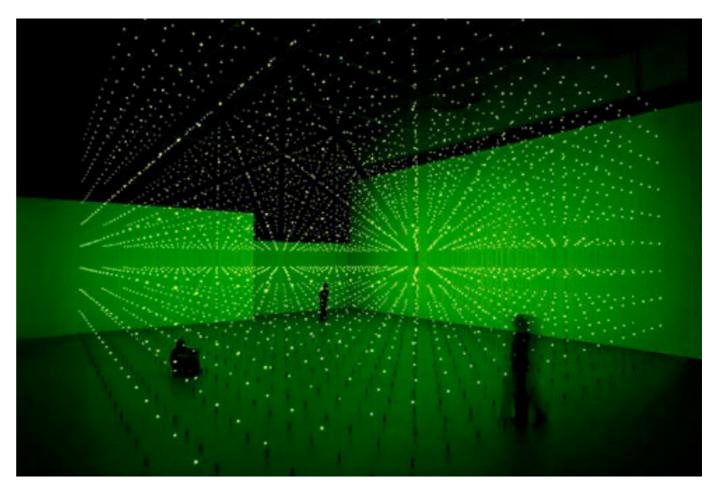
#### Next Gen Debugging Support?

#### Next Gen Debugging Support?



#### Next Gen Debugging Support?

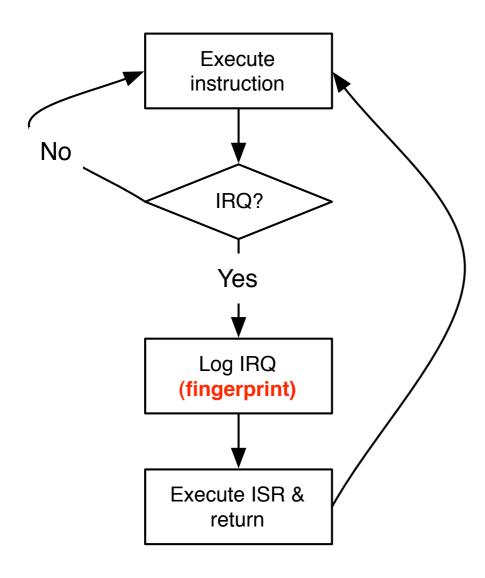




Recording at run time

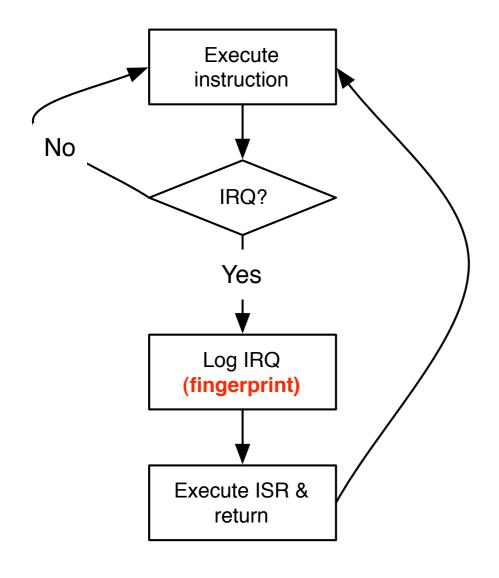
Replay offline (sim)

#### Recording at run time

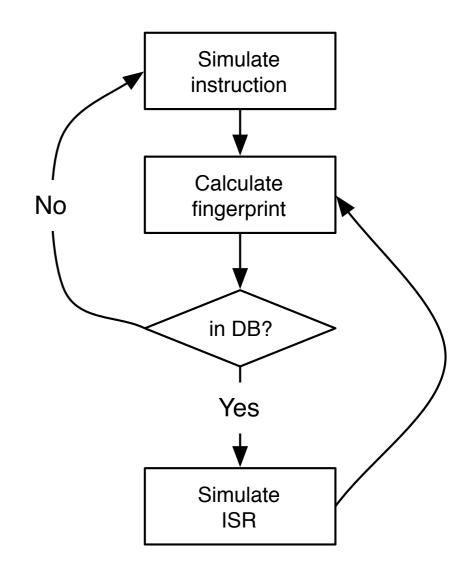


Replay offline (sim)

#### Recording at run time



#### Replay offline (sim)

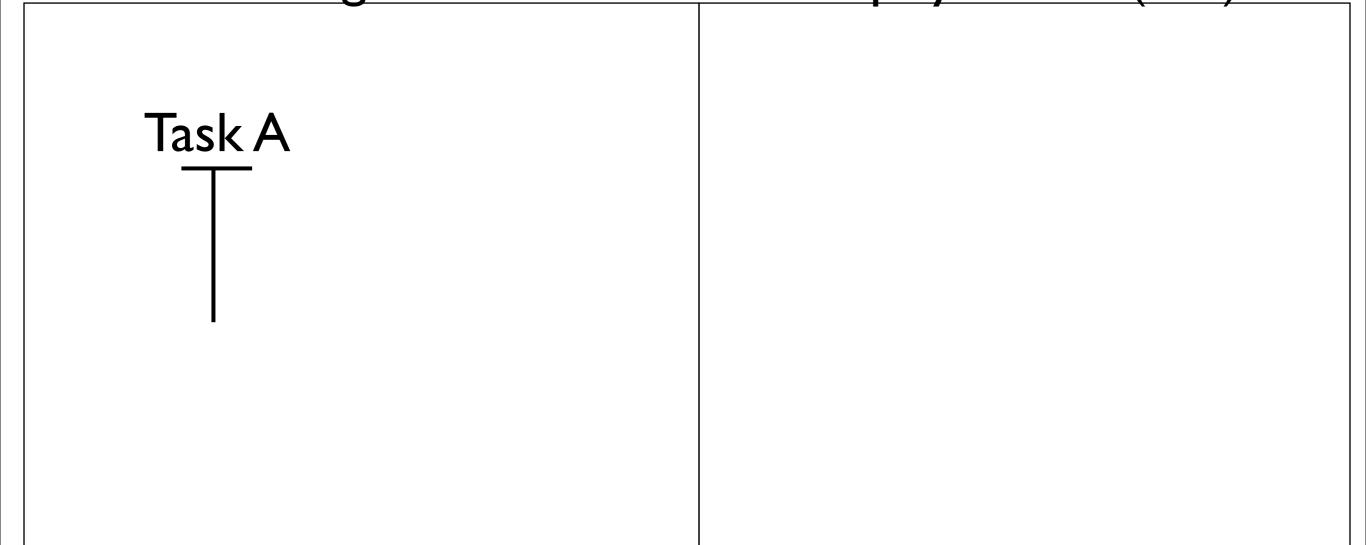


_	Recording at run time	Replay offline (sim)
1		



Recording at run time

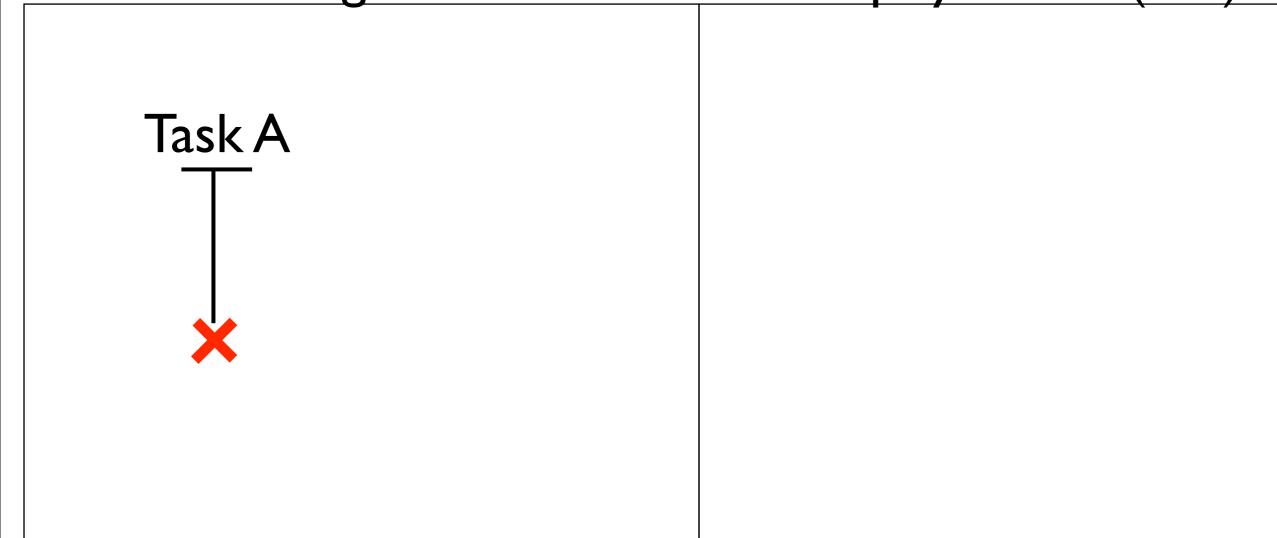
Replay offline (sim)





