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## Problem statement

Let $f \in \mathbb{Q}\left[X_{1}, \ldots, X_{n}\right]$ be squarefree with total degree $d$ and $V(f) \cap \mathbb{R}^{n}$ smooth and compact. Let $A \in \mathbb{Q}^{n^{2}}$ be a linear change of variables that we apply to $f$ obtaining $f^{A}(x)=f(A x)$. We provide bit size estimates for computing the critical points in generic coordinates $x \in V\left(f^{A}\right) \cap \mathbb{R}^{n}$ of the projection $\Pi_{1}:\left(x_{1}, \ldots, x_{n}\right) \in \mathbb{C}^{n} \mapsto x_{1} \in \mathbb{C}$. The critical points are defined by the vanishing of $f^{A}, \frac{\partial f^{A}}{\partial X_{2}}, \ldots, \frac{\partial f^{A}}{\partial X_{n}}$.

## An example of the critical points

Let $f=X_{1}^{2}+X_{2}^{2}+X_{3}^{2}-1$ and consider
$V\left(X_{1}^{2}+X_{2}^{2}+X_{3}^{2}-1\right) \subset \mathbb{C}^{3}$.
The real critical points of $\Pi_{1}$ are defined by
$X_{1}^{2}+X_{2}^{2}+X_{3}^{2}-1=X_{2}=X_{3}=0$
and thus define the algebraic set
$V\left(f, X_{2}, X_{3}\right)=V\left(f, \frac{\partial f}{\partial X_{2}}, \frac{\partial f}{\partial X_{3}}\right)$.


## Applications

Computing the real critical points of $\Pi_{1}$ in generic coordinates is an important step in computing a roadmap of a semi-algebraic set, as for instance in [1]. Roadmaps are used for deciding connectivity properties in semi-algebraic sets.
The set of real critical points of $\Pi_{1}$ in generic coordinates is finite and gives one point on each connected component of $V(f) \cap \mathbb{R}^{n}$, assuming $V(f) \cap \mathbb{R}^{n}$ is smooth and compact. Hence, computing the critical points determines an upper bound on the number of connected components and determines whether real solutions exist $[2,3]$.

## Generic coordinates

When $A$ is sufficiently generic, the Jacobian of the system of polynomials $f^{A}, \frac{\partial f^{A}}{\partial X_{2}}, \ldots, \frac{\partial f^{A}}{\partial X_{n}}$ will have full rank at all $x \in V\left(f^{A}\right) \cap \mathbb{R}^{n}$. It then follows by the Jacobian criterion [4, Theorem 16.19] that the set of critical points $V\left(f^{A}, \frac{\partial f^{A}}{\partial X_{2}}, \ldots, \frac{\partial f^{A}}{\partial X_{n}}\right)$ is finite, and the ideal $\left\langle f^{A}, \frac{\partial f^{A}}{\partial X_{2}}, \ldots, \frac{\partial f^{A}}{\partial X_{n}}\right\rangle$ is radical. In this case we say $A$ is good. Otherwise we say that $A$ is bad.

Theorem 1. The bad changes of variables are contained in a hypersurface $\Delta \subset \mathbb{C}^{n^{2}}$ of degree at most $(d+1)^{n}$.
Corollary 2. Fix $S \subset \mathbb{Q}$ with $|S| \geq \epsilon^{-1}(d+1)^{n}$ and $\epsilon>0$. Then for
$A$ in $S^{n^{2}}$ chosen randomly, $\operatorname{Pr}[A$ is good $] \geq 1-\epsilon$.

## Bit complexity

Theorem 3. Suppose that $f$ satisfies $\operatorname{deg} f \leq d$, height $f \leq s$, with $f$ given by a straight-line program $\Gamma$ of size $L$ with integer constants of height at most $b$. There exists a randomized algorithm that takes $\Gamma, d$, and $s$ as input and produces a zero-dimensional parameterization of the critical points

$$
V\left(f^{A}, \frac{\partial f^{A}}{\partial X_{2}}, \ldots, \frac{\partial f^{A}}{\partial X_{n}}\right),
$$

with probability at least $9 / 16$, where $A \in \mathbb{Q}^{n^{2}}$ is a linear change of variables chosen randomly by the algorithm. Otherwise the algorithm either produces a subset of the critical points or FAIL. In any case, the algorithm uses
$O^{\sim}\left(L b+d^{2 n}(s+d)(L+d)\right)$
boolean operations.
Running the algorithm $k$ times gives a list of outputs among which the highest cardinality set includes all critical points with probability at least $1-(7 / 16)^{k}$.

## Transversality

We prove Theorem 1 by developing a quantitative extension of Thom's weak transversality [1, Proposition B.3], specialized to the particular case of transversality to a point which can be rephrased entirely in terms of critical and regular values: $f$ is transverse to a point $\{a\}$ if and only if $\{a\}$ is a regular value of $f$, where regular /critical values are images of respectively regular /critical points. Let $\Phi: \mathbb{C}^{n} \times \mathbb{C}^{\tilde{d}} \rightarrow \mathbb{C}^{m}$ be a polynomial mapping where $n, \tilde{d}$, and $m$ are positive integers. Assume the total degree of $\Phi$ is bounded by an integer $d$. For $A \in \mathbb{C}^{\widetilde{d}}$, let $\Phi_{A}: \mathbb{C}^{n} \rightarrow \mathbb{C}^{m}$ be the induced mapping $x \mapsto \Phi(x, A)$.

