

# Experimental design for expensive path planning simulators via integer programming

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# What are path planning problems?

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## Goal

*Finding an **optimal navigation path** that minimizes a cost function while meeting specified constraints*

e.g. warehouse robotics, rover navigation tasks

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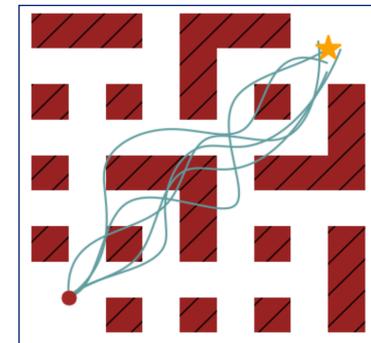
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Amazon robots (source: [amazon.com](https://www.amazon.com))



A rover navigation task (obstacles are swamps!) (Wang et al. 2018)

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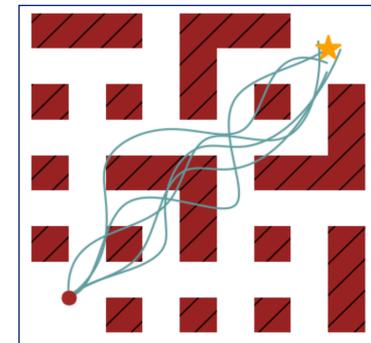
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Amazon robots (source: [amazon.com](https://www.amazon.com))

## Discrete path planning problems

- Action spaces are **discrete**
- More economical; aligns with many robotics tasks
- e.g. maze-solving problems (up, down, right, left)



A rover navigation task (obstacles are swamps!) (Wang et al. 2018)

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- Particle swarm optimization (*Kennedy and Eberhart, 1995*)
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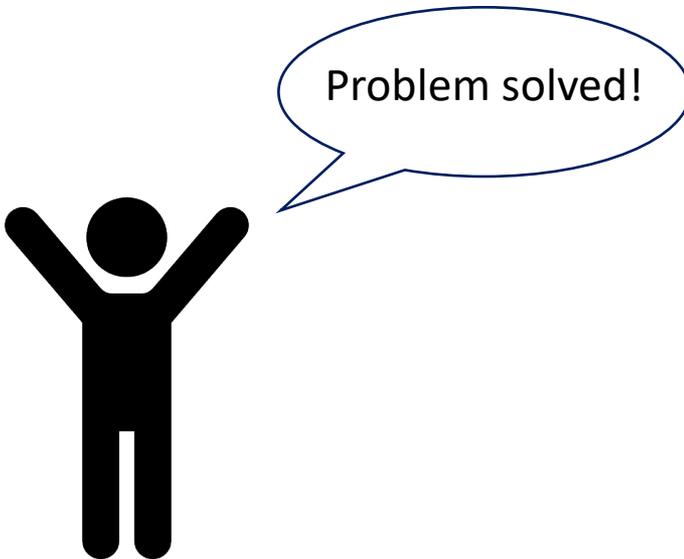
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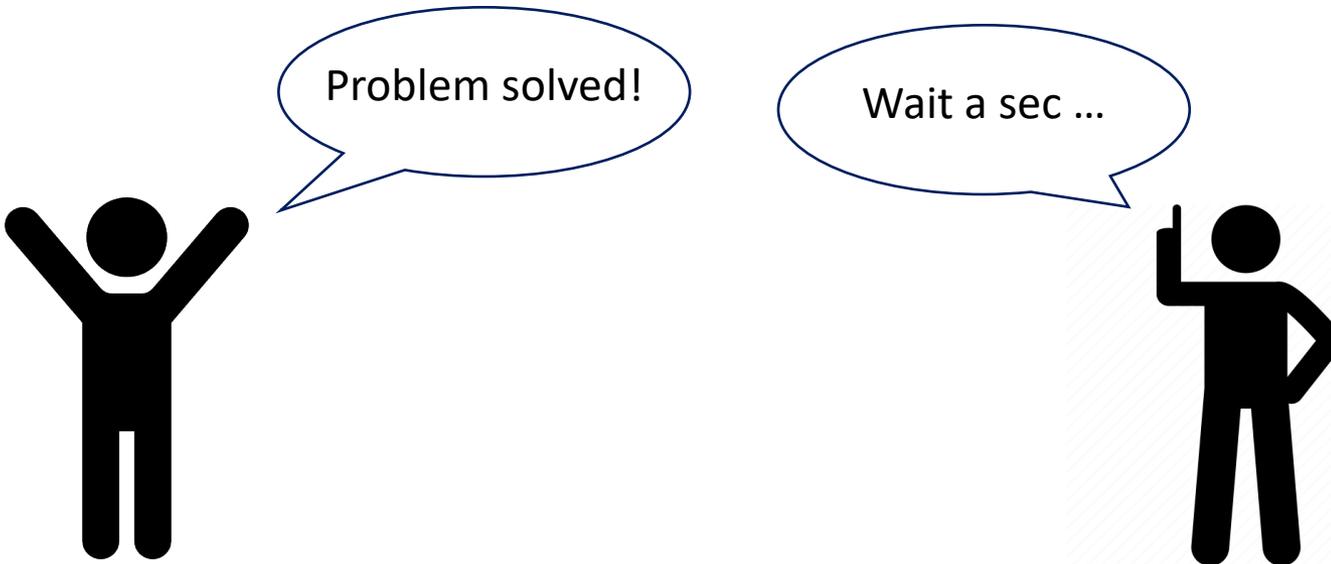