Recording at run time

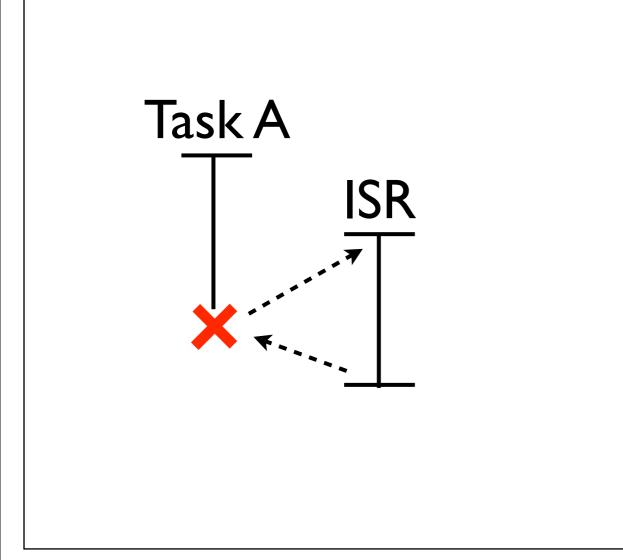
Replay offline (sim)



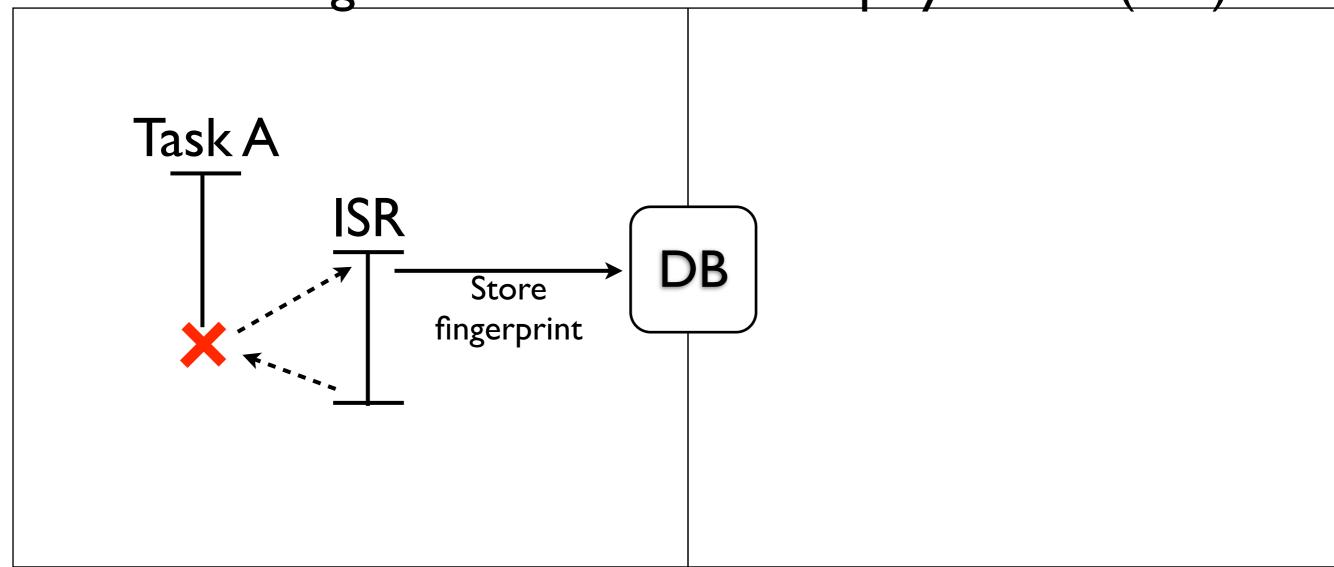


Recording at run time

Replay offline (sim)

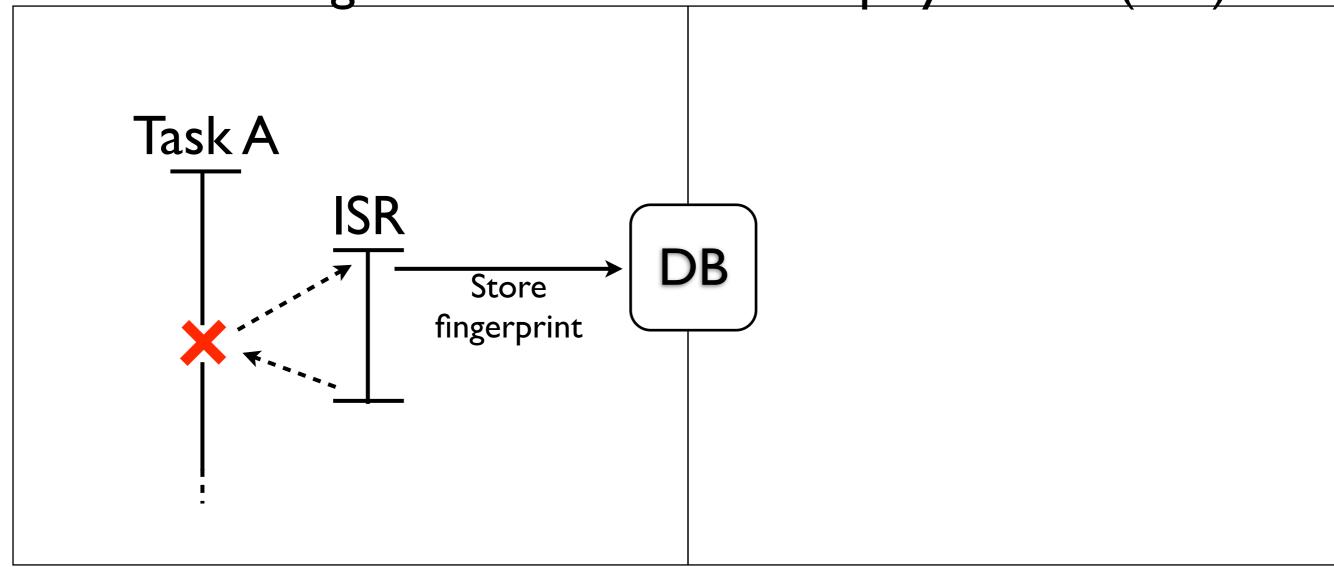


Recording at run time Replay offline (sim)





Recording at run time Replay offline (sim)





Replay offline (sim) Recording at run time Task A Task A **ISR ISR** Store fingerprint



Replay offline (sim) Recording at run time Task A Task A **ISR ISR** DB Store fingerprint

XN

Replay offline (sim) Recording at run time Task A Task A **ISR ISR** DB Store fingerprint



Replay offline (sim) Recording at run time Task A Task A **ISR ISR** DB Store fingerprint



Replay offline (sim) Recording at run time Task A Task A **ISR ISR** DB Store fingerprint



Replay offline (sim) Recording at run time Task A Task A **ISR ISR** DB Store fingerprint



Replay offline (sim) Recording at run time Task A Task A **ISR** ISR DB Store fingerprint



Replay offline (sim) Recording at run time Task A Task A **ISR** ISR DB Store fingerprint



Replay offline (sim) Recording at run time Task A Task A **ISR** ISR DB Store fingerprint



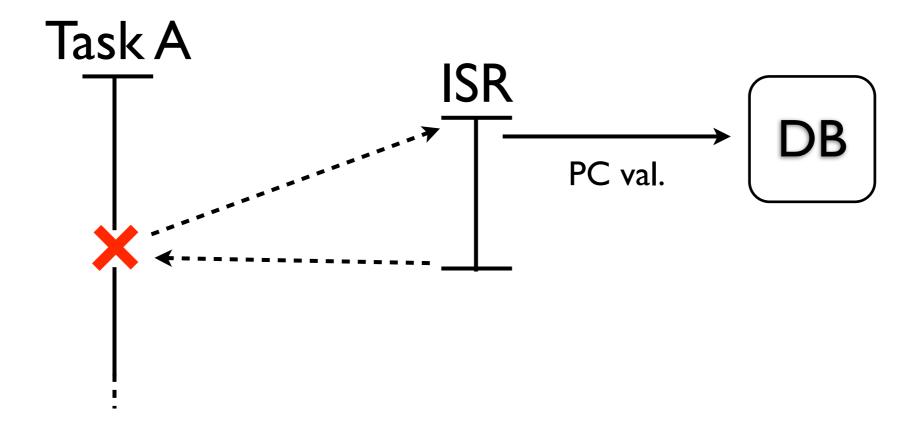
Replay offline (sim) Recording at run time Task A Task A **ISR** ISR DB Store fingerprint



