

Qudit versions of the $\pi/8$ gate: Applications in fault-tolerant QC and nonlocality

Mark Howard¹ & Jiri Vala^{1,2}

October 2012

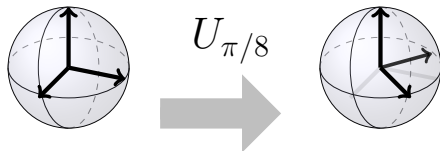
¹National University of Ireland, Maynooth, Ireland.

²Dublin Institute for Advanced Studies, Dublin, Ireland.



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The $U_{\pi/8}$ gate and its uses



$$U_{\pi/8} = \begin{pmatrix} e^{-i\frac{\pi}{8}} & 0 \\ 0 & e^{i\frac{\pi}{8}} \end{pmatrix} \propto \begin{pmatrix} 1 & 0 \\ 0 & e^{i\frac{\pi}{4}} \end{pmatrix}$$

- $\text{UQC} = \langle \text{Cliffords}, U_{\pi/8} \rangle$
 $\text{UQC} \neq \langle \text{Cliffords} \rangle$

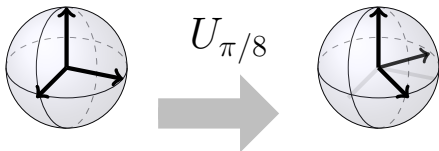


P. O. Boykin, T. Mor, M. Pulver, V. Roychowdhury and F. Vatan.

A new universal and fault-tolerant quantum basis

Information Processing Letters 75, 3 pp. 101–107, (2000).

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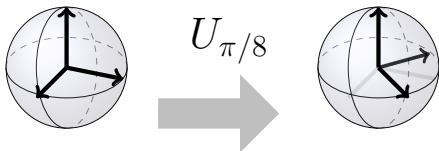
- $\text{UQC} = \langle \text{Cliffords}, U_{\pi/8} \rangle$
 $\text{UQC} \neq \langle \text{Cliffords} \rangle$
- Teleportation-based UQC



D. Gottesman and I. L. Chuang,

Demonstrating the viability of universal quantum computation using teleportation and single-qubit operations
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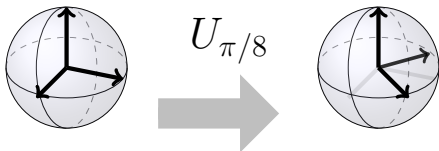
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- Teleportation-based UQC
- Transversal for R-M codes



B. Zeng, H. Chung, A. Cross and I. Chuang,

Local unitary versus local Clifford equivalence of stabilizer and graph states,
Phys. Rev. A 75, 032325 (2007).

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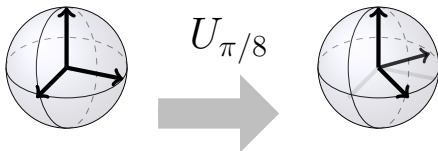
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- Teleportation-based UQC
- Transversal for R-M codes
- Topological Protection (3D)



S. Bravyi and R. Koenig

Classification of topologically protected gates for local stabilizer codes,
arXiv:1206.1609 (2012).

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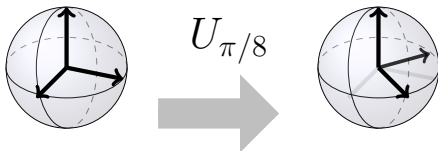
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 - Topological Protection (3D)
 - Secure assisted UQC
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A. M. Childs,

Secure assisted quantum computation
Quantum Info. Comput.5, pp. 456, (2005).

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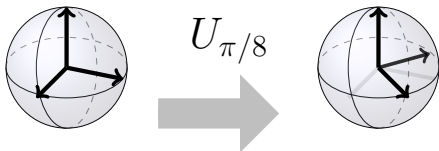
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- Measurement-based UQC with graph states



M. Silva, V. Danos, E. Kashefi and H. Ollivier,

A direct approach to fault-tolerance in measurement-based quantum computation via teleportation
New Journal of Physics 9, 6 pp. 192, (2007).