Theorem 4. Suppose that 0 is a regular value of $\Phi$. Then there exists a hypersurface $\Delta \subset \mathbb{C}^{\widehat{d}}$ of degree bounded by $(d+1)^{n}$ for which, if $A \in \mathbb{C}^{\widetilde{d}}-\Delta$ then 0 is a regular value of $\Phi_{A}$.
We use Theorem 4 to show that $\Delta$ contains that bad changes of variables.

## An example of Thom's weak transversality

Let $\Phi\left(x, \vartheta_{1}, \vartheta_{2}\right)=x^{2}+\vartheta_{1}^{2}+\vartheta_{2}^{2}-1$ so that $\Phi^{-1}(0)$ is a smooth variety which implies that $\Phi$ is trans-
verse to 0 . Now, Thom's theorem tells us that, verse to 0 . Now, Thom s theorem tells us that,
for a generic choice of $A=\left[A_{1}, A_{2}\right]$, the polyfor a generic choice of $A=\left[A_{1}, A_{2}\right]$, the poly-
nomial $\Phi\left(x, A_{1}, A_{2}\right)=\Phi_{A}(x)$ is transverse to 0 . $\Phi_{A}(x)$ is univariate and transverse to zero when $\Phi_{A} \in \Phi_{A}^{-1}(0) \Rightarrow \operatorname{grad}_{x} \Phi_{A} \neq 0$. Hence, when square$x \in \Phi_{A}^{1}(0) \Rightarrow \operatorname{grad}_{x} \Phi_{A} \neq 0$. Hence, when square-
free. The green points correspond to a generic choice of $A_{1}$ and $A_{2}$ whereas the red point correchoice of $A_{1}$ and $A_{2}$ whereas the red point corre-
sponds to a double root which is an unlucky choice.

## The bad $\vartheta$ are contained in $\triangle$

We let $\Phi$ be the mapping $(x, \vartheta) \mapsto\left(f^{\vartheta}(x), \frac{\partial f^{\vartheta}}{\partial X_{2}}(x), \ldots, \frac{\partial f^{\vartheta}}{\partial X_{n}}(x)\right)$ so that $\Phi^{-1}(0)$ defines the critical points in generic coordinates.
We show that 0 is a regular value of $\Phi$ and thus, by Theorem 4, a hypersurface $\Delta \subset \mathbb{C}^{n^{2}}$ exists with the property that, if $A \in \mathbb{C}^{n^{2}}-\Delta$, then 0 is a regular value of $\Phi_{A}$, which means that jac ${ }_{x} \Phi_{A}$ has full rank for all $x \in V\left(f^{A}\right)$ and therefore $A$ is good. Hence, $\Delta$ contains the bad changes of variables

## Proving Theorem 4

Put $X=\Phi^{-1}(0)$ and let $\pi:(x, \vartheta) \in \mathbb{C}^{n} \times \mathbb{C}^{\widetilde{d}} \mapsto \vartheta \in \mathbb{C}^{\tilde{d}}$.
The classical proof of Thom's weak transversality goes by showing that if $\vartheta \in \mathbb{C}^{\tilde{d}}$ is such that 0 is not a regular value of $\Phi_{\vartheta}$ then $\vartheta$ is a critical value of $\left.\pi\right|_{X}$. It then follows from Sard's lemma [1] that the critical values of $\left.\pi\right|_{X}$ are contained in a hypersurface $\Delta \subset \mathbb{C}^{\tilde{d}}$.
We first bound the degree of an algebraic set $\Delta^{\prime}$ containing the critical points $(x, \vartheta)$ of $\left.\pi\right|_{X}$. We show that $\operatorname{deg} \Delta \leq \operatorname{deg} \Delta^{\prime} \leq(d+1)^{n}$.
Let $M=\left[\begin{array}{c}\operatorname{jac}\left(\left.\pi\right|_{X}\right) \\ \operatorname{jac}(\Phi)\end{array}\right]$. We prove that $M(x, \vartheta)$ has full rank $\widetilde{d}+m$ if and only if $(x, \vartheta)$ is a regular point of $\left.\pi\right|_{X}$. Hence, $\Delta^{\prime}$ is defined by the minors of $M(x, \vartheta)$ of order $\widetilde{d}+m$.
Next, we observe that
$M(x, \vartheta)=\left[\begin{array}{c}\operatorname{jac}_{(x, \vartheta)}\left(\left.\pi\right|_{X}\right) \\ \operatorname{jac}_{(x, \vartheta)}(\Phi)\end{array}\right]=\left[\begin{array}{cc}\mathbf{0}_{\widetilde{d} \times n} & \mathbf{I}_{\widetilde{d}} \\ \operatorname{jac}_{(x, \vartheta)}(\Phi)[; 1, n] & \operatorname{jac}_{(x, \vartheta)}(\Phi)[; n+1, \widetilde{d}]\end{array}\right]$,
and we show that $\Delta^{\prime}$ is also defined by the minors of the sub-matrix $J=$ $\operatorname{jac}_{(x, \vartheta)}(\Phi)[; 1, n]$ of order $m$. We then introduce Lagrange multipliers $L=\left(L_{1}, \ldots, L_{m}\right)$ and let $G_{1}, \ldots, G_{n}$ be the equations defined by the Lagrange system

$$
\left[L_{1}, \ldots, L_{m}\right] J(x, \vartheta)=\left[G_{1}(x, \vartheta, L), \ldots, G_{n}(x, \vartheta, L)\right]
$$

We let $\mathfrak{Z}$ denote the algebraic set defined by the vanishing of $G_{1}, \ldots, G_{n}$, and show that $\operatorname{deg} \Delta \leq \operatorname{deg} \Delta^{\prime} \leq \operatorname{deg} \mathfrak{Z} \leq(d+1)^{n}$.

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