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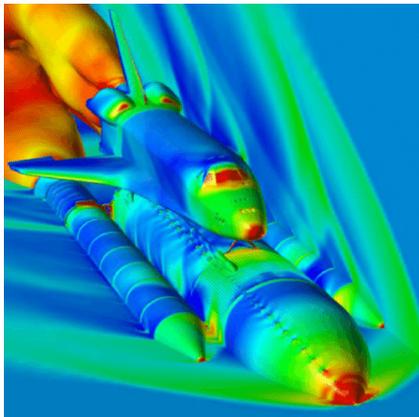
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# Experiments can be expensive!

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- **Path planning** is now needed for **highly sophisticated** systems, e.g., space rover, unmanned aerial vehicle, ... etc.
- High-fidelity simulators can be **computationally expensive (days)**!



Rocket propulsion



unmanned aerial vehicle  
(Narayanan et al. 2024)

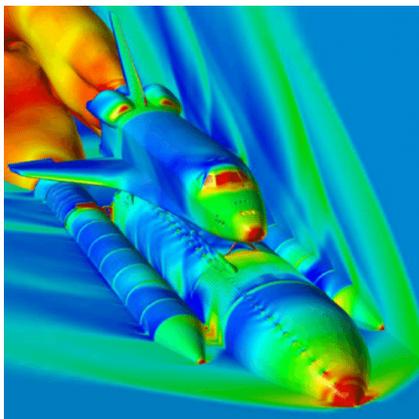
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We should be more careful when *designing initial runs* and *selecting sequential runs*.

Especially when the problem is **high-dimensional**!



Rocket propulsion



unmanned aerial vehicle  
(Narayanan et al. 2024)

# In this talk

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## Goals:

- Construct **optimal space-filling initial designs** via integer programming (IP)
- Propose simple, straightforward IP-based **sequential designs for prediction and optimization** that **avoids exhaustive search**

**Action space:**  $q$  decisions to make;  $m$  different choices of actions

$$\mathcal{A} = \{x; x \in \{a_1, \dots, a_m\}^q\}$$

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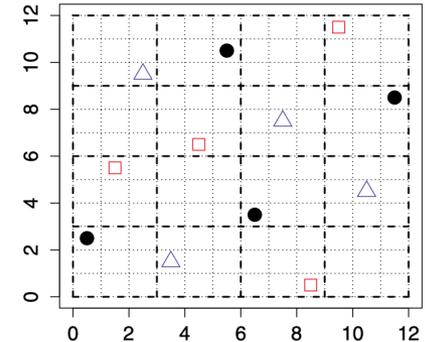
both  $q$  and  $m$  can be high-dimensional!

# Literature Review

# Popular experimental designs

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- **Orthogonal Arrays (restricted run sizes)**
- **Latin hypercube designs**
  - Sliced-LHD (exchangeable algorithms)  
(*Ba et al., 2015*)
- **Maximin designs**



Example of sliced-LHD (*Ba et al., 2015*)