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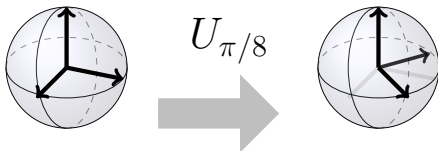
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- Measurement-based UQC with graph states
 - Optimal CHSH game with $(|00\rangle + |11\rangle)/\sqrt{2}$



M. Howard and J. Vala,

Nonlocality as a benchmark for universal quantum computation in Ising anyon topological quantum computers, Phys. Rev. A **85**, 022304 (2012).

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- Measurement-based UQC with graph states
 - Optimal CHSH game with $(|00\rangle + |11\rangle)/\sqrt{2}$
 - Blind UQC

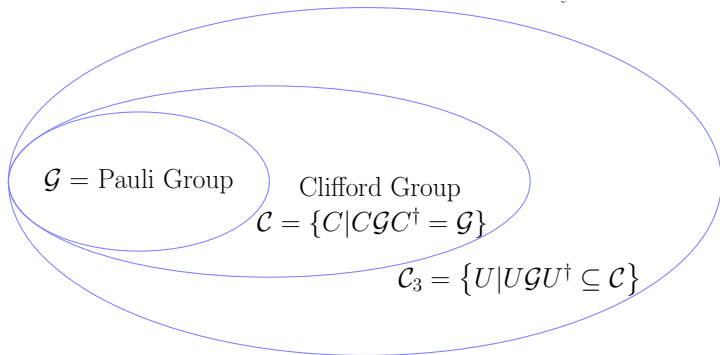


A. Broadbent, J. Fitzsimons and E. Kashefi,

Universal blind quantum computation,

Annual IEEE Symposium on Foundations of Computer Science, pp. 517–526, (2009).

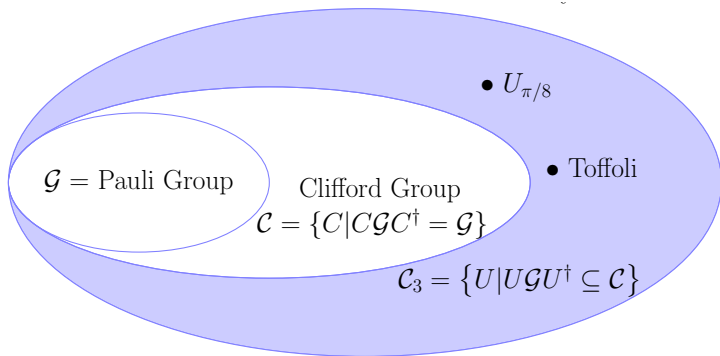
Structure of Pauli/Clifford groups



D. Gottesman and I. L. Chuang,

Demonstrating the viability of universal quantum computation using teleportation and single-qubit operations
Nature 402, 6760 pp. 390–393, (1999).

Structure of Pauli/Clifford groups



We will focus on **single, p-level** particles

- Generalized σ_x/σ_z : $X|j\rangle = |j+1 \bmod p\rangle$ $Z|j\rangle = \omega^j|j\rangle$ ($\omega = e^{2\pi i/p}$)
- Displacement operators, $D_{(x|z)} = \omega^{2^{-1}xz} X^x Z^z$, form Pauli group \mathcal{G}



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Single-Qudit Clifford Gates

$$\mathcal{C} = \{C_{(F|\vec{\chi})} \mid F \in SL(2, \mathbb{Z}_p), \vec{\chi} \in \mathbb{Z}_p^2\},$$

- $F = \begin{pmatrix} \alpha & \beta \\ \gamma & \delta \end{pmatrix}$ with unit determinant. $\vec{\chi} = \begin{pmatrix} x \\ z \end{pmatrix}$ is a vector of length 2.
- All elements of $F, \vec{\chi}$ are from $\mathbb{Z}_p = \{0, 1, \dots, p-1\}$
- Explicit recipe for constructing a Clifford unitary (where $\tau = \omega^{2^{-1}}$):

$$C_{(F|\vec{\chi})} = D_{(x|z)} V_F$$

$$V_F = \begin{cases} \frac{1}{\sqrt{p}} \sum_{j,k=0}^{p-1} \tau^{\beta^{-1}(\alpha k^2 - 2jk + \delta j^2)} |j\rangle \langle k| & \beta \neq 0 \\ \sum_{k=0}^{p-1} \tau^{\alpha \gamma k^2} |\alpha k\rangle \langle k| & \beta = 0. \end{cases}$$



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Properties of the extended Clifford group with applications to SIC-POVMs and MUBs,
arXiv:quant-ph/0909.5233, (2009).