#### Benefits

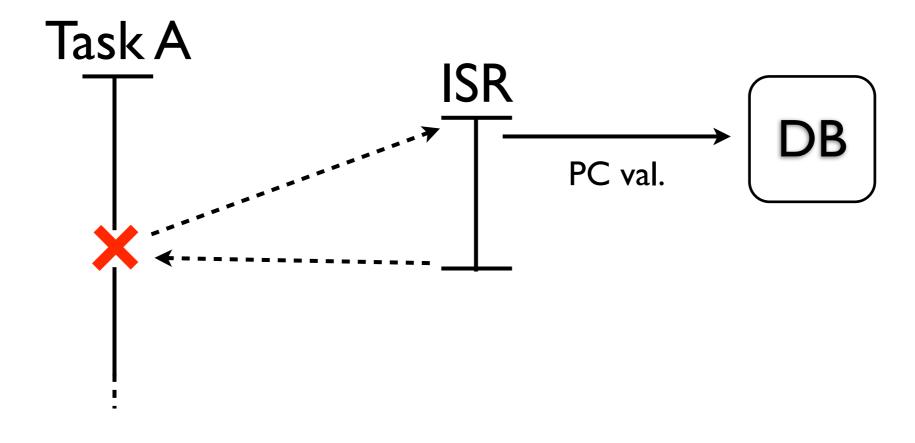
- Communicate and document bugs (Show the trace)
- Go forward and backwards in time (Thoroughly analyze how the bug happened)
- Retest specific scenario for fixes (Exercise the system with the specific trace)

#### First try: record program counter



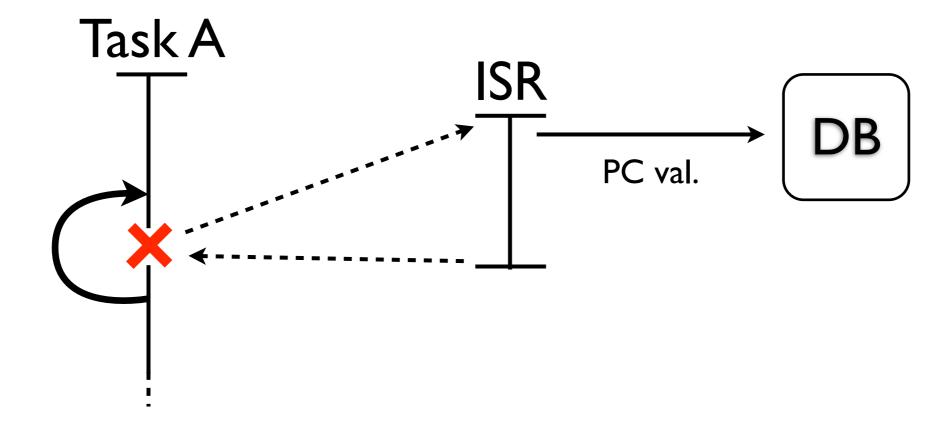


#### First try: record program counter



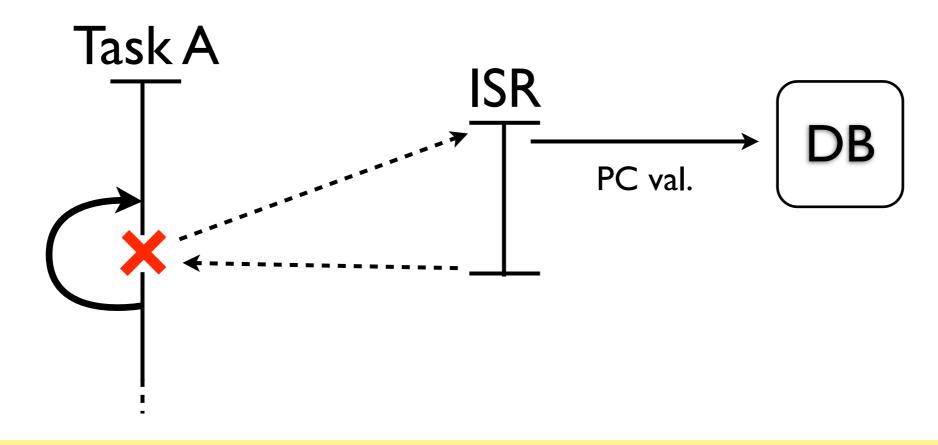


#### First try: record program counter





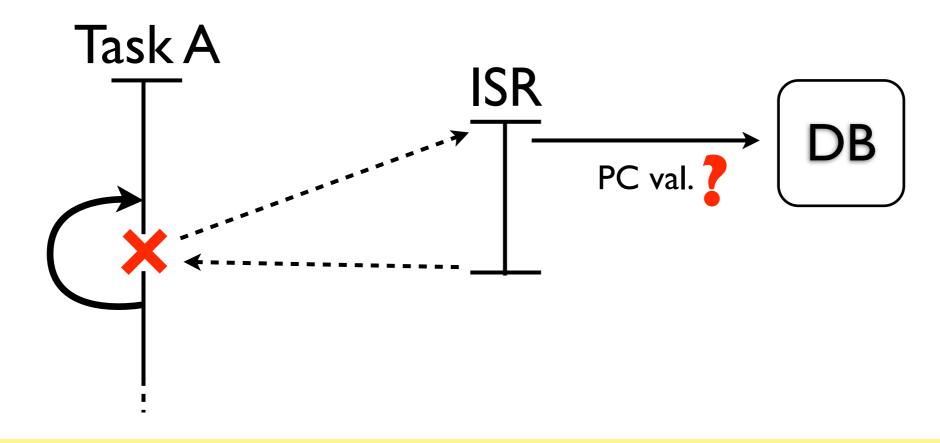
#### First try: record program counter



Second try: record program counter and state information



#### First try: record program counter



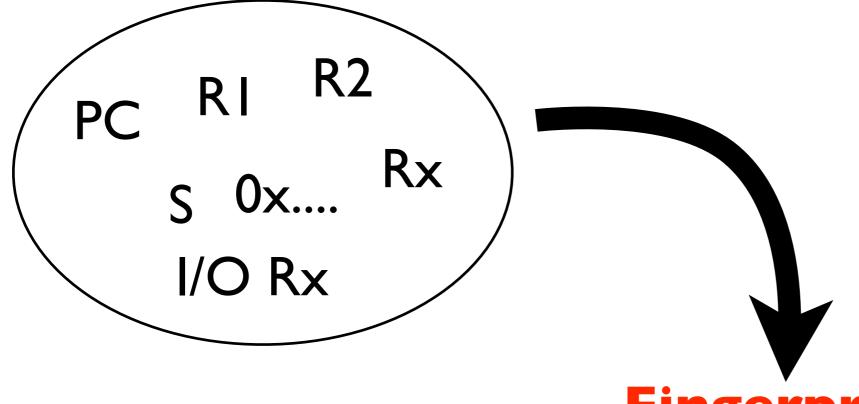
Second try: record program counter and state information