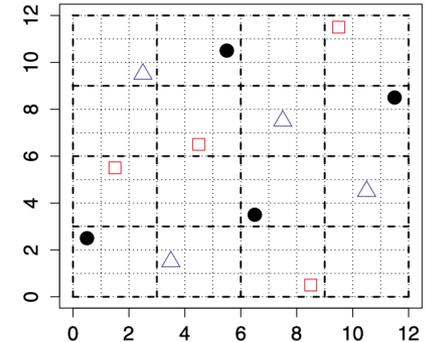
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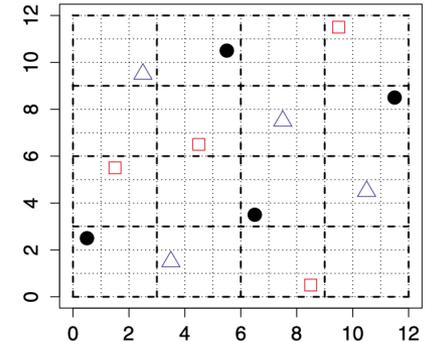
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# Qualitative Gaussian Processes

---

## Hamming-distance based kernel (Qian et al. 2008)

- Hamming distance

$$d(x_i, x_j) = \sum_{k=1}^q \mathbf{1}(x_{ik} \neq x_{jk})$$

- Exchangeable kernel

$$k_{\theta} (x_i, x_j) = \tau^2 \exp \left\{ - \sum_{l=1}^q \theta_l \mathbf{1}(x_{il} \neq x_{jl}) \right\}$$

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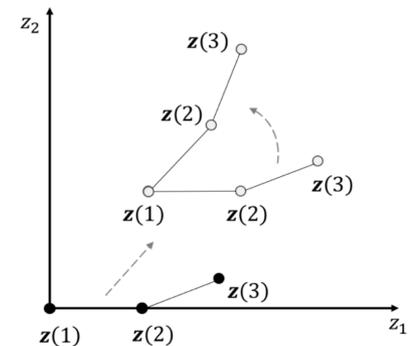
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## Latent variable kernel (Zhang et al. 2020)

- **Idea**: mapping qualitative variables to *latent continuous spaces* and applying continuous kernels
- **Assumption**: there *exists* latent continuous spaces for qualitative variables
- **Applications**: material science, physics-based simulation



LVGP (Zhang et al. 2020)

# Sequential design for qualitative variables

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## Exhaustive search

- Pros: ensured to achieve global optimum
- Concerns: infeasible in high-dim cases or when the decision space is large

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- **Concerns**: dummy variables are not really quantitative variables; can be problematic

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## Latent variable Bayesian optimization (Zhang et al. 2020b)

- **Pros:** flexible, efficient, nice UQ
- **Concerns:** can be slow in high-dim cases; need sufficient training samples, latent space assumption might not hold

# Proposed methods

# Experimental designs via IP

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- **Convert inputs to dummy variables**

e.g. Suppose action set = {a, b, c, d} and design = (b,c,a,d,a). Then

$$I(b, c, a, d, a) = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \end{pmatrix}.$$

- Each  $I(x)$  must satisfy  $\sum_{j=1}^m I(x)_{i,j} = 1$  for all  $i$ .

- Hamming distance between designs can be represented as

$$d(x_i, x_j) = q - \text{tr}\{I(x_i)I(x_j)^\top\}.$$

# Experimental designs via IP

---

Solving a **maximin design** with minimal Hamming distance  $\geq C$  is the same as solving:

$$\begin{aligned} \max_{\mathcal{X}_n} 1, \quad \text{subject to} \quad & q - \sum_{k=1}^q (I_i I_{i'}^\top)_{kk} \geq C, \quad \forall 1 \leq i < i' \leq n \\ & I_i \mathbf{1}_m - \mathbf{1}_q = \mathbf{0}_q, \quad \forall 1 \leq i \leq n, \quad I_i \in \{0,1\}^q \times \{0,1\}^m \quad \forall 1 \leq i \leq n. \end{aligned}$$

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## Advantages

1. Guaranteed **convergence to global optimum** (IP+AP)
2. Can be solved by **state-of-art IP optimizers** (Gurobi)
3. Providing **optimization gaps** (not available in heuristic algorithms)

# How to select initial threshold $C$ ?

## The Gilbert–Varshamov bound (Gilbert, 1952; Varshamov, 1957)

Let  $N_m(q, s)$  denote the maximum size of path planning designs with action space  $(m, q)$  and minimum pairwise Hamming distance  $s$ . Then,

$$N_m(q, s) \geq \frac{m^q}{\sum_{j=0}^{s-1} \binom{q}{j} (q-1)^j}.$$

## Proposition (Initial threshold $C$ )

By Gilbert-Varshamov bound, a lower bound threshold  $S(m, q, n)$  of DPP designs of size  $n$  and action space  $(m, q)$  can be estimated by

$$S(m, q, n) = \arg \max_{s \in \{1, 2, \dots, q\}} \left\{ m^q / \sum_{j=0}^{s-1} \binom{q}{j} (q-1)^j \geq n \right\}.$$

