## Note

# Labelled trees and factorizations of a cycle into transpcsitions 

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#### Abstract

Goulden, I.P. and S. Pepper, Labelled irees and factorizations of a cycle into transpositions, Discrete Mathematics 113 (1993) 263-268. Moszkowski has previously given a direct bijection between labelled trees on $n$ vertices and factorizations of the cycle $(12 \cdots n)$ in $S_{n}$ into $n-1$ transpositions. By considering a quadratic recurrence equation and its combinatorial interpretation for trees and for transposition factorizations, we derive another such bijection in a straightforward manner.


Let $A_{n}$ be the set of labelled trees on $n$ vertices and $B_{n}$ be the set of ( $n-1$ )-tuples of transpositions in $S_{n}$ whose ordered product is the cycle $(12 \cdots n)$. Cayley [1] proved that $\left|A_{n}\right|=n^{n-2}$, and Dénes [2] proved that $\left|B_{n}\right|=\left|A_{n}\right|$, by giving a bijection between sets of cardinality $(n-1)!\left|A_{n}\right|$ and $(n-1)!\left|B_{n}\right|$. Dénes posed the problem of finding a direct bijection between $A_{n}$ and $B_{n}$; the first such bijection was given by Moszkowski [4].
Jackson [3] was able to enumerate factorizations in $S_{n}$ in several more general situations. The method was based on symmetric group characters, and yielded simple binomial summations as solutions. For example, the number of $k$-tuples of transpositions in $S_{n}$ whose ordered product is the cycle $(12 \cdots n)$ was shown to be

$$
t(n, k)=\frac{1}{n!}\left(\frac{n}{2}\right)^{k} \sum_{i=0}^{n-1}(-1)^{i}\binom{n-1}{i}(n-2 i-1)^{k} .
$$

The form of this summation suggests that a direct combinatorial explanation should exist, though none is known. Since $t(n, n-1)=\left|B_{n}\right|$, in looking for this

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direct explanation, it is of interest to know as much about the combinatorics of $\boldsymbol{B}_{\boldsymbol{n}}$ as possible. Accordingly, in this paper we give a second bijection between $A_{n}$ and $B_{n}$, one which is easy to implement and has a very simple proof.

Our method is to give direct combinatorial proofs that $\left|A_{n}\right|$ and $\left|B_{n}\right|$ satisfy the recurrence equation

$$
\begin{equation*}
c_{n}=\sum_{k=1}^{n-1}(n-k)\binom{n-2}{k-1} c_{k} c_{n-k}, \quad n \geqslant 2 ; \quad c_{i}=1 \tag{*}
\end{equation*}
$$

A comparison of these proofs yields the required bijection.
To prove that $\left|A_{n}\right|$ satisfies the recurrence relation, it is convenient to define a particular edge-deletion operation. Let $t$ be a labelled tree on $m$ vertices, with vertex labels $\left\{\beta_{1}, \ldots, \beta_{m}\right\}$, where $\beta_{1}<\cdots<\beta_{m}$ and $m \geqslant 2$. Consider the path in $t$ from vertex $\beta_{m}$ to $\beta_{m-1}$. Suppose that vertex $j$ is the vertex adjacent to vertex $\beta_{m}$ in this path. Now remove edge ( $\beta_{m}, j$ ) from $t$ to obtain two trees. One of these trees, call it $t^{\prime}$, contains vertex $\beta_{m}$, and the other tree, call it $t^{\prime \prime}$, contains vertex $\beta_{m-1}$. Define

$$
F(t)=\left(t^{\prime}, t^{\prime \prime}, j\right) .
$$

Proposition 1. $\left|A_{n}\right|, n \geqslant 1$ satisfies (*).
Proof. Consider an arbitrary $a \in A_{n}$. If $F(a)=\left(a^{\prime}, a^{\prime \prime}, j\right)$ then $a^{\prime}$ has vertex labels $\alpha \cup\{n\}$ for some $\alpha \subseteq\{1,2, \ldots, n-2\}$, and $a^{\prime \prime}$ has vertex labels $\bar{\alpha} \cup\{n-1\}$.
Clearly $|\alpha|=k-1$ for some $k=1, \ldots, n-1$. There are then $\binom{n-2}{k-1}$ choices for $\alpha,\left|A_{k}\right|$ choices for $a^{\prime},\left|A_{n-k}\right|$ choices for $a^{\prime \prime}$ and $n-k$ choices for $j$. This is reversible and the result follows.

Proposition 2. $\left|B_{n}\right|, n \geqslant 1$ satisfies (*).
Proof. Consider an arbitrary $\left(b_{1}, b_{2}, \ldots, b_{n-1}\right) \in B_{n}$, so $b_{1} b_{2} \cdots b_{n-1}=$ $(12 \cdots n)$. Now $b_{n-1}=(i, i+k)$ for some $k=1, \ldots, n-1$ and $i=1, \ldots, n-k$. Thus, multiplying from left to right,

$$
b_{1} b_{.} \cdots b_{n-2}=(12 \cdots n)(i, i+k)=c_{1} c_{2}
$$

where $c_{1}$ is the $k$-cycle $(i, i+1, \ldots, i+k-1)$ and $c_{z}$ is the $(n-k)$-cycle $(i+k, i+k+1, \ldots, n, 1, \ldots, i-1)$. Since $c_{1}$ and $c_{2}$ consist of disjoint elements, we have

$$
\prod_{j \in \alpha} b_{j}=c_{1}, \quad \prod_{l \in \bar{\alpha}} b_{l}=c_{2}
$$

for some $\alpha \subseteq\{1,2, \ldots, n-2\},|\alpha|=k-1$, where $b_{j}$ and $b_{l}$ commute for $j \in \alpha$, $l \in \bar{\alpha}$.