# Fingerprinting

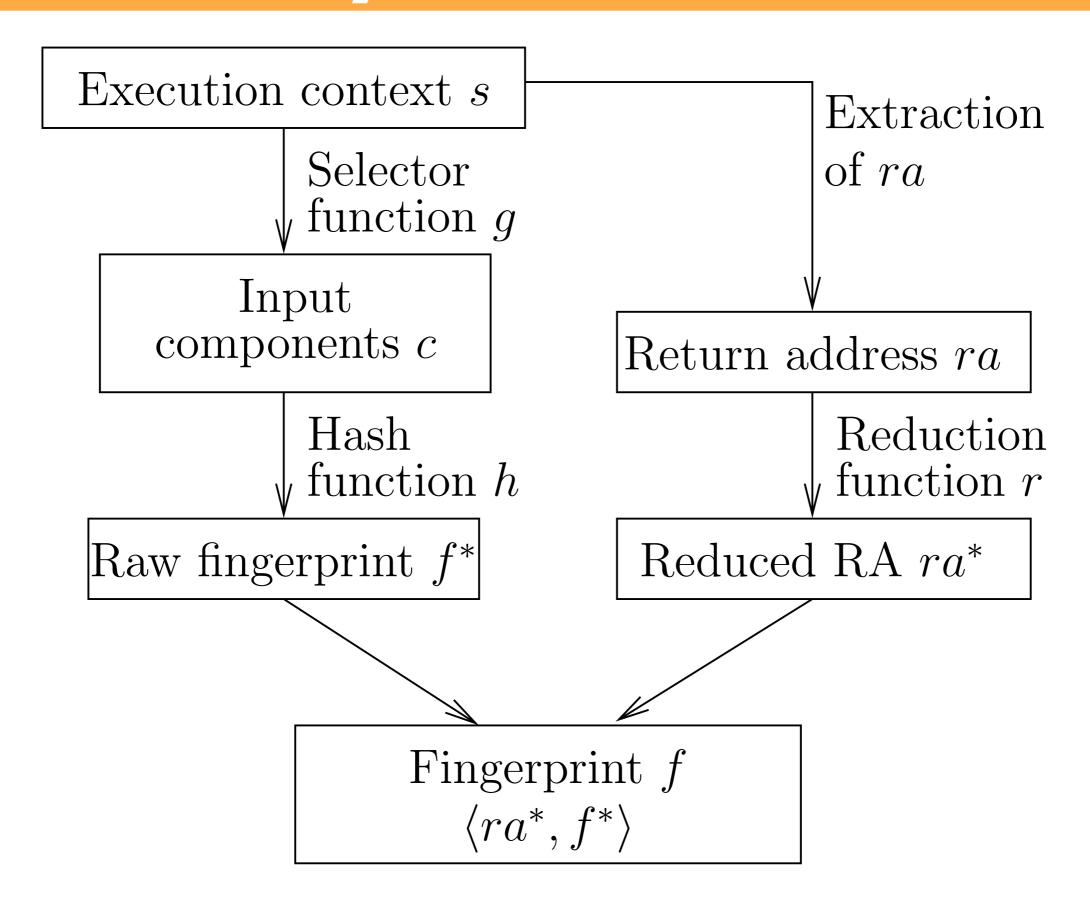
#### Compress system state

System state 5|2+ bits



Fingerprint <32 bits 0100100100100100100100100

# System Model



#### Caveats

False input duplicates (selected bad subset)

$$s \neq s' \land g(s) = g(s')$$

• False loop positives (falsely believe IRQ happened here)

$$s \neq s' \land s.ra = s'.ra$$
 but also  $(h \cdot g)(s) = (h \cdot g)(s')$ 

• False RA positives (falsely believe IRQ happened here)

$$s \neq s' \land s.ra \neq s'.ra$$
 but also  $r(s.ra) = r(s'.ra) \land (h \cdot g)(s) = (h \cdot g)(s')$ 

#### Caveats

False input duplicates (selected bad subset)

$$s \neq s' \land g(s) = g(s')$$

• False loop positives (falsely believe IRQ happened here)

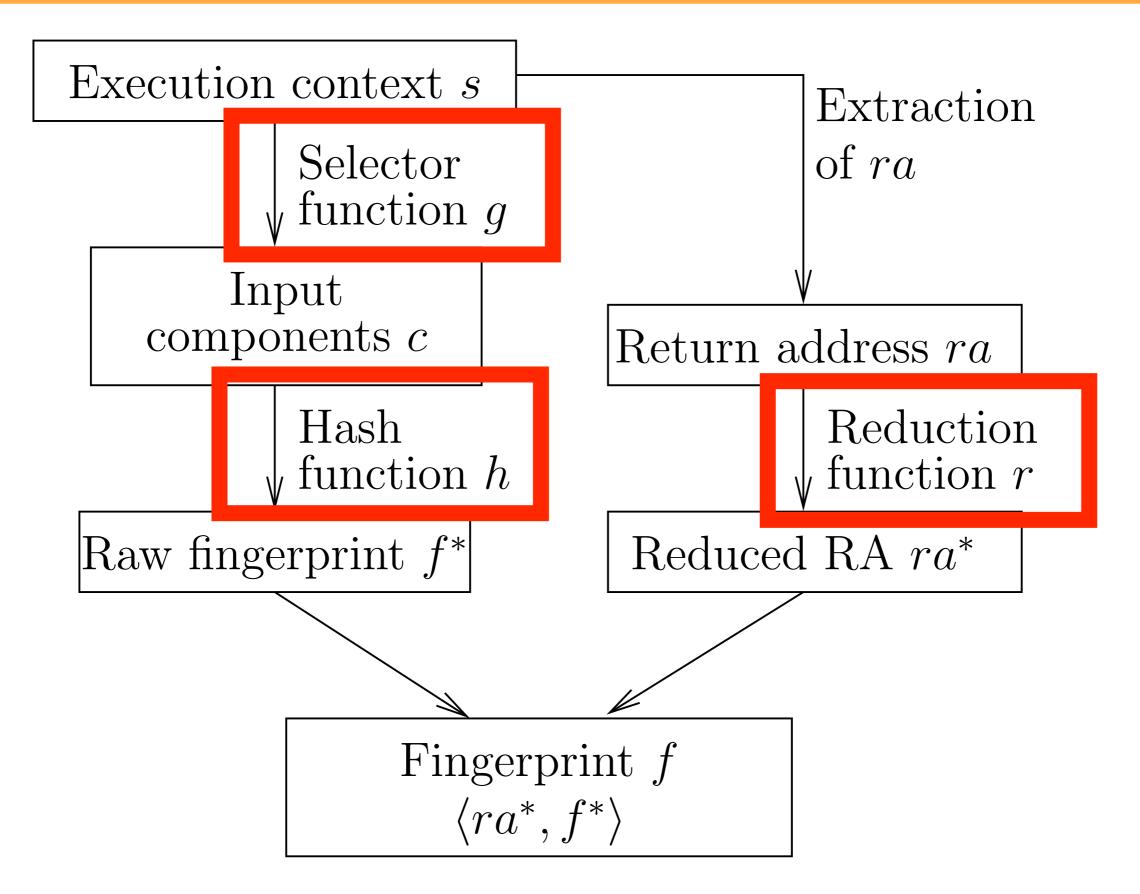
$$s \neq s' \land s.ra = s'.ra$$
 but also  $(h \cdot g)(s) = (h \cdot g)(s')$ 

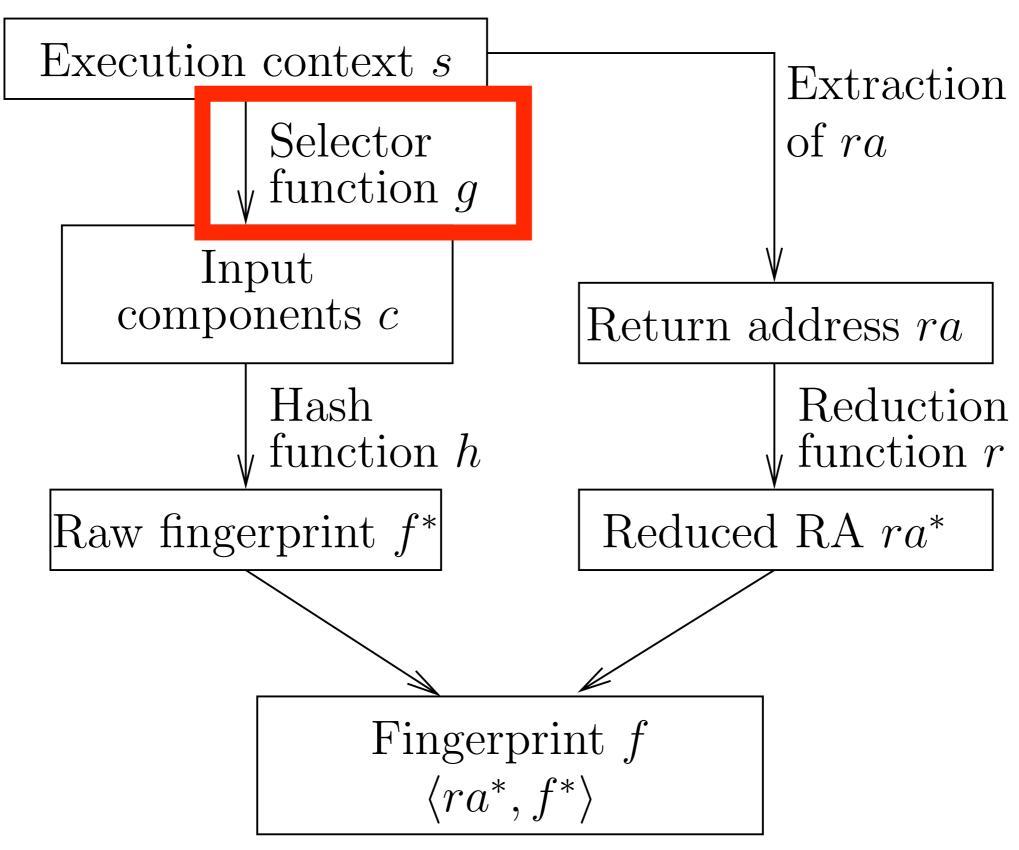
• False RA positives (falsely believe IRQ happened here)

$$s \neq s' \land s.ra \neq s'.ra$$
 but also  $r(s.ra) = r(s'.ra) \land (h \cdot g)(s) = (h \cdot g)(s')$ 

### Calls for good design of fingerprint

# Designing the Fingerprint





### Selector Function

#### Selector Function

• Many elements to choose from: data registers, control registers, GPIO, IO status, general SDRAM, stack, heap, ...