# Constructing initial designs $D(m, q, n)$

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Use existing maximin design  $D(m_0, q_0, n_0)$ , where  
 $m_0 \leq m, q_0 \leq q, n_0 \leq n$  as warmstarts



optional

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**optional**

Search for a design with minimal pairwise distance  $C = S(m, q, n)$

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If design exists

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If design exists

Search for a design with minimal pairwise distance =  $C$

↓ if design **not** exists or  $C+1 = q-1$

Return design

# Sequential Prediction

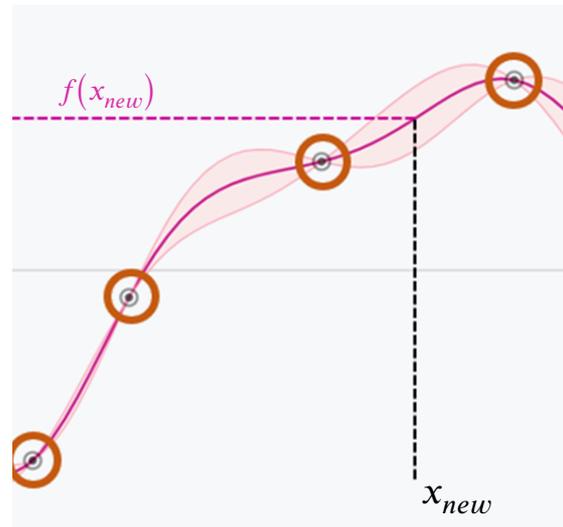
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**Goal:** optimizing prediction performance

**Active learning Mackay** (ALM; MacKay, 1992) for prediction

Explore the unknown!

$$x_{new} = \arg \max_x \sigma_n^2(x)$$



# IP-based ALM for qualitative kernel

---

Let  $\mathcal{X}_n = (X_1, X_2, \dots, X_n)$  be the observed data and  $\{\theta_k\}_{k=1}^q$  be the set of lengthscale parameters.

$$\begin{aligned}
 \max_x \sigma_n^2(x) &= \max_x \tau^2 \{K(x, x) - K(x, \mathcal{X}_n)K_{nn}^{-1}K(\mathcal{X}_n, x)\} \\
 &= \min_{I(x)} K(x, \mathcal{X}_n)K_{nn}^{-1}K(\mathcal{X}_n, x) \\
 &= \min_{I(x)} \sum_{i=1}^n \sum_{j=1}^n \exp \left[ - \sum_{k=1}^q \left\{ I(x)I(X_i)^\top \right\}_{kk} \theta_k \right] (K_{nn}^{-1})_{ij} \exp \left[ - \sum_{k=1}^q \left\{ I(x)I(X_j)^\top \right\}_{kk} \theta_k \right] \\
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subject to

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**A mixed integer programming problem (MIP)!**

**Able to achieve global optimum without exhaustive search!**

# Sequential Optimization

---

**Goal:** finding  $x^* = \arg \max f(x)$

## Upper confidence bound (UCB)

Exploration - exploitation

$$\max_{I(x)} \mu_n \{I(x)\} + \lambda \sigma_n \{I(x)\}$$

Posterior mean: 
$$\mu_n \{I(x)\} = \hat{\mu}_n + \sum_{i=1}^n \sum_{j=1}^n \exp \left[ - \sum_{k=1}^q \{I(x)I(X_i)^\top\}_{kk} \theta_k \right] (K_{nn}^{-1})_{ij} (y_{n,j} - \hat{\mu}_n)$$

Posterior variance: 
$$\sigma_n^2 \{I(x)\} = \hat{\tau}^2 \left\{ 1 - \exp \left( - \sum_{k=1}^q \left[ I(x) \{ I(X_i) + I(X_j) \}^\top \right]_{kk} \theta_k \right) (K_{nn}^{-1})_{ij} \right\}$$

