

AMATH 840, Lecture 08, Jan 31, 2022  
RIP  $m \gtrsim O(s \ln(\frac{n}{s}))$

Today ① Examples of matrices that satisfy the RIP cond.

② Sparse Optimization Methods Coherence  $m \gtrsim O(s^2)$

i) Proximal Methods: nondiff  $\|x\|_1$

ii) Lagrangian Methods: constrained  $\rightarrow$  unconstrained  
Image processing  $\rightarrow$   $\|\nabla x\|_1 \rightarrow$  ADMM

References ① A Mathematical Intro. to Compressive Sensing  
by Foucart and Raufut, Chapters 9, 12, 14.

② Convex Optimization, by Boyd & Vandenberghe

③ Slides from | EE 236C, by Vandenberghe, UCLA  
| CMSC 764, by T. Goldstein, Uni. of Maryland

# ① Recall RIP condition

Def (i) The restricted isometry constant  $\delta_s$  of a matrix  $A \in \mathbb{C}^{m \times n}$

is the smallest  $\delta \geq 0$  s.t

$$(1 - \delta) \|x\|_2^2 \leq \|Ax\|_2^2 \leq (1 + \delta) \|x\|_2^2, \quad \forall x: s\text{-sparse.}$$

(ii) The matrix  $A$  is said to satisfy the RIP condition if

$\delta_{ks} < \delta_{*}$ , where constants  $k$  and  $\delta_{*}$  depend on the

algorithms/models:

BP	IHT	HTP	OMP
$\delta_{2s} < \frac{4}{\sqrt{41}}$	$\delta_{6s} < \frac{1}{\sqrt{3}}$	$\delta_{6s} < \frac{1}{\sqrt{3}}$	$\delta_{13s} < \frac{1}{6}$

6.12

Theorem: If  $A \in \mathbb{C}^{m \times n}$  satisfy the RIP condition for BP

$$\delta_{2s} < \frac{4}{\sqrt{41}}, \text{ then}$$

① Every  $s$ -sparse vector  $w$  is exactly recovered from (BP).

$$\min_{z \in \mathbb{C}^n} \|z\|_1 \quad \text{st} \quad Az = Aw$$

② Moreover, the reconstruction is stable when sparsity is replaced by compressibility and the reconstruction is robust when measurement error occurs. More precisely, for any  $w \in \mathbb{C}^n$  and  $y \in \mathbb{C}^m$  with  $\|y - Aw\|_2 \leq \eta$ , a solution  $w^\#$  of BPDN

$$\min \|z\|_1 \quad \text{st} \quad \|y - Az\|_2 \leq \eta$$

approximates  $w$  with errors:  $\|w - w^\#\|_2 \leq \frac{c}{\sqrt{s}} \zeta_s(w)_1 + D\eta$

## ② Examples of matrices that satisfy RIP condition

Concentration Inequalities (Bernstein Ineq) / Covering number

Example 1: Gaussian Random Matrices: the entries of  $A$  are

independent standard Gaussian random variables

Example 2: Bernoulli Random Matrices: the entries of  $A$  are

independent Rademacher variables (taking values  $\pm 1$  with prob.  $\frac{1}{2}$ )

Example 3: Subgaussian Random Matrices: the entries of  $A$  are

independent mean-zero subgaussian random variables with

variance 1 and  $E(e^{sA_{j,k}}) \leq e^{\frac{\sigma^2 s^2}{2}}$ ,  $\forall s \in \mathbb{R}$ ,  
 $j \in [m], k \in [n]$ .

Here  $\sigma$  is a constant.