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$$C\left(\begin{bmatrix} 1 & 0 \\ \gamma & 1 \end{bmatrix} \middle| \begin{bmatrix} x \\ z \end{bmatrix}\right) \in \text{SU}(p) \quad \forall p > 3,$$

$$\det\left(C\left(\begin{bmatrix} 1 & 0 \\ \gamma & 1 \end{bmatrix} \middle| \begin{bmatrix} x \\ z \end{bmatrix}\right)\right) = \tau^{2\gamma} \text{ for } p = 3,$$

U_v as qudit version of $U_{\pi/8}$ $(p > 3)$

Define
$$U_v = U(v_0, v_1, \dots) = \sum_{k=0}^{p-1} \omega^{v_k} |k\rangle\langle k| \quad (v_k \in \mathbb{Z}_p)$$

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Easy to show
$$U_v D_{(x|z)} U_v^\dagger = D_{(x|z)} \sum_k \omega^{(v_{k+x} - v_k)} |k\rangle\langle k|$$

and in particular
$$U_v D_{(1|0)} U_v^\dagger = D_{(1|0)} \sum_k \omega^{(v_{k+1} - v_k)} |k\rangle\langle k| \quad *$$

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Must have
$$\begin{aligned} U_v D_{(1|0)} U_v^\dagger &= \omega^{\epsilon'} C \left(\begin{bmatrix} 1 & 0 \\ \gamma' & 1 \end{bmatrix} \middle| \begin{bmatrix} 1 \\ z' \end{bmatrix} \right) \\ &= \omega^{\epsilon'} D_{(1|z')} \sum_{k=0}^{p-1} \tau^{\gamma' k^2} |k\rangle\langle k| \quad ** \end{aligned}$$

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$$= \omega^{\epsilon'} D_{(1|z')} \sum_{k=0}^{p-1} \tau^{\gamma' k^2} |k\rangle\langle k| \quad **$$

Equating * and **:
$$\omega^{v_{k+1} - v_k} = \omega^{\epsilon'} \tau^{z'} \omega^{kz'} \tau^{k^2 \gamma'} \quad (\forall k \in \mathbb{Z}_p),$$

$$\Rightarrow v_k = 12^{-1} k(\gamma' + k(6z' + (2k - 3)\gamma')) + k\epsilon' \quad (v_0 = 0)$$

U_v as qudit version of $U_{\pi/8}$ $(p = 3)$

$$\text{For } p = 3 : \quad \det \left(\zeta^{2\gamma'} C \left(\left[\begin{array}{c|c} 1 & 0 \\ \gamma' & 1 \end{array} \right] \middle| \left[\begin{array}{c} 1 \\ z' \end{array} \right] \right) \right) = 1 \quad (\zeta = e^{\frac{2\pi i}{9}})$$

$$\Rightarrow U_v = \sum_{k=0}^2 \zeta^{v_k} |k\rangle\langle k| \quad (v_k \in \mathbb{Z}_9)$$

$$v = (v_0, v_1, v_2) = (0, 6z' + 2\gamma' + 3\epsilon', 6z' + \gamma' + 6\epsilon') \pmod{9}$$

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Examples:

$$p = 3 : \begin{pmatrix} 1 & 0 & 0 \\ 0 & e^{\frac{2\pi i}{9}} & 0 \\ 0 & 0 & e^{-\frac{2\pi i}{9}} \end{pmatrix}_{[z'=1, \gamma'=2, \epsilon'=0]} \quad p = 5 : \begin{pmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & e^{-\frac{4\pi i}{5}} & 0 & 0 & 0 \\ 0 & 0 & e^{-\frac{2\pi i}{5}} & 0 & 0 \\ 0 & 0 & 0 & e^{\frac{4\pi i}{5}} & 0 \\ 0 & 0 & 0 & 0 & e^{\frac{2\pi i}{5}} \end{pmatrix}_{[z'=1, \gamma'=4, \epsilon'=0]}$$

Group Structure of U_v gates

Can create

- p^3 gates $U_v(z', \gamma', \epsilon')$ varying over $z', \gamma', \epsilon' \in \mathbb{Z}_p$.
- $p^2(p-1)$ non-Clifford U_v varying over $z', \epsilon' \in \mathbb{Z}_p, \gamma' \in \mathbb{Z}_p^*$.