#### Selector Function

 Many elements to choose from: data registers, control registers, GPIO, IO status, general SDRAM, stack, heap, ...

#### Can't take all:

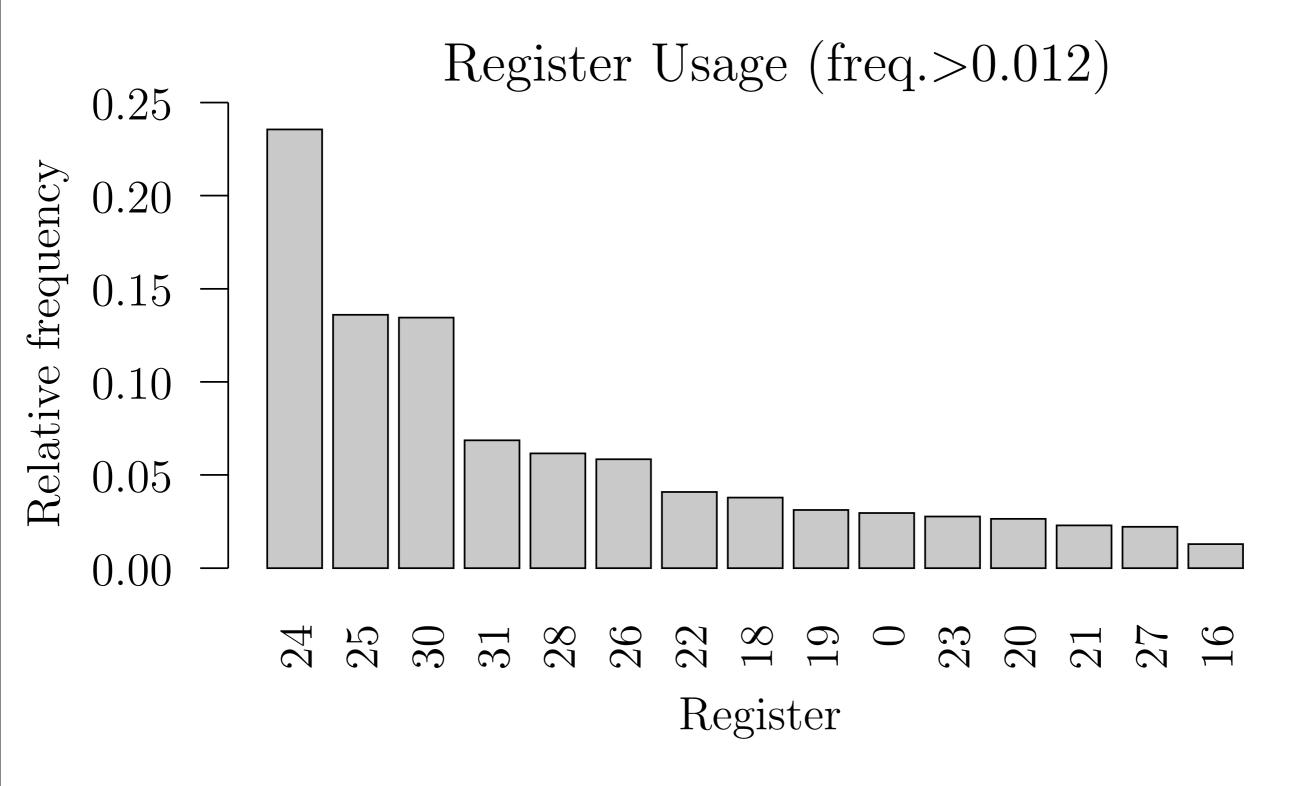
- Too much computation overhead
- Not all are used equally often