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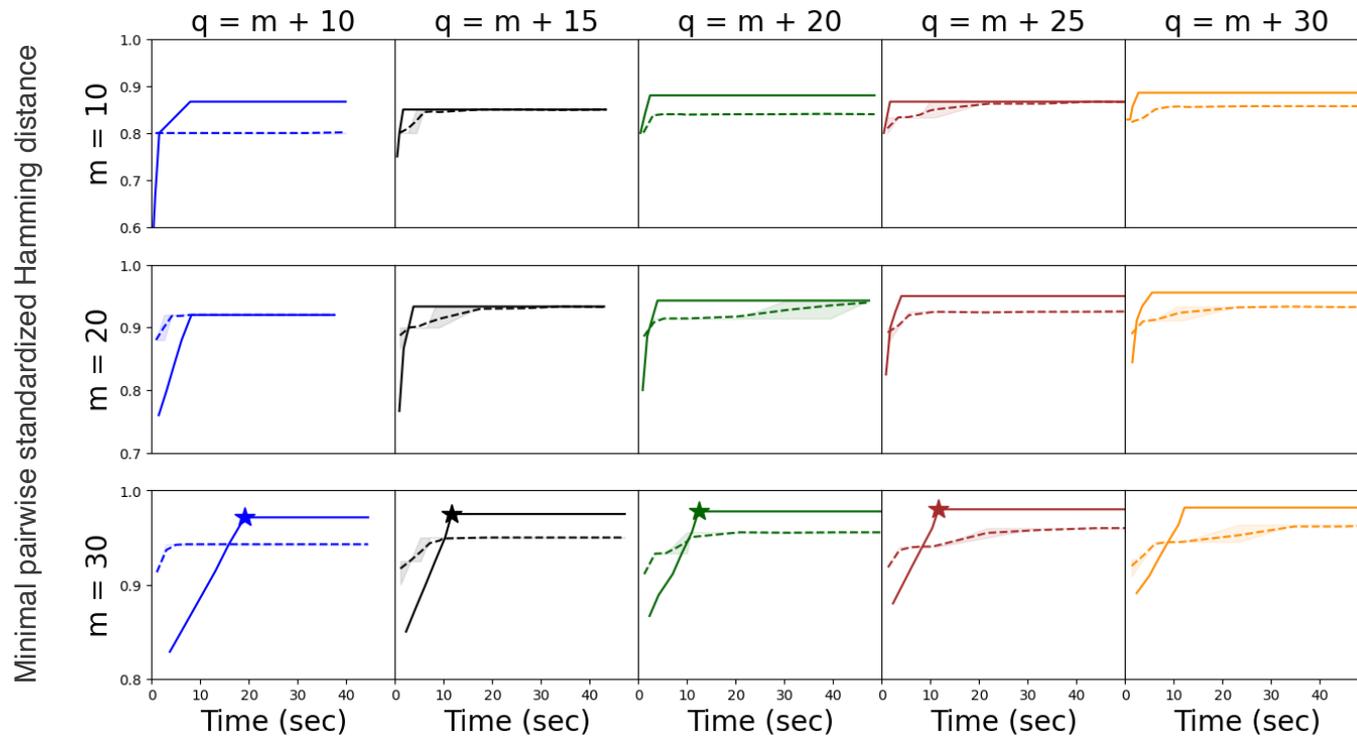
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# Numerical Results

# Maximin designs

Maximin distance achieved by **proposed IP optimizer (Gurobi)** v.s. **metaheuristic algorithms** (Stokes et al. ,2023) when design size is 50.



**solid lines:** IP optimizer (Gurobi)

**dashed lines:** metaheuristic algorithm (differential evolution) by Stokes et al. (2023)

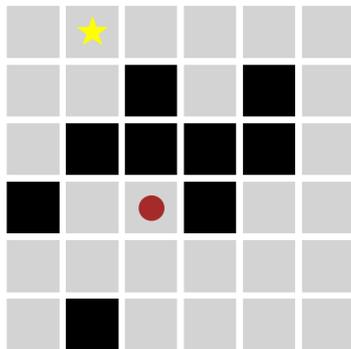
# Simulation - Maze problems

## Goal:

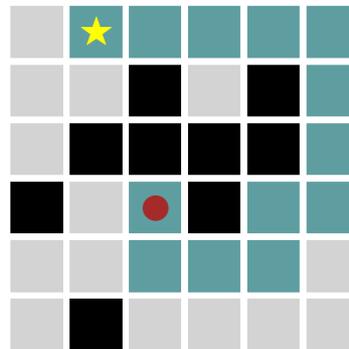
Predict the number of steps needed to reach the goal after taking  $q$  steps.