$$U_v D_{(1|0)} U_v^\dagger = \omega^{\epsilon'} C\left(\begin{bmatrix} 1 & 0 \\ \gamma' & 1 \end{bmatrix} \middle| \begin{bmatrix} 1 \\ z' \end{bmatrix}\right) \quad \text{but} \quad C\left(\begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \middle| \begin{bmatrix} 1 \\ z' \end{bmatrix}\right) = D_{(1|z')}$$

Easy to show p^3 gates $\{U_v\}$ form a finite Abelian group.

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Use Fund. Thm. of finite Abelian groups to characterize $\{U_v\}$

	Group name	No. elements of order				Min. no. of generators
		1	p	p^2	p^3	
$p = 2$	\mathbb{Z}_8	1	1	2	4	1
$p = 3$	$\mathbb{Z}_9 \times \mathbb{Z}_3$	1	8	18	0	2
$p > 3$	\mathbb{Z}_p^3	1	$p^3 - 1$	0	0	3

Group Structure

“All primes are odd except two, which is the oddest prime of all!”

Can create

- p^3 gates $D_{(1|0)}$ (ϵ') varying over $z', \gamma', \epsilon' \in \mathbb{Z}_p$.
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Easily check that $\{U_v\}$ form a finite Abelian group.

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Easily

$\{U_v\}$

Perhaps, in Quantum Information, three is the **second oddest** prime!

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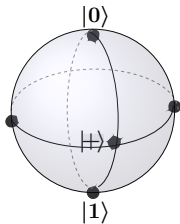
Easy to show p^3 gates $\{U_v\}$ form a finite Abelian group.

$$U_v(z_1, \gamma_1, \epsilon_1) U_v(z_2, \gamma_2, \epsilon_2) = U_v(z_1 + z_2, \gamma_1 + \gamma_2, \epsilon_1 + \epsilon_2)$$

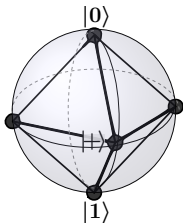
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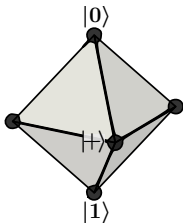
Qubit Geometry: Magic States and $U_{\pi/8}$



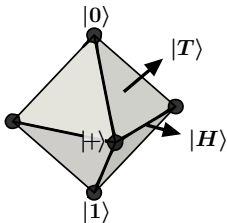
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Qubit Geometry: Magic States and $U_{\pi/8}$



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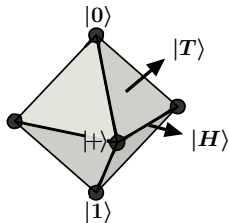


$$|T\rangle\langle T| = \frac{1}{2} \left(\mathbb{I} + \frac{\sigma_x + \sigma_y + \sigma_z}{\sqrt{3}} \right)$$

$$|H\rangle\langle H| = \frac{1}{2} \left(\mathbb{I} + \frac{\sigma_x + \sigma_y}{\sqrt{2}} \right)$$

- Both $|T\rangle$ and $|H\rangle$ are eigenvectors of Clifford gates
- $|T\rangle$ is the most non-stabilizer qubit state
- $|H\rangle$ is the most non-stabilizer qubit state in the equatorial plane
 - Note that $|H\rangle = U_{\pi/8}|+\rangle = \frac{1}{\sqrt{2}}\text{diag}(U_{\pi/8})$

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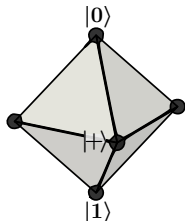
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 - Note that $|H\rangle = U_{\pi/8}|+\rangle = \frac{1}{\sqrt{2}}\text{diag}(U_{\pi/8})$
- In [vDH'11], states $\frac{|0\rangle - |1\rangle}{\sqrt{2}}$ were found to be maximally non-stabilizer, and Clifford eigenvectors, in all odd prime dimensions (similar to $|T\rangle$?)
- We will argue that $|\psi_{U_v}\rangle \equiv U_v|+\rangle$ is the qudit analogue of $|H\rangle$.