#### ldea 1

### Pick the most frequently used ones

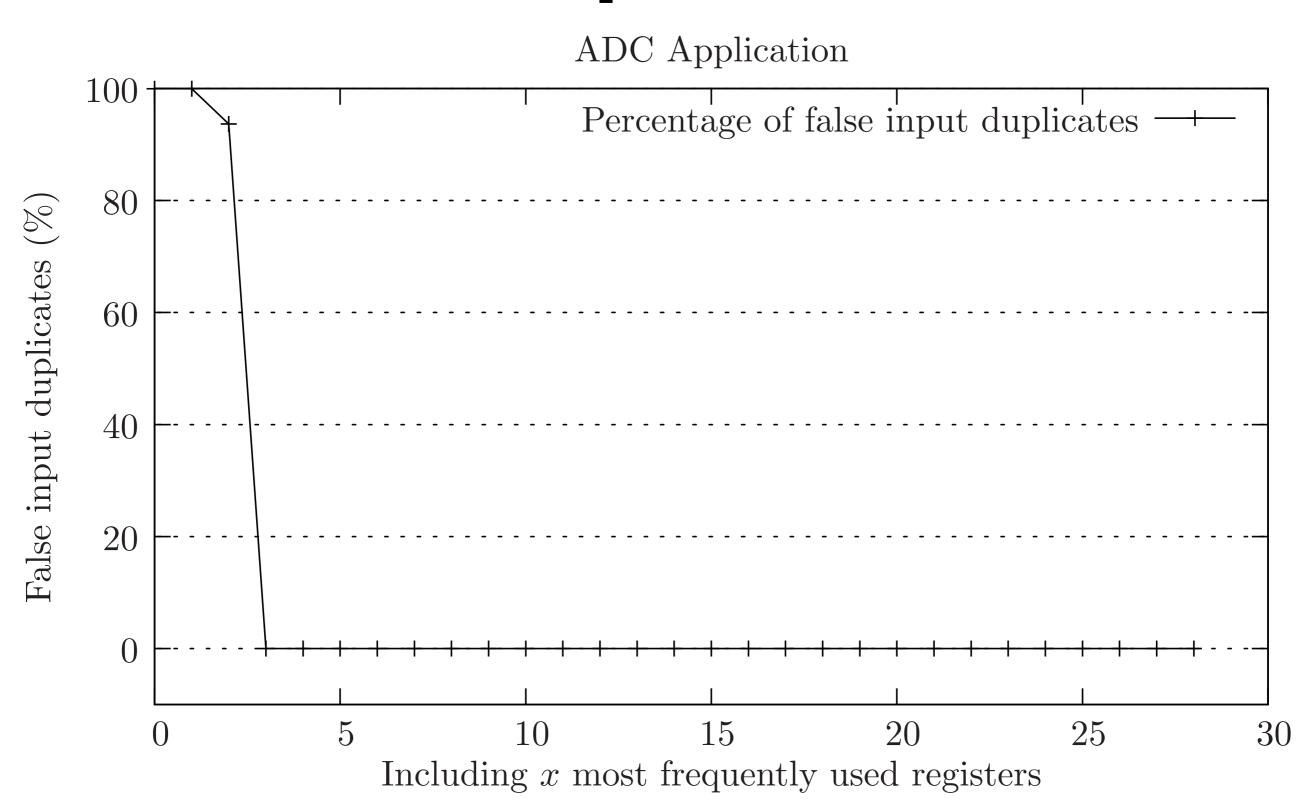
#### ldea 1

### Pick the most frequently used ones

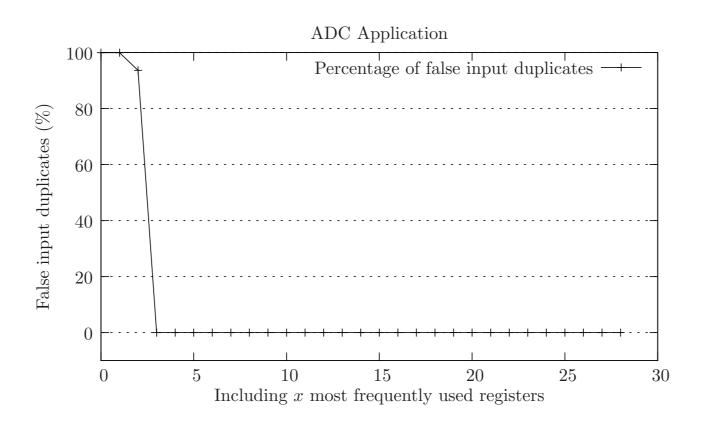


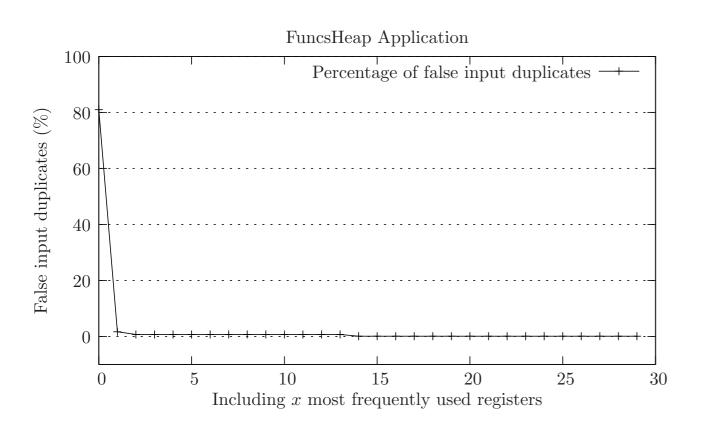
#### Does it work?

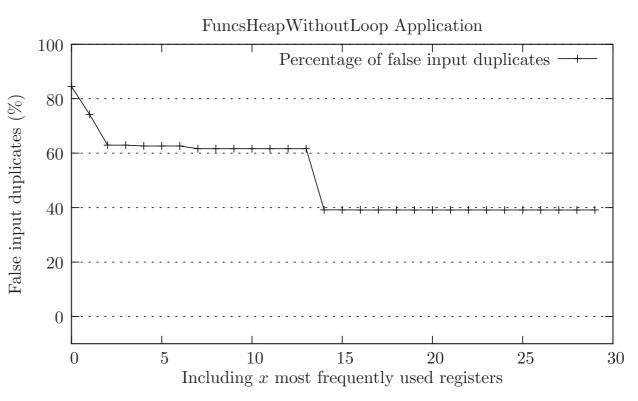
### Yes, quite well.