Thus there are $\binom{n-2}{k-1}$ choices for $\alpha,\left|B_{k}\right|$ choices for $\left(b_{\alpha_{1}}, \ldots, b_{\alpha_{k-1}}\right),\left|B_{n-k}\right|$ choices for ( $b_{\bar{\alpha}_{1}}, \ldots, b_{\bar{\alpha}_{n-k}}$ ) and $n-k$ choices for $i$, where $k=1, \ldots, n-1$. This is reversible and the result follows.

After comparing these combinatorial proofs, we are in a position to describe a recursive algorithm that provides the bijection.

Algorithm 3. The inputs are a labelled tree $t$ on $m$ vertices, with vertex labels $\left\{\beta_{1}, \ldots, \beta_{m}\right\}, \beta_{1}<\cdots<\beta_{m}, m \geqslant 1$, and an $m$-tuple of integers ( $l_{1}, \ldots, l_{m}$ ). If $m=1$ then STOP. Otherwise let $F(t)=\left(t^{\prime}, t^{\prime \prime}, j\right)$, let $k$ be the number of vertices in $t^{\prime}$, and suppose that $j$ is the $r$ th smallest label in $t^{\prime \prime}$. Then output the transposition $b_{\beta_{m-1}}=\left(l_{r}, l_{r+k}\right)$, perform Algorithm 3 with tree $t^{\prime}$ and $k$-tuple ( $l_{r}, l_{r+1}, \ldots, l_{+k-1}$ ) as inputs, and perform Algorithm 3 with tree $t^{\prime \prime}$ and ( $m-k$ )-tuple $\left(l_{r+k}, \ldots, l_{m}, l_{1}, \ldots, l_{r-1}\right)$ as inputs.

Theorem 4. For arbitrary $a \in A_{n}$, if we perform Algorithm 3 with tree $a$ and $n$-iuple $(1,2, \ldots, n)$ as inputs, we obtain $\left(b_{1}, \ldots, b_{n-1}\right) \in B_{n}$, and this is reversible, providing the required bijection.

Proof. Every $i=1, \ldots, n-1$ will appear as the second largest label in one of the sub-trees on which the algorithm is performed, so $b_{1}, \ldots, b_{n-1}$ are all defined.
Comparison of Proposition 1 and 2 shows that in the bijection, removal of an edge from a tree and the two resulting sub-trees correspond to factoring out the last transposition and the two resulting cycles. In applying Algorithm 3, we output this last transposition and recursively consider the sub-trees and the list of elements on the corresponding cycle.

We now give an example of the bijection, in which the operation of Algorithm 3 is represented by a binary tree. Each node represents one application of the algorithm, giving the input tree $t$ and iist, and output transposition. The edge to be deleted from $t$ is doubled. The left of spring treats the sub-tree $t^{\prime}$ and the right offspring treats the sub-tree $t^{\prime \prime}$. Nodes representing applications of the algorithm which STOP have been deleted.

## Example 5. See Fig. 1.

This tells us that corresponding to

$\in A_{7}$ there is $((6,7),(4,2),(6,1),(3,4),(2,5),(2,6)) \in B_{7}$, and indeed, to check, multiplying these transpositions in this order yields the cycle ( $12 \cdots 7$ ), as required.
It is straightforward to reverse this bijection, as follows. Given an $(n-1)$-tuple of transpositions $\left(b_{1}, \ldots, b_{n-1}\right)$ such that $b_{1} \cdots b_{n-1}=(12 \cdots n)$, we can re-


Fig. 1.
construct the binary tree created by Algorithm 3 in two passes, one down, and a second up.

On the first pass, for each vertex of the binary tree we identify the list of vertices in the cycle, the transposition, and the value of $r$. The list is passed from above, and the transposition is that of maximum subscript containing two elements in the list. We begin with the list $(1,2, \ldots, n)$ (and thus the transposition $b_{n-1}$ ) at the root vertex. The value of $r$ is straightforward at each vertex.

Simultaneously, we assign the labels $1,2, \ldots, n$ to the 'STOP' vertices of the binary tree as follows. We pass each vertex a label from above, which is assigned to the vertex if it is a STOP vertex. Otherwise, the label is passed to the left offspring of the vertex, and the subscript of the transposition corresponding to the vertex is passed to the right offspring of the vertex. We begin by passing the label ' $n$ ' to the root vertex.

The second pass up the tree is now used to join the labelled STOF vertices as specified iteratively by the valucs of $r$ at each internal vertex.

Finaily, this reversal of the bijection is illustrated by an example.


Fig. 2.

Example 6. Given $((3,4),(2,6),(3,6),(2,7),(1,2),(5,6)) \in B_{7}$, we complete the first pass down to give the binary tree in Fig. 2, in which the labels on the edges give the passed labels, and the circled vertices are the labelled STOP vertices.

The second pass up this binary tree immediately yields the corresponding tree in $A_{7}$ :


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