## Action space ( $m^q = 5^7, 5^{12}$ ):

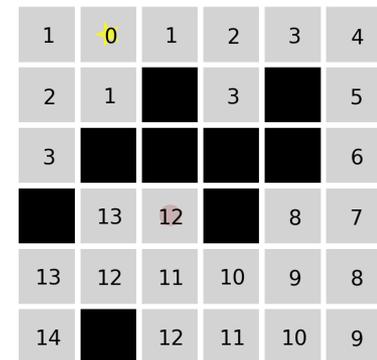
- Number of actions ( $m$ ) = 5 (left, right, up, down, stay)
- Step lengths ( $q$ ) = 7 or 12



A 6x6 maze

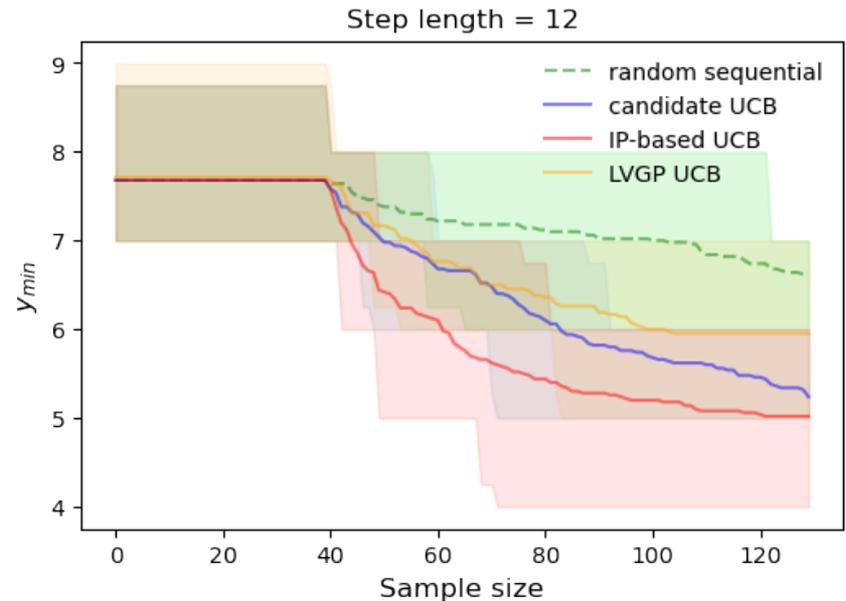
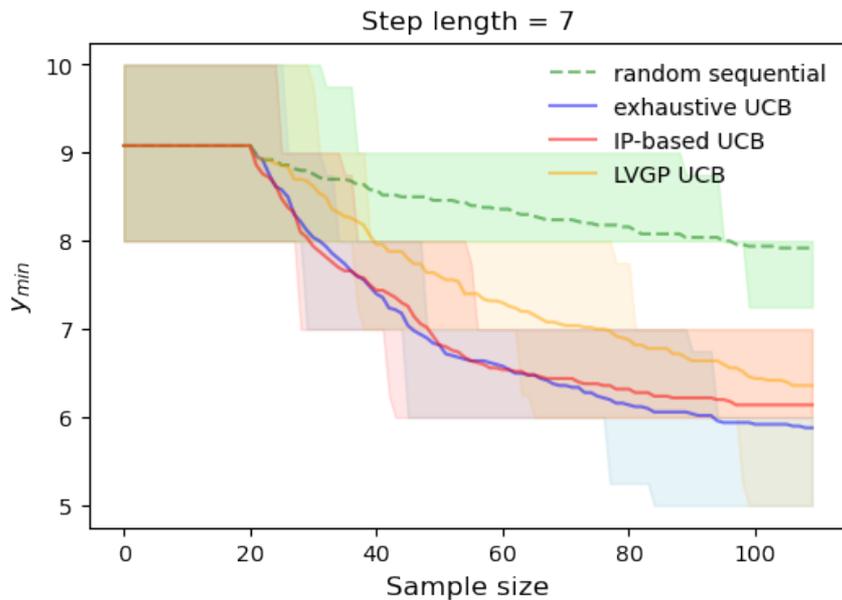