Remark: Gaussian and Bernoulli random matrices are subgaussian, since:

⊛ If  $X$  is  $\mathcal{N}(0, \sigma^2)$ , for any  $t > 0$ , we have:

$$\mathbb{E}(e^{sX}) = e^{\sigma^2 s^2 / 2}, \quad \forall s \in \mathbb{R}$$

⊛ If  $X$  is a Rademacher random variable, then

$$\begin{aligned} \mathbb{E}(e^{sX}) &= \frac{1}{2}(e^s + e^{-s}) \\ &= \frac{1}{2} \left( \sum_{k=0}^{\infty} \frac{s^k}{k!} + \sum_{k=0}^{\infty} \frac{(-s)^k}{k!} \right) = \sum_{k=0}^{\infty} \frac{s^{2k}}{(2k)!} \\ &\leq 1 + \sum_{k=1}^{\infty} \frac{s^{2k}}{2^k k!} = e^{s^2/2} \end{aligned}$$

Theorem 9.2 let  $A \in \mathbb{C}^{m \times n}$  subgaussian random matrix

Then there exists a constant  $C$  (depending on the subgaussian parameter  $\delta$ ) s.t. the restricted isometry constant of  $\frac{1}{\sqrt{m}} A$  satisfies  $\delta_s \leq \delta$  with probability at least  $1 - \varepsilon$ , provided

$$m \geq C \delta^{-2} \left[ s \ln\left(\frac{en}{s}\right) + \ln(2\varepsilon^{-1}) \right]$$

letting  $\varepsilon = 2 \exp(-\delta^2 m / (2C))$  yields :

$$m \geq 2C \delta^{-2} s \ln\left(\frac{en}{s}\right)$$

$$\begin{aligned} & \mathbb{E} \left( \left\| \frac{1}{\sqrt{m}} Ax \right\|_2^2 \right) \\ &= \mathbb{E} \left( \|x\|_2^2 \right) \end{aligned}$$

## Example 4: Structured Random Matrices / Random Sampling Matrices

Def: Bounded orthonormal system (BOS)

Let  $\mathcal{D} \subset \mathbb{R}^d$  be endowed with a probability measure  $\nu$ :

$\Phi := \{ \phi_1, \dots, \phi_n \}$ , where  $\phi_k: \mathcal{D} \rightarrow \mathbb{C} \quad \forall k \in [n]$  and

$$\textcircled{1} \quad \int_{\mathcal{D}} \phi_j(t) \overline{\phi_k(t)} d\nu(t) = \delta_{jk}$$

$$\textcircled{2} \quad \|\phi_j\|_{\infty} := \sup_{t \in \mathcal{D}} |\phi_j(t)| \leq \underset{\substack{\uparrow \\ \text{a constant}}}{K} \quad \forall j \in [n]$$

Then  $\Phi$  is called a bounded orthonormal system with constant  $K$

Def (cont'd) let  $t_1, \dots, t_m \in \mathcal{D}$  be some sampling points

The matrix  $A = (A_{l,k} = \Phi_k(t_l)) \in \mathbb{C}^{m \times n}$   
 is called a sampling matrix.

$$A = \begin{pmatrix} \Phi_1(t_1) & \Phi_2(t_1) & \dots & \Phi_n(t_1) \\ \vdots & \vdots & & \vdots \\ \Phi_1(t_m) & \Phi_2(t_m) & \dots & \Phi_n(t_m) \end{pmatrix}$$

### Examples of BOS

$$k = \{ -q, -q+1, \dots, 0, \dots, q \}$$

① Trigonometric polynomials:  $\mathcal{D} = [0, 1]$ ,  $\Phi_k(t) = e^{2\pi i k t}$

② Real Trigonometric polynomials,  $\mathcal{D} = [0, 1]$

$$\Phi_0(t) = 1, \quad \Phi_{2k}(t) = \sqrt{2} \cos(2\pi k t), \quad \Phi_{2k-1}(t) = \sqrt{2} \sin(2\pi k t)$$

$$k \geq 1.$$



③ Partial discrete Fourier transform

$$F \in \mathbb{C}^{n \times n}, \quad F_{l,k} = \frac{1}{\sqrt{n}} e^{2\pi i (l-1)(k-1)/n}$$