W. van Dam and M. Howard,

Noise thresholds for higher-dimensional systems using the discrete Wigner function
 Phys. Rev. A. **83**, 032310, (2011).

Geometry: $|\psi_{U_v}\rangle$ as the qudit analogue of $|H\rangle$

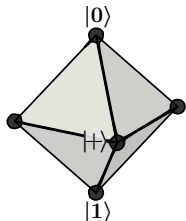


$STAB =$ Convex hull of qudit stabilizer states

$$= \left\{ \rho \mid \rho = \sum_{i=1}^{p(p+1)} q_i |\psi_{STAB}^{(i)}\rangle \langle \psi_{STAB}^{(i)}|, \sum_{i=1}^{p(p+1)} q_i = 1 \right\}$$

$$= \left\{ \rho \mid \min_{u \in \mathbb{Z}_p^{p+1}} \text{Tr}[A(u)\rho] \geq 0 \right\}$$

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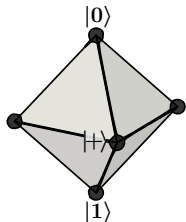


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 \end{aligned}$$

Define $\begin{cases} U_\theta = e^{i\theta_k} |k\rangle \langle k| \ (\theta_k \in \mathbb{R}) \\ |\psi_{U_\theta}\rangle = \frac{e^{i\theta_k}}{\sqrt{p}} |k\rangle = U_\theta |+ \rangle \end{cases}$ so that $\begin{cases} \{U_v\} \subset \{U_\theta\} \\ \{|\psi_{U_v}\rangle\} \subset \{|\psi_{U_\theta}\rangle\} \end{cases}$

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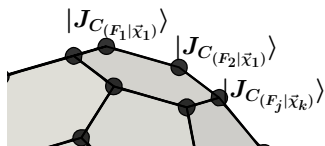
- $|\psi_{U_v}\rangle$ is farthest outside $STAB$ of all $|\psi_{U_\theta}\rangle$ (for $p = 2, 3, 5, 7$ at least)
- $|\psi_{U_v}\rangle$ is also a Clifford eigenvector:

$$C\left(\begin{bmatrix} 1 & 0 \\ \gamma' & 1 \end{bmatrix} \middle| \begin{bmatrix} 1 \\ z' \end{bmatrix}\right) |\psi_{U(z', \gamma', \epsilon')}\rangle = \omega^{-\epsilon'} |\psi_{U(z', \gamma', \epsilon')}\rangle.$$

Geometry: U_v as the qudit analogue of $U_{\pi/8}$

$$\text{For } \begin{cases} U \in U(p) \\ \mathcal{E} : \rho_{in} \mapsto \rho_{out} \end{cases} \quad \text{Jamiołkowski state} = \begin{cases} |J_U\rangle \equiv (\mathbb{I} \otimes U) \sum_{j=0}^{p-1} \frac{|jj\rangle}{\sqrt{p}} \\ \varrho_{\mathcal{E}} = [\mathcal{I} \otimes \mathcal{E}] \left(\sum_{j,k=0}^{p-1} \frac{|j,j\rangle\langle k,k|}{p} \right) \end{cases}$$

\mathcal{CLIFF} = Convex hull of qudit Clifford gates



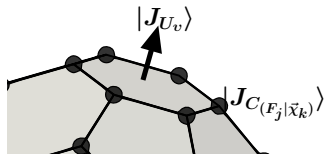
$$= \left\{ \varrho_{\mathcal{E}} \mid \varrho_{\mathcal{E}} = \sum_{j,k=1}^{j=p(p^2-1), k=p^2} q_{j,k} |J_{C(F_j|\vec{x}_k)}\rangle\langle J_{C(F_j|\vec{x}_k)}| \right\}$$

$$= \left\{ \varrho_{\mathcal{E}} \mid \min_{W \in \mathcal{W}} \text{Tr} [W \varrho_{\mathcal{E}}] \geq 0 \right\}$$

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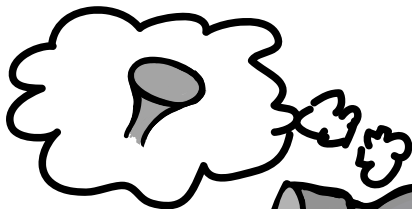
- Seems that U_v is the most non-Clifford $U \in U(p)$ (for $p = 2, 3, 5, 7, ?$)



W. van Dam and M. Howard,

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 Phys. Rev. A. **83**, 032310, (2011).