optimal path



Cost values at each grid

# Sequential optimization



Maximin initial designs were obtained by IP optimizers

# Sequential prediction

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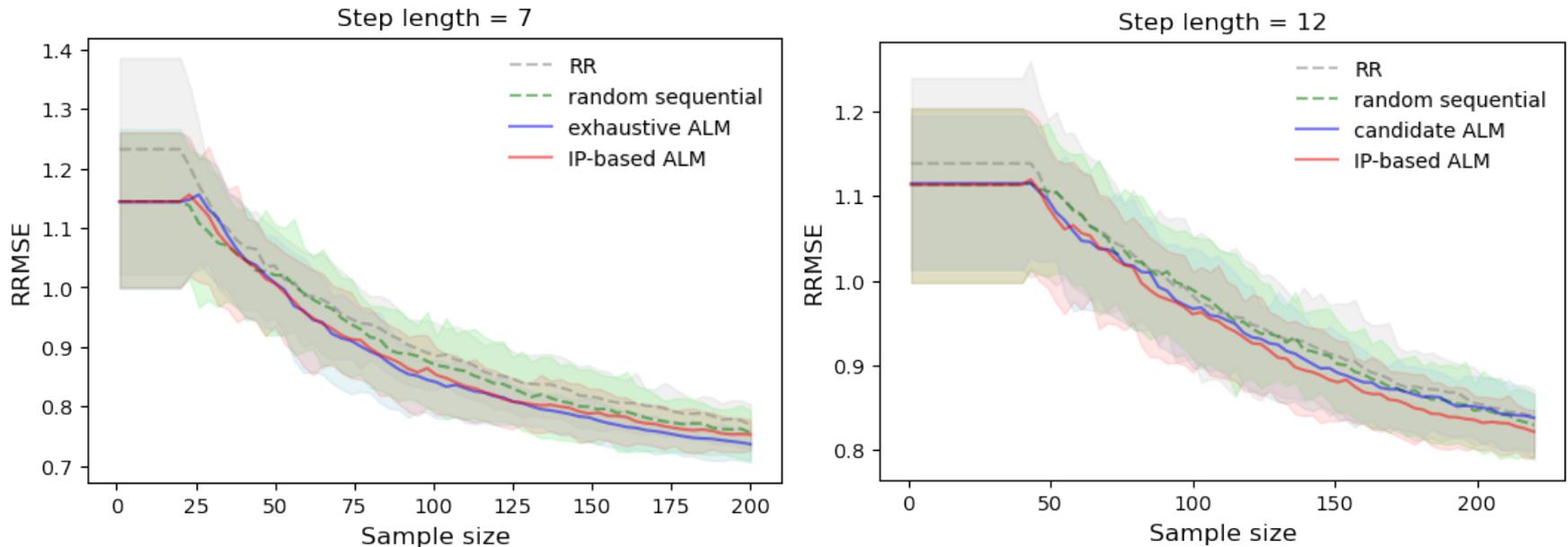
Evaluation criteria: RRMSE

$$\text{RRMSE} = \sqrt{\frac{N^{-1} \sum_{i=n+1}^{n+N} (\hat{y}_i - y_i)}{(n-1)^{-1} \sum_{i=1}^n (y_i - \bar{y}_n)^2}}.$$

# Sequential prediction

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- RR: random initial design + random sequential design
- Others: use maximin initial design found by IP optimizer

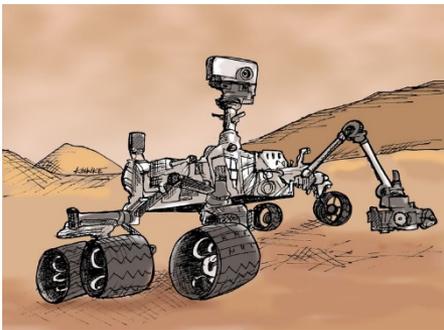
# Rover navigation task

## Goal:

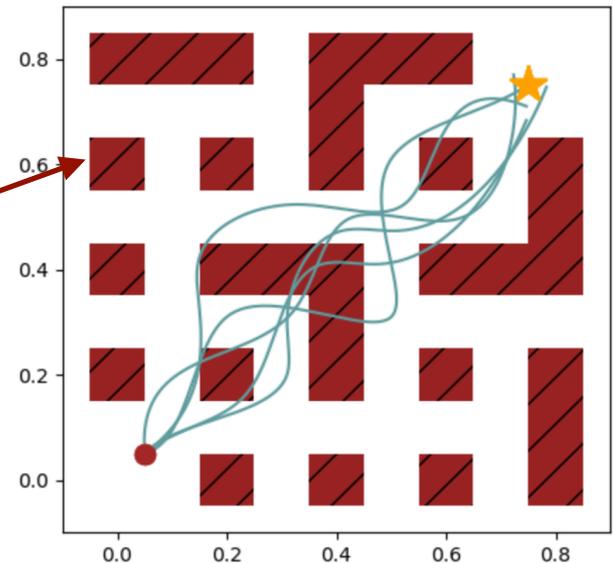
Reaching the goal in limited steps and minimizing the time trapped in the swamps