$l, k \in [n]$

The normalized columns of  $F$ ,  $\sqrt{n} F_k \in \mathbb{C}^n$  form an orthonormal system w.r.t. the discrete uniform measure on  $[n]$

④ Legendre polynomials  $\mathcal{D} = [-1, 1]$

$$P_0(x) = 1, \quad P_1(x) = x, \quad P_2(x) = \frac{1}{2}(3x^2 - 1), \dots$$

↓ normalized

Setting Consider functions of the form

$$f(t) = \sum_{k=1}^n w_k \Phi_k(t)$$

$$\left\{ \begin{array}{l} f: \mathcal{D} \subset \mathbb{R} \rightarrow \mathbb{C} \\ \Phi_k: \mathcal{D} \rightarrow \mathbb{C} \end{array} \right\} \text{ BOS}$$

Suppose we are given sample values  $(t_J, f(t_J))_{J=1}^m \subset \mathcal{D} \times \mathbb{C}$ .

$$y_l = f(t_l) = \sum_{k=1}^n w_k \Phi_k(t_l)$$

$$\begin{pmatrix} y_1 \\ \vdots \\ y_m \end{pmatrix} = A \begin{pmatrix} w_1 \\ \vdots \\ w_n \end{pmatrix}$$

where  $A \in \mathbb{C}^{m \times n}$

$$(A_{l,k} = \Phi_k(t_l))$$

$$l \in [m], k \in [n]$$

$$\vec{y} = A \vec{w}$$

**Goal** : Given sample values  $\vec{y}$  and sampling matrix  $A$

find  $\vec{w} \in \mathbb{C}^n$  . ( $\Leftrightarrow$  find function  $f$ )

If  $f$  is  $s$ -sparse w.r.t  $\{\phi_1, \dots, \phi_n\}$ , find  $s$ -sparse vector  $w \in \mathbb{C}^n$  s.t  $\vec{y} = A \vec{w}$ .

**Remark**: Assume the sampling points  $t_1, \dots, t_m$  are selected independently at random according to the prob. measure  $\nu$ . Then the sampling matrix  $A$  is called a random sampling matrix associated to a BAS with constant  $K \geq 1$ .

Theorem 12.22 let  $w \in \mathbb{C}^n$  and let  $A \in \mathbb{C}^{m \times n}$

be the random sampling matrix associated to a BAS with constant  $K \geq 1$ . For  $y = Aw$  with  $\|w\|_2 \leq \sqrt{m}\eta$  for some  $\eta \geq 0$ , let  $w^\#$  be a solution of

$$\min_{z \in \mathbb{C}^n} \|z\|_1 \quad \text{s.t.} \quad \|Az - y\|_2 \leq \sqrt{m}\eta$$

depend on the  $\nu_{\max}$  degree of Legendre poly.

If  $m \geq C K^2 S \ln(n) \ln(\epsilon^{-1})$ ,

then with probability at least  $1 - \epsilon$ , the reconstruction error satisfies

$$\|w - w^\#\|_2 \leq C_1 \sigma_S(w)_1 + C_2 \sqrt{S}\eta$$

$C_1, C_2$  are global constants

**Theorem 12.32** : let  $A \in \mathbb{C}^{m \times n}$  be the random sampling matrix associated to a BOS with constant  $\kappa \geq 1$ . Under some conditions on  $m$  ( $\Theta(\ln(s))$ ), the matrix  $A$  satisfies the RIP condition with high prob.

↓ Robust

Stable prob measure

build →

$\{\phi_k\}_{k=1}^n$

BOS

Drawback Depend ↻

$\{t_k\}_{k=1}^m$

↻ randomly

↑

Input data

Empirical Distributions

soln ①

sem ② Transfer randomness<sup>of</sup> input data  $\rightarrow$  randomness of the Random

$$A = \left( \sigma \langle t_k, \overset{\downarrow \text{Random weight}}{\omega_k} \rangle \text{controllable} \right) \text{Feature}$$

A also satisfies RIP-like conditions