Applications?



$$U_{\pi/8} = \begin{pmatrix} e^{-i\frac{\pi}{8}} & 0 \\ 0 & e^{i\frac{\pi}{8}} \end{pmatrix} \propto \begin{pmatrix} 1 & 0 \\ 0 & e^{i\frac{\pi}{4}} \end{pmatrix}$$

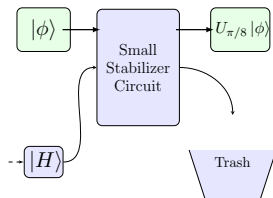


- UQC = $\langle \text{Cliffords}, U_{\pi/8} \rangle$
- Teleportation-based UQC
- Transversal for some Reed-Muller codes
- Secure assisted UQC

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- Measurement-based UQC with graph states
 - Optimal CHSH game with $(|00\rangle + |11\rangle)/\sqrt{2}$
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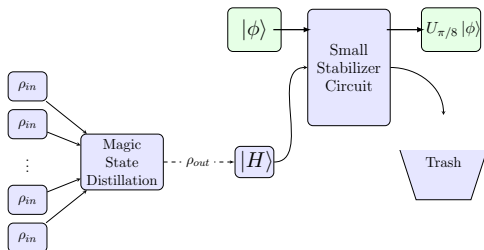
Application to Fault-tolerant QC

- We argued $|\psi_{U_v}\rangle$ was the qudit analogue of $|H\rangle = |\psi_{U_{\pi/8}}\rangle$.
- The key feature of $|H\rangle$ is that it is suitable for magic state distillation



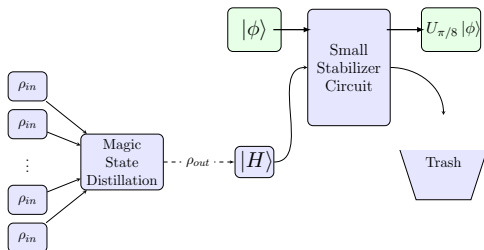
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- Campbell, Anwar and Browne have shown $|\psi_{U_v}\rangle$ are qudit magic states.
- Nebe, Rains and Sloane have shown $\langle \text{Cliffords}, U_v \rangle$ enables UQC
- Geometrical features are good news!



E. T. Campbell, H. Anwar and D. E. Browne,

Magic state distillation in all prime dimensions using quantum Reed-Muller codes
arXiv:1205.3104v1, (2012).

UQC using perfect Cliffords + noisy U_v gates

$$\mathcal{E}_{(|\psi_{U_v}\rangle, \varepsilon)}(\rho) \equiv (1 - \varepsilon)U_v\rho U_v^\dagger + \varepsilon\frac{\mathbb{I}}{p} \quad (\varepsilon \approx \text{depolarizing error rate})$$

- For what noise rates, ε , does $\mathcal{E}_{(|\psi_{U_v}\rangle, \varepsilon)}$ + Cliffords enable UQC?

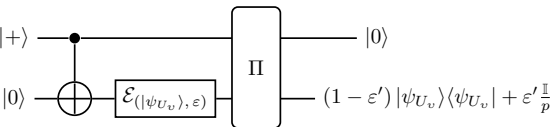
$$|+\rangle \text{ --- } \boxed{\mathcal{E}_{(|\psi_{U_v}\rangle, \varepsilon)}} \text{ --- } (1 - \varepsilon)|\psi_{U_v}\rangle\langle\psi_{U_v}| + \varepsilon\frac{\mathbb{I}}{p}$$

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- Postselection “dilutes” the noise: $\varepsilon' < \varepsilon$.

Exact relationship:

$$\varepsilon' = \frac{\varepsilon}{p - (p - 1)\varepsilon}$$

- Use this result, along with routines in [CAB'12], to find allowable ε

Noise thresholds for UQC

	Lower Bound	Upper Bound
$p = 2$	45.32%	45.32%
$p = 3$	58.15%	78.63%
$p = 5$	80.61%	95.20%
$p = 7$	72.24%	97.63%

Table: Bounds on threshold ε for UQC using noisy U_v + ideal Cliffords

How are these values found?