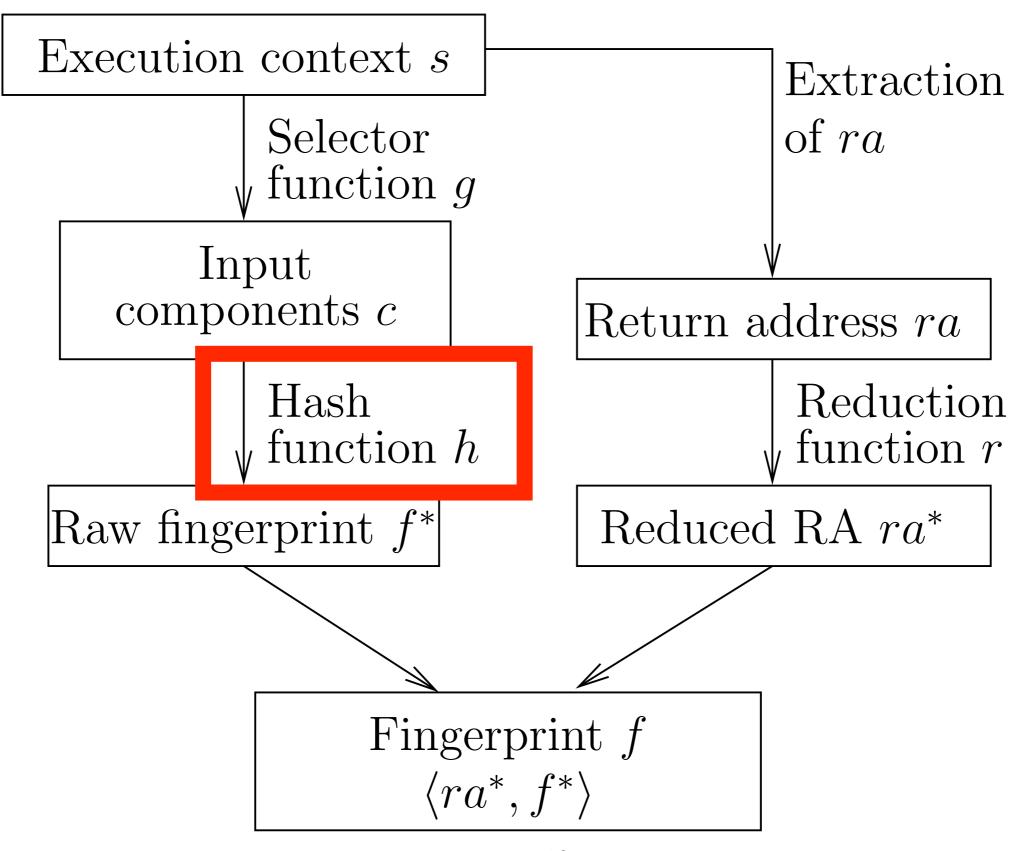


## Does it work?

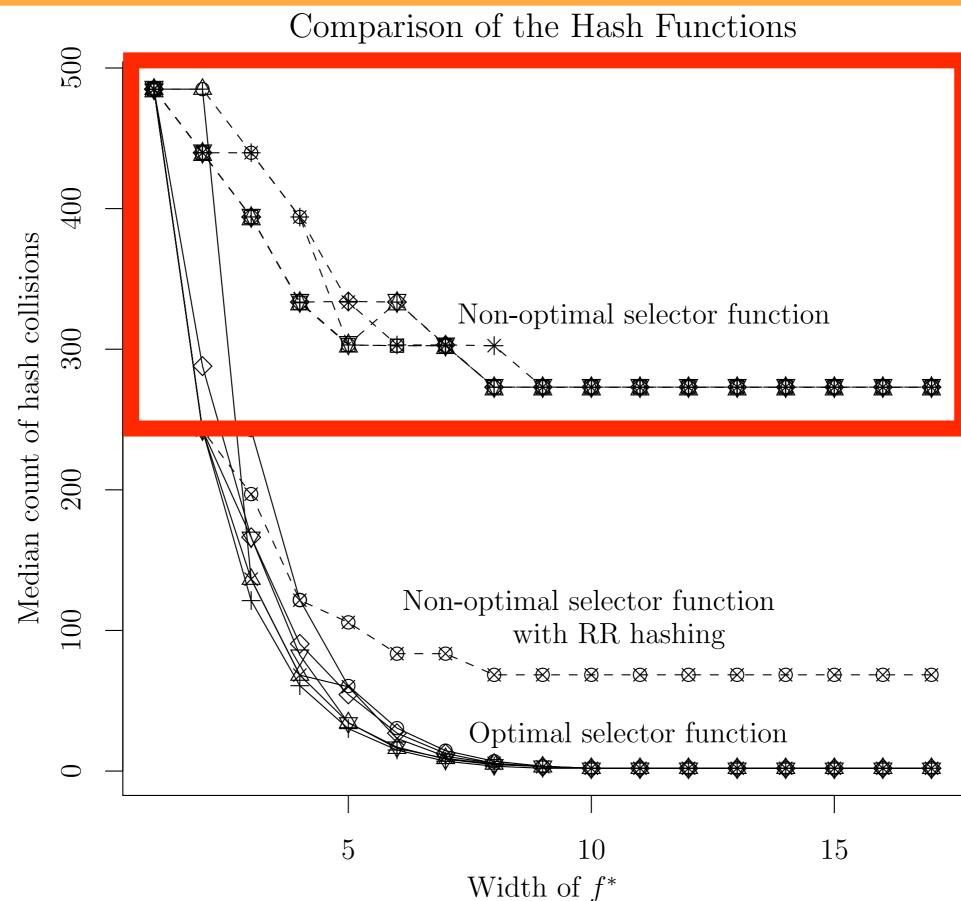








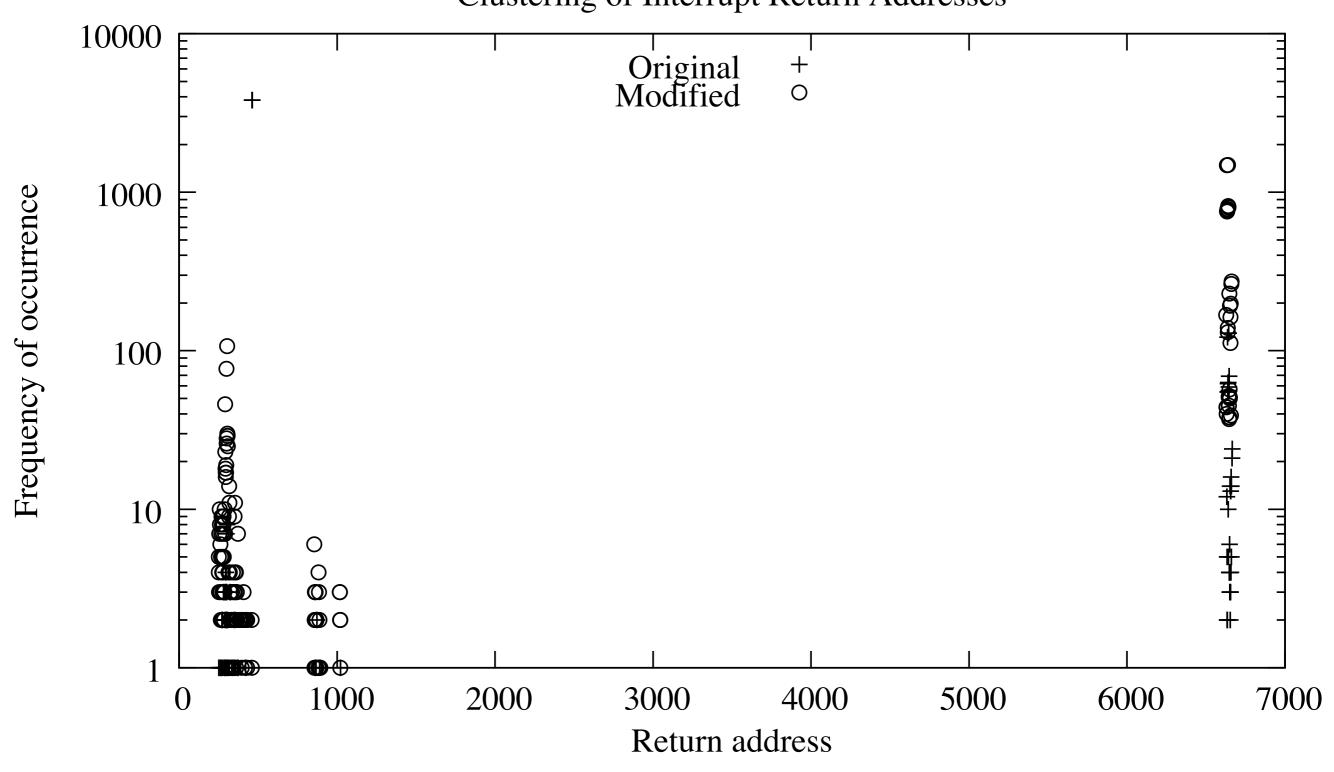
# All Perform Equally Badly



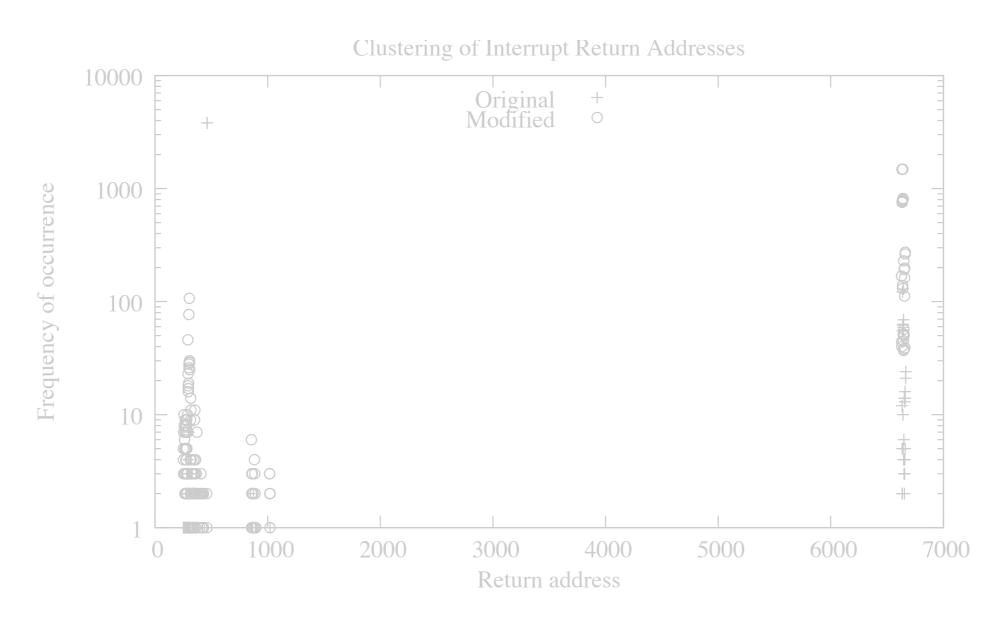
### Observation

#### Return addresses tend to cluster

Clustering of Interrupt Return Addresses



#### Return addresses tend to cluster



#### Return addresses tend to cluster

Clustering of Interrupt Return Addresses



#### Return addresses tend to cluster

Clustering of Interrupt Return Addresses

- Tight loops: busy waiting, iterative calculations, array/matrix operations
- Blocking operations: non-preemptive peripheral access, waiting for user input, costly instructions

Return address

# Why?

#### Return addresses tend to cluster

Clustering of Interrupt Return Addresses

- **Tight loops:** busy waiting, iterative calculations, array/matrix operations
- Blocking operations: non-preemptive peripheral access, waiting for user input, costly instructions
- Disabled interrupts: generic concurrency control, I/O access

# Why?

#### Return addresses tend to cluster

Clustering of Interrupt Return Addresses

- **Tight loops:** busy waiting, iterative calculations, array/matrix operations
- Blocking operations: non-preemptive peripheral access, waiting for user input, costly instructions
- Disabled interrupts: generic concurrency control, I/O access

Context specific hashing => RR hashing

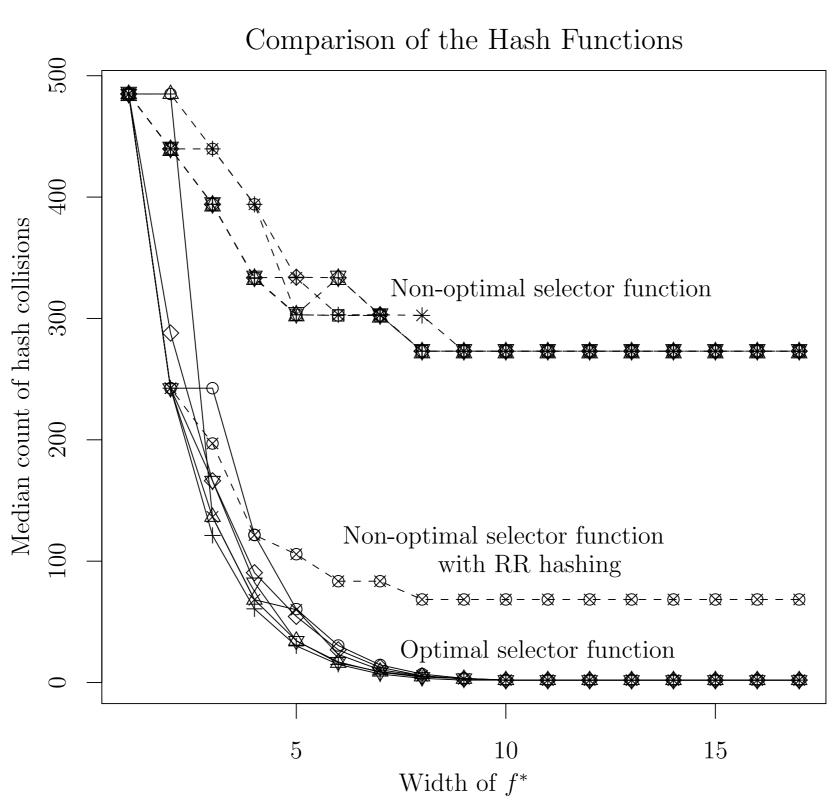
Use different hash functions to minimize collisions