## Action space ( $m^q = 9^8$ ):

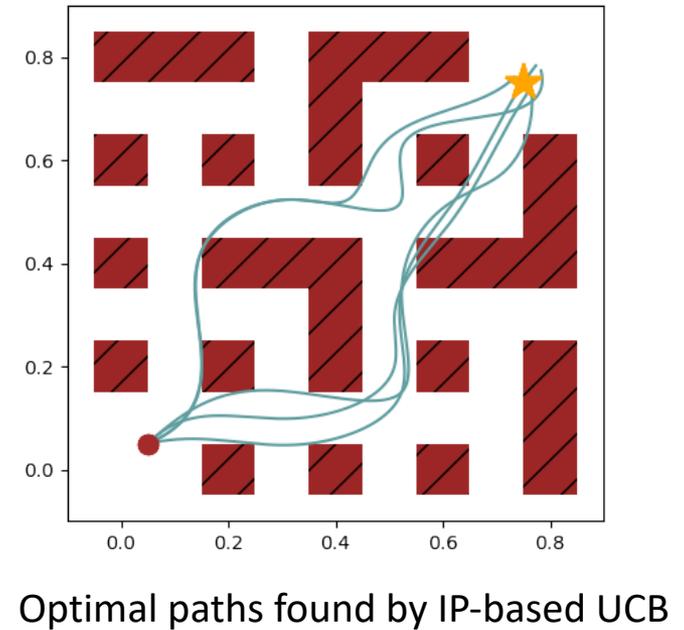
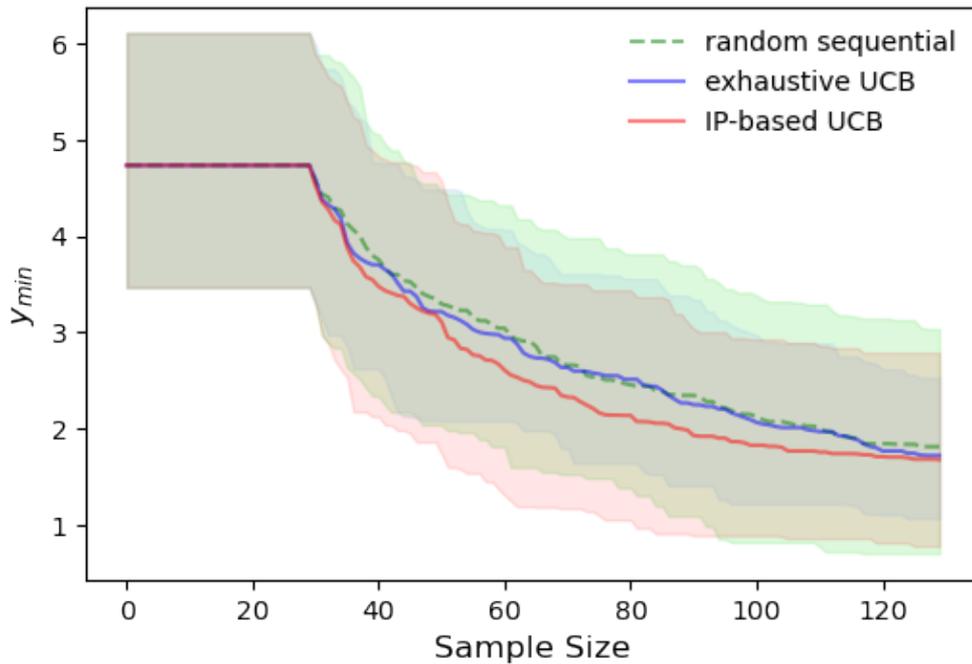
- Number of actions ( $m$ ) = 9:  
4 angles ( $0, \pi/6, \pi/4, \pi/2$ )  $\times$  2 lengths (0.1, 0.2) + 1 (stay)
- Step lengths ( $q$ ) = 8



**Swamps**  
(can pass through with greatly reduced speed)



# Sequential optimization



Thank you!  
Any questions?