Lower bound:

- Postselection circuit & the best performing MSD routines given in [CAB'12].

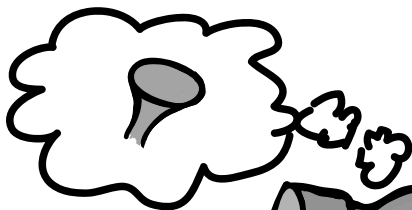
Upper bound:

- Explicit facets of \mathcal{CLIFF} for which $\text{Tr} \left(W \left[(1 - \varepsilon) |J_U\rangle\langle J_U| + \varepsilon \frac{\mathbb{I}}{p^2} \right] \right) = 0$

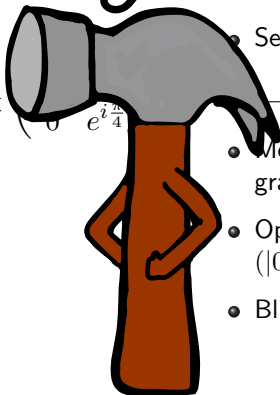
Note: U_v also maximally robust to phase damping noise ($p = 2, 3, 5, 7, ?$)

Open Question: Can the gaps be closed?

Applications?



$$U_{\pi/8} = \begin{pmatrix} e^{-i\frac{\pi}{8}} & 0 \\ 0 & e^{i\frac{\pi}{8}} \end{pmatrix} \propto \begin{pmatrix} 1 & 0 \\ 0 & e^{i\frac{\pi}{4}} \end{pmatrix}$$

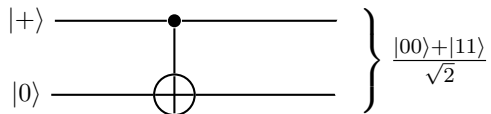


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CHSH Bell Inequality

$$\langle A_0 B_0 \rangle + \langle A_0 B_1 \rangle + \langle A_1 B_0 \rangle - \langle A_1 B_1 \rangle \leq 2 \quad (\lambda(A_j), \lambda(B_k) = \pm 1)$$



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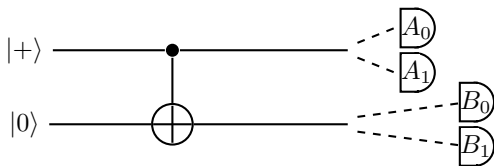
$$(\lambda(A_j), \lambda(B_k) = \pm 1)$$

$$A_0 = X$$

$$A_1 = Y$$

$$B_0 = (X - Y)/\sqrt{2}$$

$$B_1 = (X + Y)/\sqrt{2}$$



- $2\sqrt{2} \not\leq 2$

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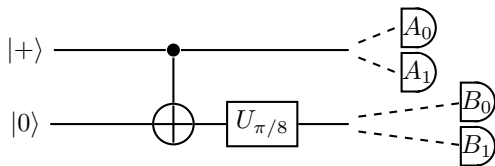
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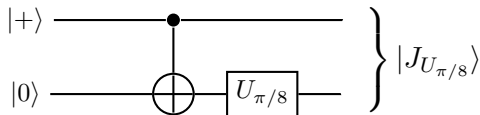
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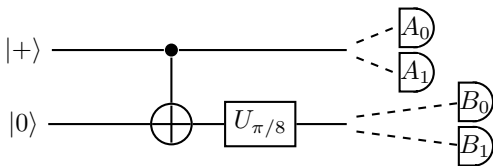
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Maximizing quantity $\langle \mathcal{B} \rangle$ is related to maximizing

$$\sum_{\substack{a+b=st \pmod{2} \\ (a,b,s,t \in \mathbb{Z}_2)}} p(a, b|s, t)$$

where $p(a, b|s, t)$ is a conditional prob.
settings
outcomes

CHSH Bell Inequality

$$\langle \mathcal{B} \rangle \leq 2 \quad \mathcal{B} = \sum_{j,k \in \mathbb{Z}_2} (-1)^{jk} A_j B_k$$

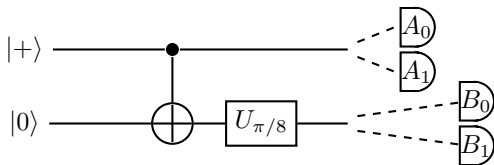
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$$\langle \mathcal{B} \rangle \leq 2 \quad \mathcal{B} = \sum_{j,k \in \mathbb{Z}_2} (e^{\frac{2\pi i}{2}})^{jk} A_j B_k$$