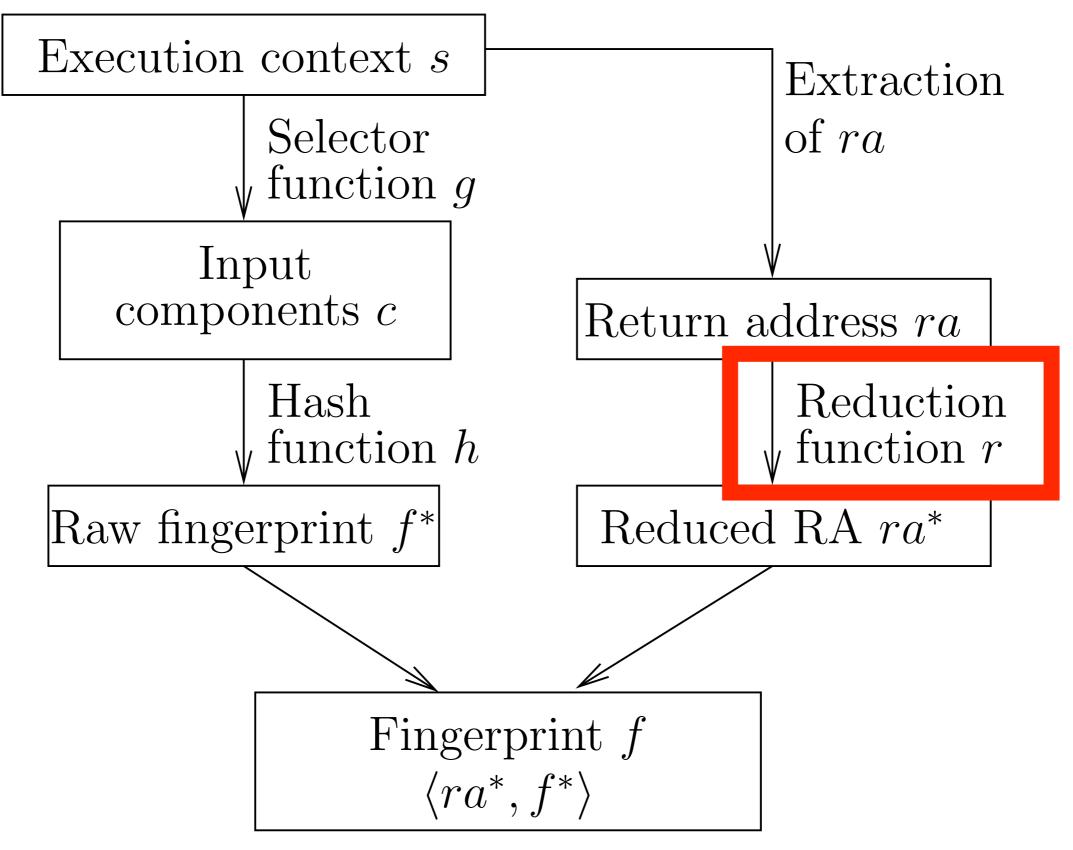
- Use different hash functions to minimize collisions
- Insert markers in the code for different hash functions

- Use different hash functions to minimize collisions
- Insert markers in the code for different hash functions
- In the replay use the correct hash function for according to the marker

### Does it work?

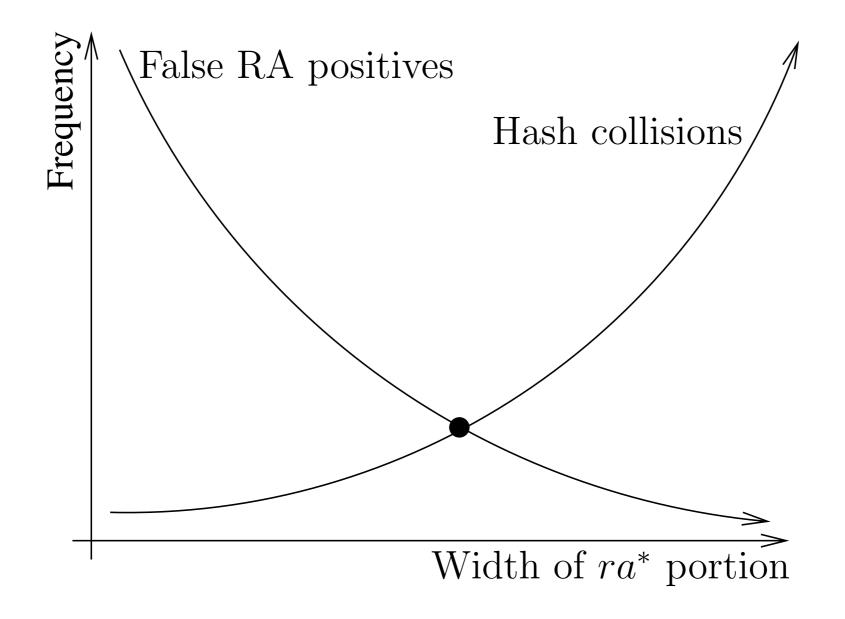
### RR hashing works surprisingly well





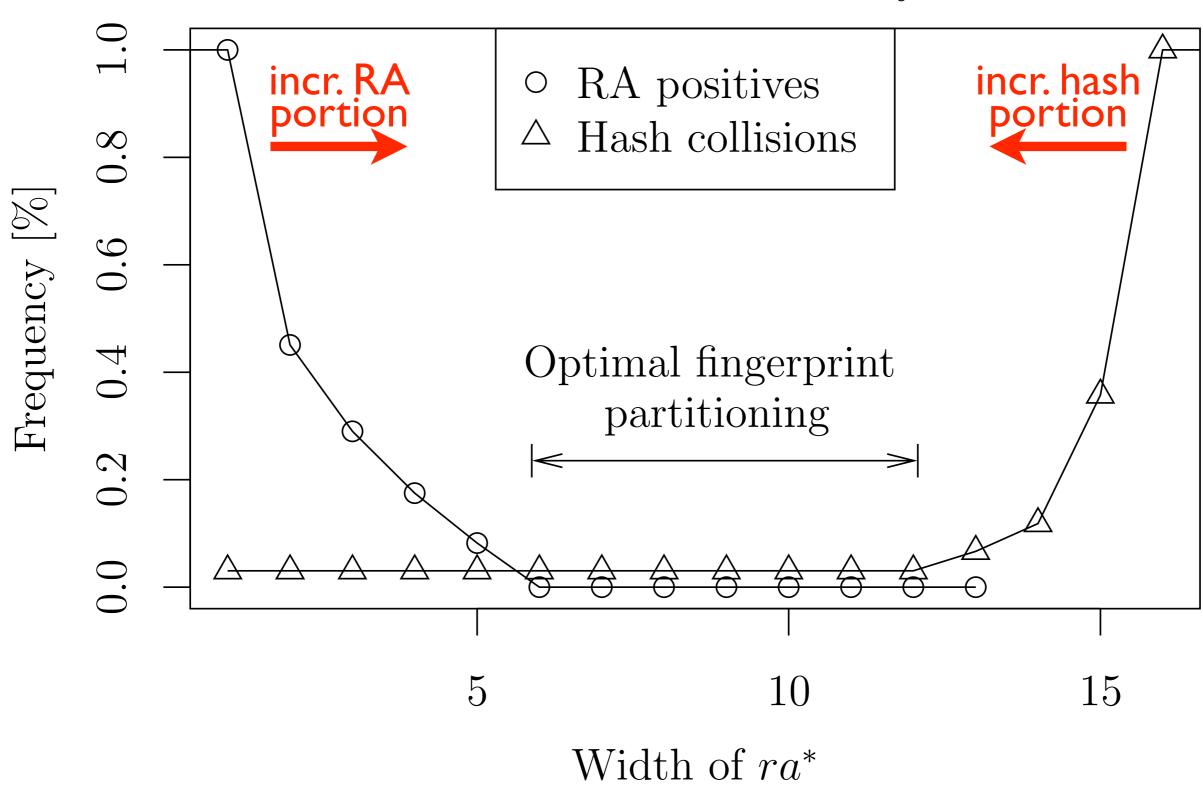
### How best to use the FP bits?

- Give all to the hash value? (minimize hash collisions)
- Give all to the RA? (minimize false RA positives)



# Engineering the Fingerprint

Trade-off Between  $ra^*$  and  $f^*$  Width



#### Conclusions

 Debugging embedded systems is painful => Replay debugging can help [ST08]

Frequency-based selection function



Round robin hashing



Tradeoff engineering for RA vs hash



Future work: Lots, see discussion section

# Questions?

Work has been sponsored by:

Graduate Students' Exchange Program (GSEP) funded by Foreign Affairs and International Trade Canada (DFAIT),
Ontario Research Fund Research Excellence (ORF-RE) &

CIMIT under U.S. Army Medical Research Acquisition Activity Cooperative Agreement W81XWH-07-2-0011