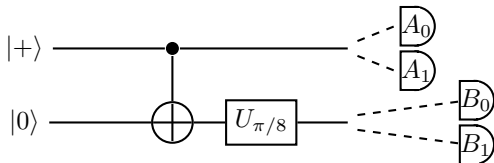
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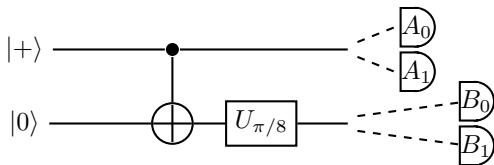
$$\langle \mathcal{B} \rangle \leq 2 \quad \mathcal{B} = \sum_{j,k \in \mathbb{Z}_2} \omega^{jk} A_j B_k \quad (\lambda(A_j), \lambda(B_k) = \pm 1)$$

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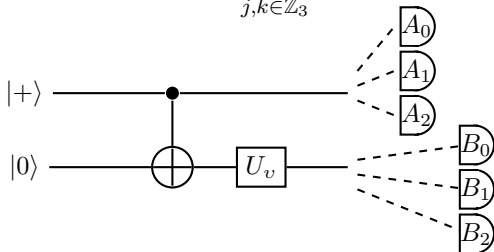
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Generalized CHSH Bell Inequality for $p = 3$

$$\langle \mathcal{B} \rangle \leq 4.5 \quad \mathcal{B} = \sum_{\substack{n \in \mathbb{Z}_3^* \\ j, k \in \mathbb{Z}_3}} \omega^{njk} A_j^n B_k^n \quad (\lambda(A_j), \lambda(B_k) = \{\omega^0, \omega^1, \omega^2\})$$

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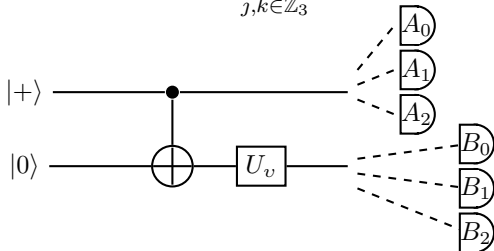
$$A_j = \omega^{j(j+1)} X Z^j$$

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- $\langle \mathcal{B} \rangle = 5.117 \not\leq 4.5$

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Generalized CHSH Bell Inequalities

Originates with Ji *et al.*, Reexamined by Liang *et al.*



Se-Wan Ji, Jinhyoung Lee, James Lim, Koji Nagata, and Hai-Woong Lee

Multisetting Bell inequality for qudits,

Phys. Rev. A 78, 052103 (2008).



Yeong-Cherng Liang, Chu-Wee Lim, Dong-Ling Deng

Reexamination of a multisetting Bell inequality for qudits,

Phys. Rev. A 80, 052116 (2009).

Observation 1: For all the cases ($p \leq 17$) that we have checked, U_v are optimal in the sense:

$$\text{Tr}(|J_{U_v}\rangle\langle J_{U_v}| \mathcal{B}) = \lambda_{\max}(\mathcal{B})$$

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Observation 2: Suggests a natural generalization to $p \geq 3$ of the qubit result relating capability of operation \mathcal{E} + stabilizer ops

Violation of
CHSH
inequality \Rightarrow Universal
QC
(via MSD)



M. Howard and J. Vala,

Nonlocality as a benchmark for universal quantum computation in Ising anyon topological quantum computers,

Phys. Rev. A 85, 022304 (2012).

Open Questions & Thanks

- Are qudits better in any way?
 - How does one fairly compare qubits and qudits?
- Physical system that enables topologically protected qudit Cliffords?
(cf. Ising anyons for $p = 2$)
- What does the rest of \mathcal{C}_3 look like?
- What does the diagonal subset of \mathcal{C}_k look like?
- Can we close the gap between upper and lower bounds on noise thresholds?
- Are there applications in MUBs, SICs, Unitary designs?
- Can stronger statements be made relating nonlocality and UQC in the Clifford computer/magic state model of computation?

Thanks to Earl Campbell for many helpful discussions & comments

